

Transversal Helly numbers,  
pinning theorems  
and projection of simplicial complexes

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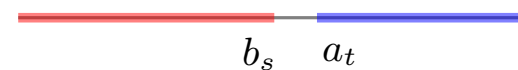
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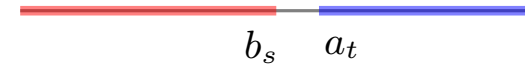


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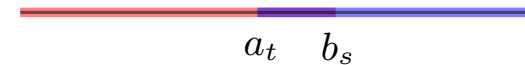
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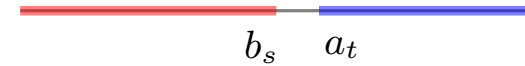


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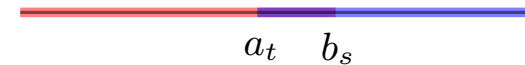
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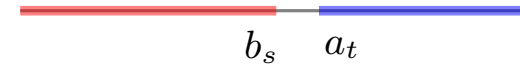
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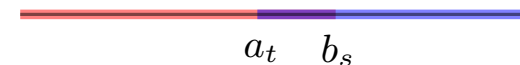
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This is what **Helly numbers** capture:

situations where empty intersection of **arbitrary large families** can be traced back to **constant-size** sub-families.

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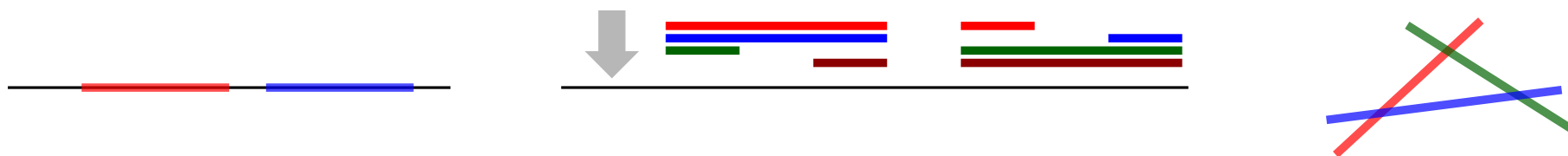
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In general, Helly numbers may be arbitrarily large...

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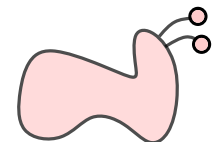
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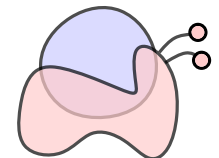
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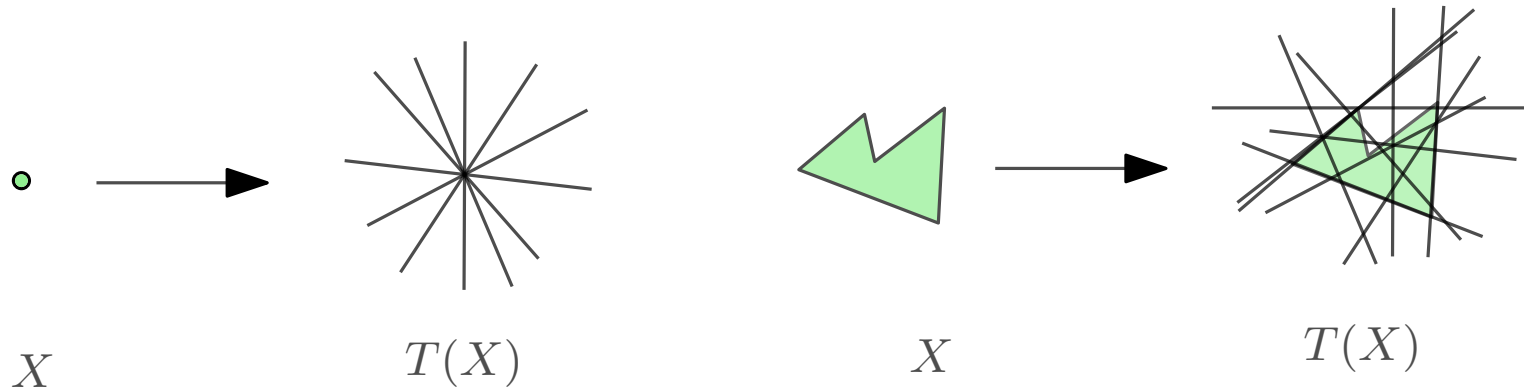
Helly numbers also arise naturally in discrete geometry, topology, algebra...



This presentation discusses Helly numbers of **sets of line transversals**.

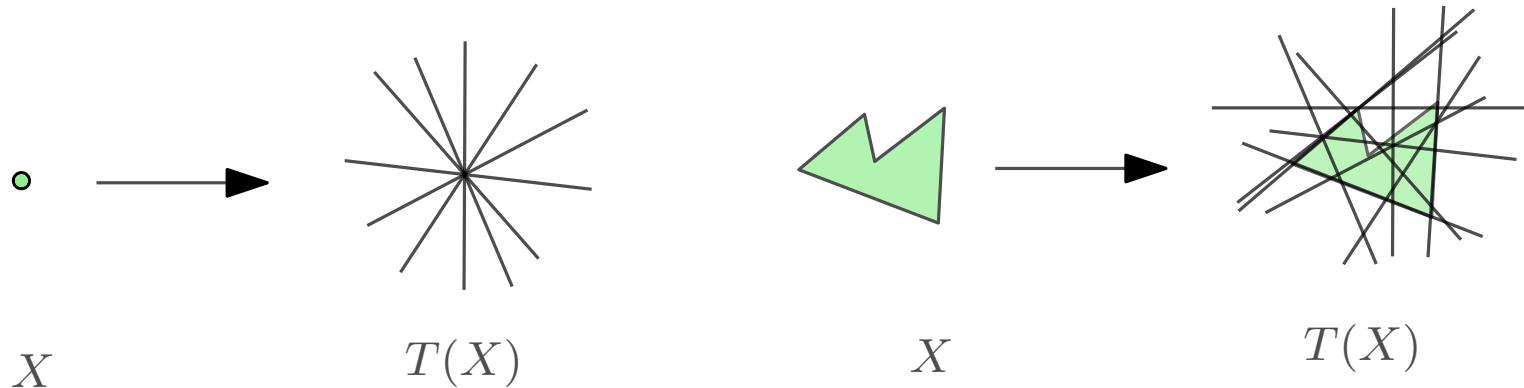
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**Conjecture (Danzer, 1957).** For any  $d \geq 2$  there exists  $\mathcal{H}_d \in \mathbb{N}$  such that the following holds: for any  $n \in \mathbb{N}$ , for any family  $\{B_1, \dots, B_n\}$  of pairwise disjoint unit balls in  $\mathbb{R}^d$ , the Helly number of  $\{T(B_1), \dots, T(B_n)\}$  is at most  $\mathcal{H}_d$ .

"if any  $\mathcal{H}_d$  balls in a family can be stabbed by a line, the whole family can be stabbed by one and the same line."

In this presentation...

An overview of a proof of Danzer's conjecture

Show how "everything fits together" (high-level).

A follow-up: a new homological conditions for bounding Helly numbers

Show a "nice machinery in motion" (more in-depth).

In this presentation...

Quick panorama of my research activity of these last years

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Some research perspectives

# Panorama of research activities

At the interface between computer science and mathematics.

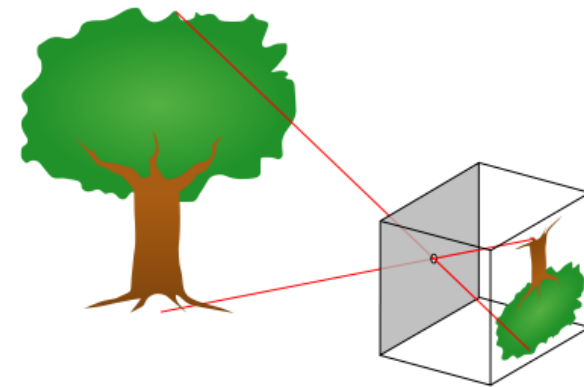
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How can line geometry help understand light propagation and models of imaging systems?

Lehtinen et al. 2008



Wikipedia

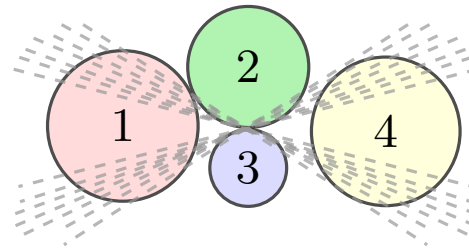
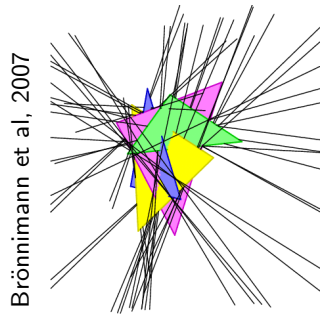


- ★ Shadow boundaries & **topological** visual event surfaces.
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[PhD Demouth], [Msc Batog], [PhD Batog], [Msc Jang]  
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# Geometric transversal theory

How does the geometry of an object determine the structure of its geometric transversals?



- ★ Geometric permutations & topology of sets of line transversals
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- ★ Pinning theorems

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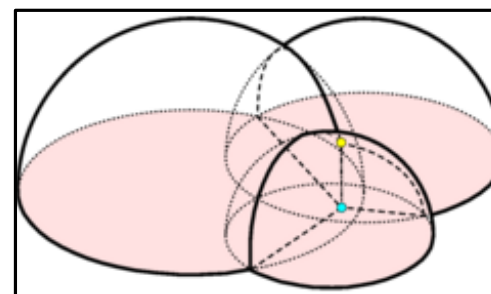
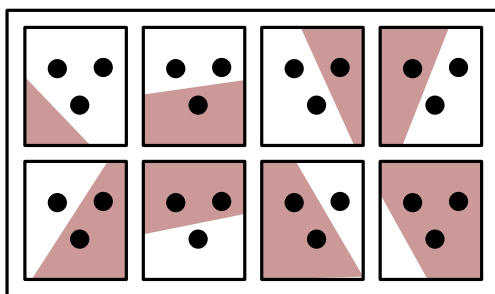
[SoCG 2005], [SoCG 2007], [DCG]<sup>x4</sup>, [IJM]

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## Combinatorics of geometric structures

How does the geometry shape the combinatorial structure underlying geometric objects?



- ★ Helly numbers for approximate covering
- ★ asymptotic of shatter functions of hypergraphs and families of permutations
- ★ Helly numbers from generalized nerve theorems

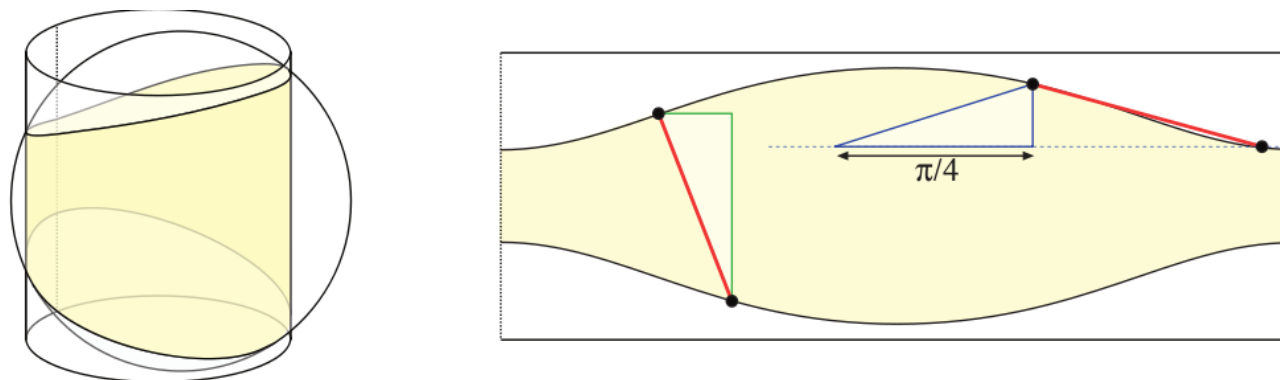
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How can probabilistic analysis help refine overly pessimistic worst-case analysis?



- ★ Delaunay triangulation of random samples of a surface
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[ANR "Projet Blanc" 2012-2016 with stochastic geometers]

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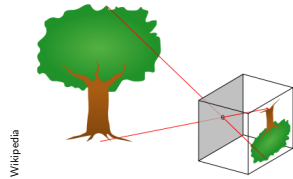
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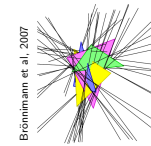
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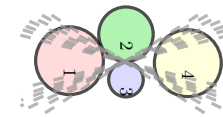
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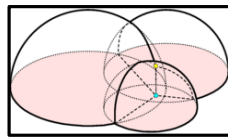
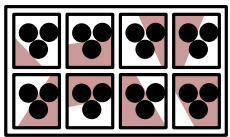


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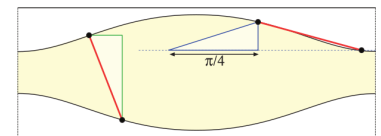
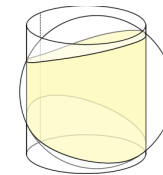
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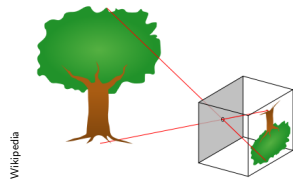
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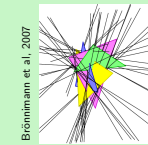
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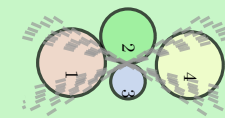
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Brömmann et al., 2007

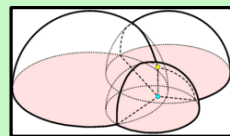
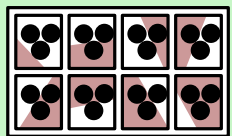


- ★ Geometric permutations & topology of sets of line transversals
- ★ Proof of Danzer's conjecture
- ★ Pinning theorems

[~Msc Koenig], ([PhD Ha])  
 [SoCG 2005], [SoCG 2007], [DCG]x4, [IJM]  
 Collaboration with O. Cheong, A. Holmsen, S. Petitjean, C. Borcea, B. Aronov, G. Rote

## Combinatorics of geometric structures

How does the geometry shape the combinatorial structure underlying geometric objects?



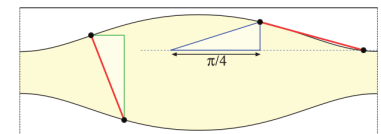
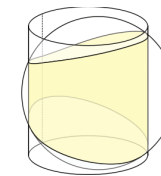
- ★ Helly numbers for approximate covering
- ★ asymptotic of shatter functions of hypergraphs and families of permutations
- ★ Helly numbers from generalized nerve theorems

[PhD Demouth]  
 [SoCG 2008]

13-2 Collaboration with O. Devillers, M. Glisse, O. Cheong, C. Nicaud, E. Colin de Verdière, G. Ginot

## Complexity of random geometric structures

How can probabilistic analysis help refine overly pessimistic worst-case analysis?



- ★ Delaunay triangulation of random samples of a surface
- ★ Smoothed complexity of convex hulls

[ANR "Projet Blanc" 2012-2016 with stochastic geometers]  
 [SODA 2008]  
 Collaboration with O. Devillers, J. Erickson, M. Glisse, D. Attali

# An overview of the proof of Danzer's conjecture

Where does Danzer's conjecture come from?

Helly's theorem reformulates as

*"The Helly number of sets of **point transversals** to convex sets is at most  $d + 1$ ."*

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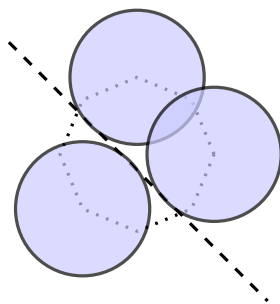
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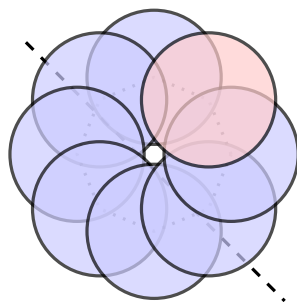
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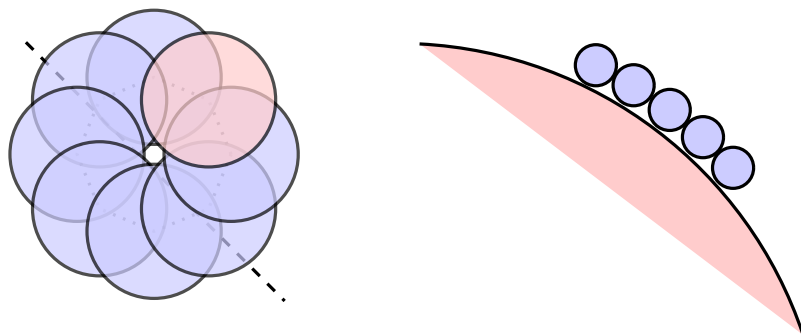
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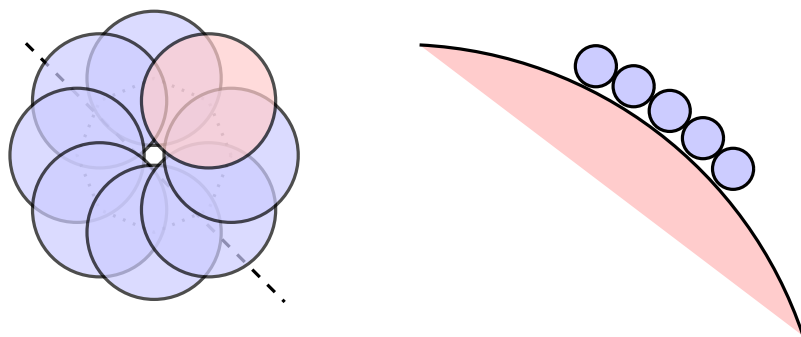
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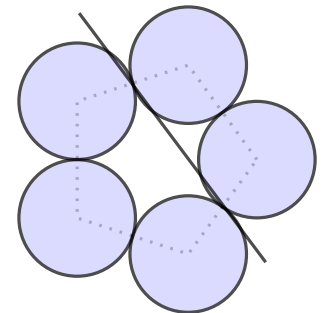
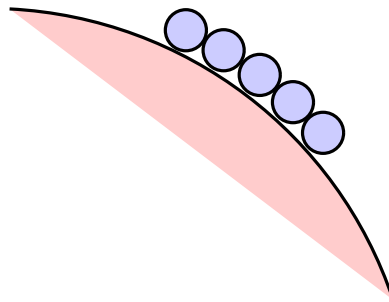
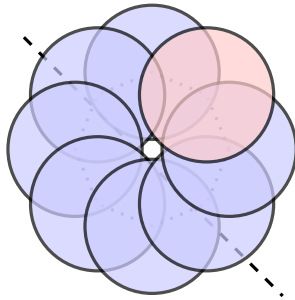
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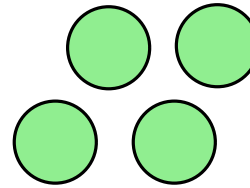


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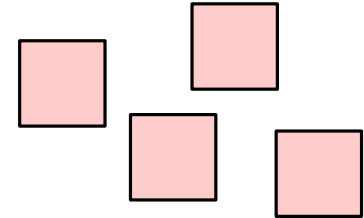
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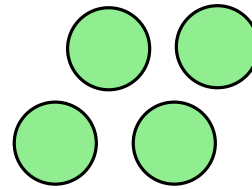


to

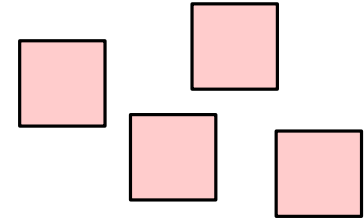


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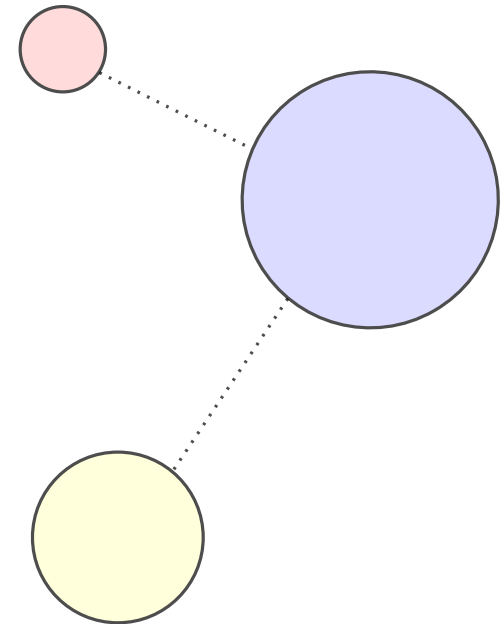
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## Results on Danzer's conjecture up to 2004. (1/2)

★ **True** for collections of thinly distributed balls in  $\mathbb{R}^d$ .

[Hadwiger 1959] and [Grünbaum 1960]

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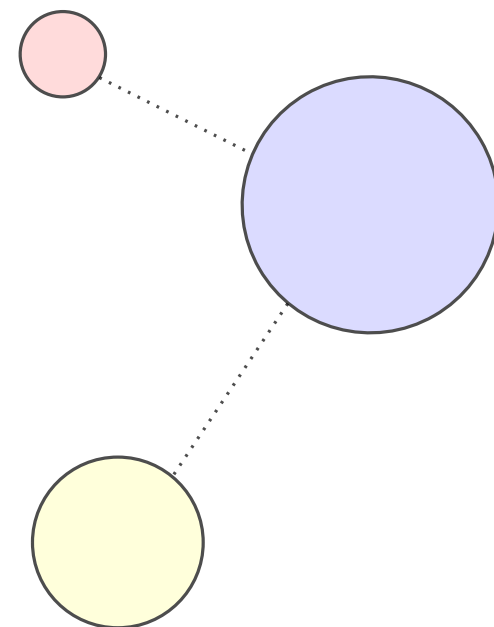
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Given  $\mathcal{F} = \{B_1, \dots, B_n\}$  put  $T(\mathcal{F}) = \bigcap_i T(B_i)$

Let  $\pi$  map each line to its **orientation** in  $\mathbb{RP}^{d-1}$

$\mathcal{K}(\mathcal{F}) = \pi(T(\mathcal{F}))$  is the **cone of directions**.

If  $\mathcal{F}$  is thinly distributed then  $\mathcal{K}(\mathcal{F})$  is **convex** (Hadwiger).



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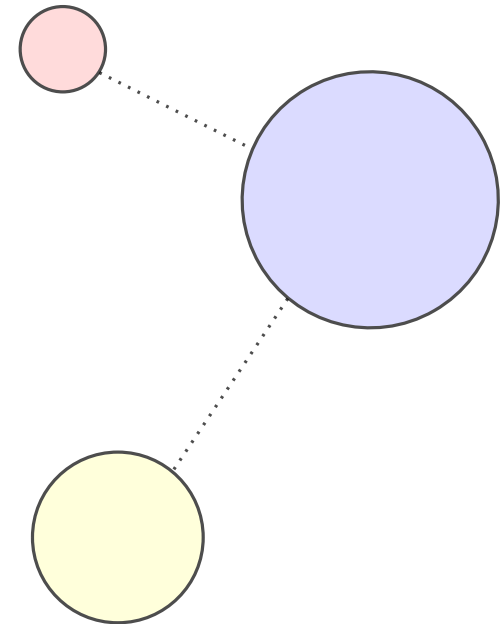
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Thus  $\{T(B_1), \dots, T(B_n)\}$  form a good cover (Grünbaum).

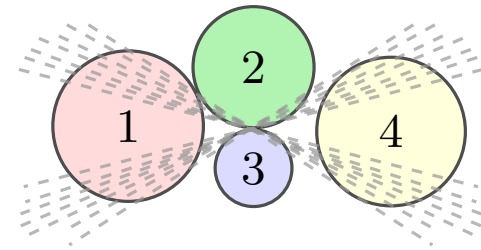
Helly's topological theorem  $\Rightarrow$  Helly number of  $\{T(B_1), \dots, T(B_n)\} \leq 2d - 1$ .



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[Holmsen-Katchalski-Lewis 2003]

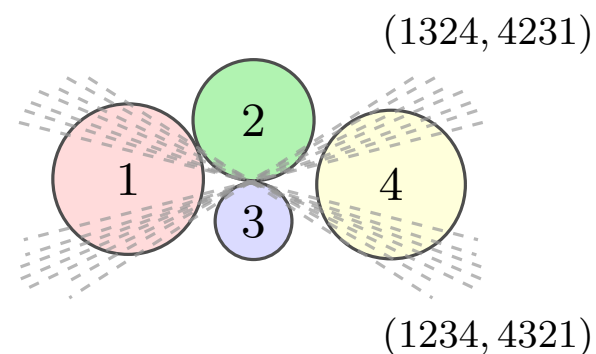


If  $\mathcal{F} = \{B_1, \dots, B_n\}$  is a family of disjoint unit balls in  $\mathbb{R}^3$  then each *connected component* of  $\mathcal{K}(\mathcal{F})$  is **convex**.

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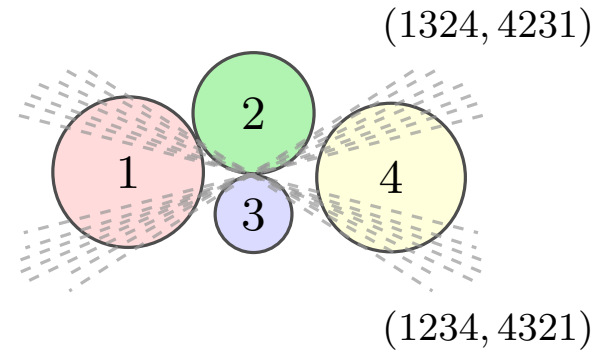
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Combinatorial restrictions on geometric permutations of disjoint unit balls

$\Rightarrow$  Helly number of  $\{T(B_1), \dots, T(B_n)\} \leq 46$ .



## Ingredients of our proof

- ★ **Generalized** the convexity structure of **cones of directions**.

[SoCG 2007] [DCG]

Joint work with C. Borcea and S. Petitjean

- ★ **Clarified** the structure of sets of geometric permutations.

[CGTA]

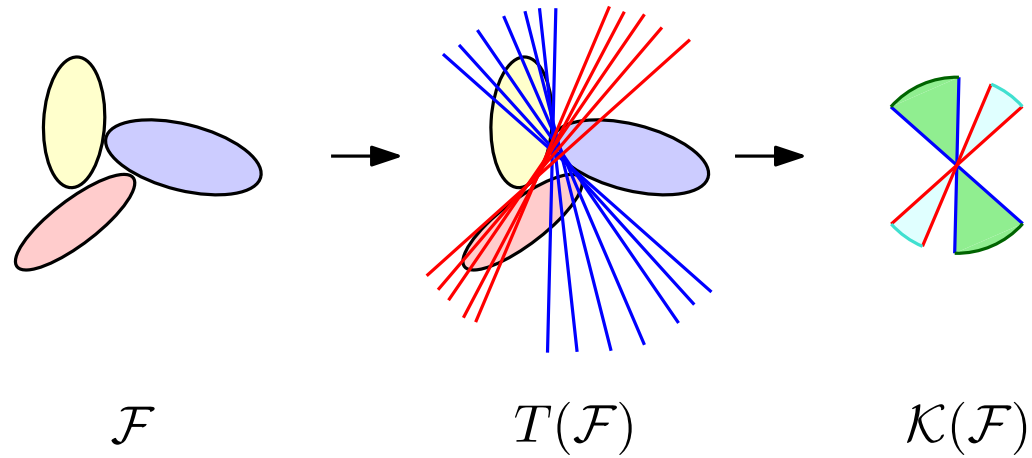
Joint work with O. Cheong and H.S. Na

- ★ Added a new ingredient: **pinning theorems**.

[SoCG 2005] [DCG]

Joint work with O. Cheong and A. Holmsen

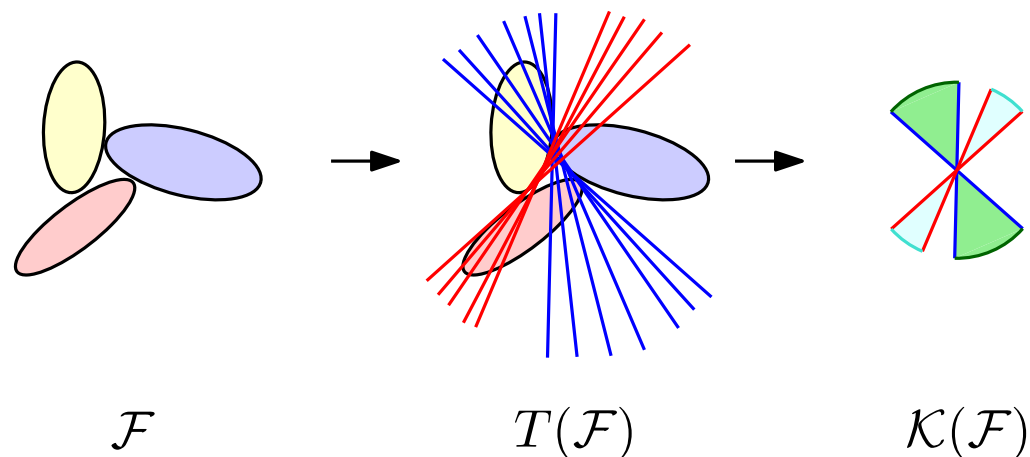
What are cones of directions?



[SoCG 2007] [DCG]

Joint work with C. Borcea and S. Petitjean

What are cones of directions?



How to prove that cones of directions are convex?

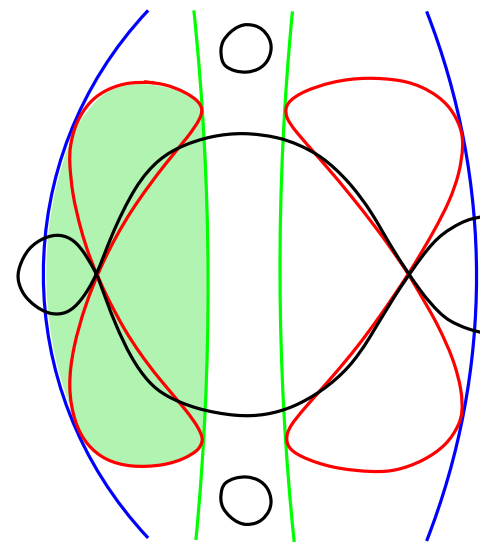
★ we analyzed the geometry of the curves bounding  $\mathcal{K}(\mathcal{F})$

Arcs of **conics** and **sextics**.

Track the **inflexion/singular** points.

Characterization of the arcs of sextic on  $\partial\mathcal{K}(\mathcal{F})$

No inflexion/singular point on  $\partial\mathcal{K}(\mathcal{F})$



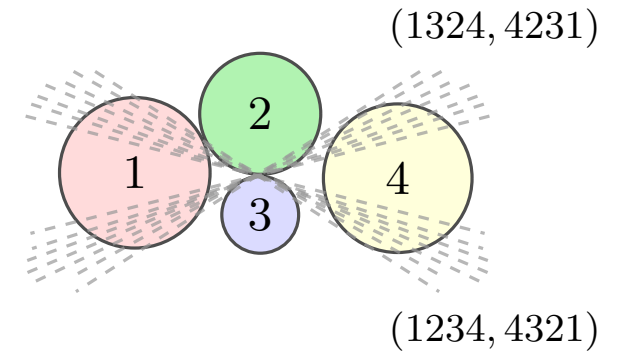
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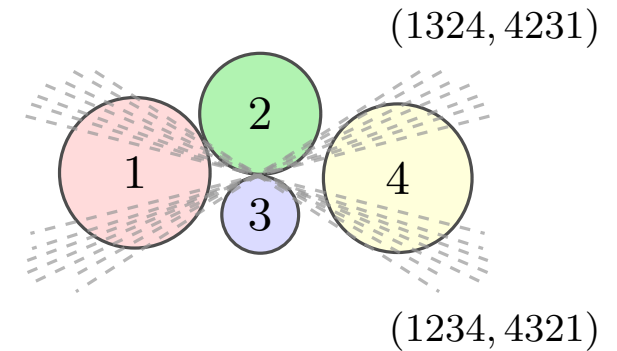
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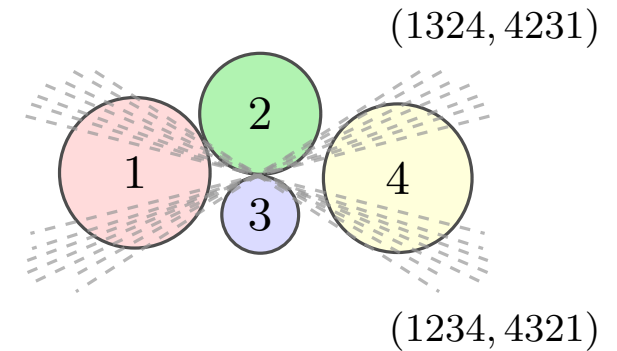
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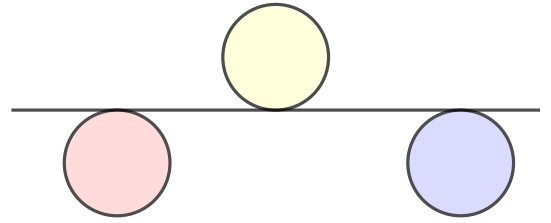
★ geometry  $\Rightarrow$  excluded **pairs** of patterns (in the Stanley-Wilf sense).

★ combinatorial extrapolation.

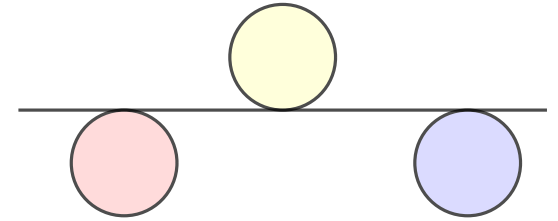
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Joint work with O. Cheong and H.S. Na

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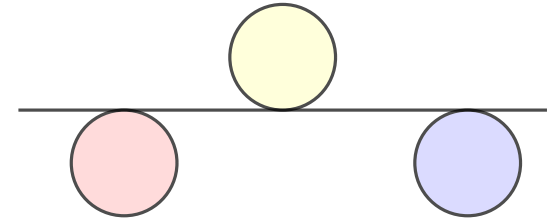
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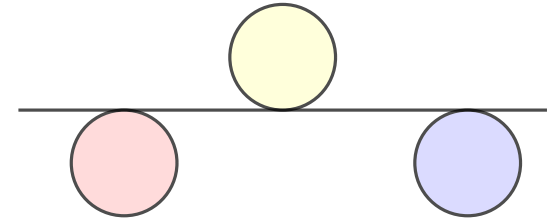
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[SoCG 2005] [DCG]

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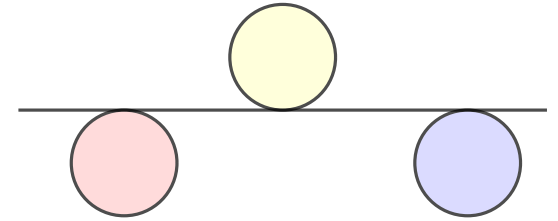
**Pinning theorem**, a **local** analogue of a Helly number.

★ Argue that **locally** near  $\ell$  the  $T(B_i)$  form a **good cover**.

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**Pinning theorem**, a **local** analogue of a Helly number.

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- ★ Conclude using Helly's topological theorem.

[SoCG 2005] [DCG]

Joint work with O. Cheong and A. Holmsen

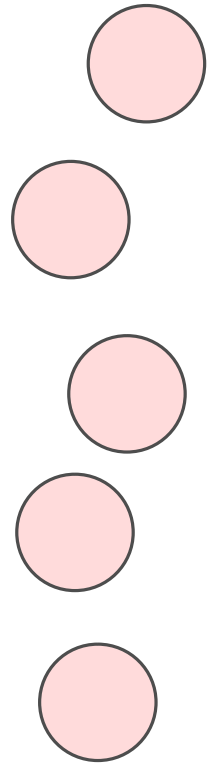
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[DCG]

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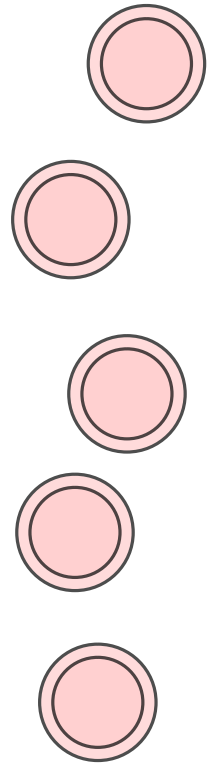
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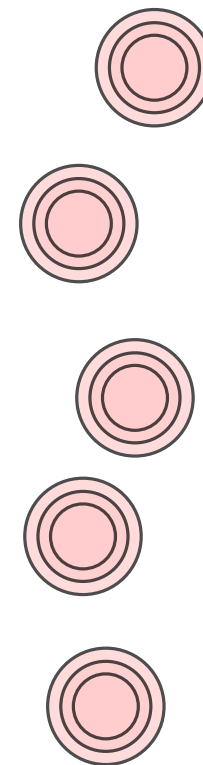
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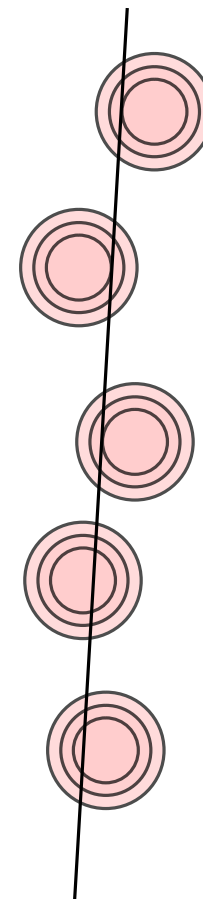


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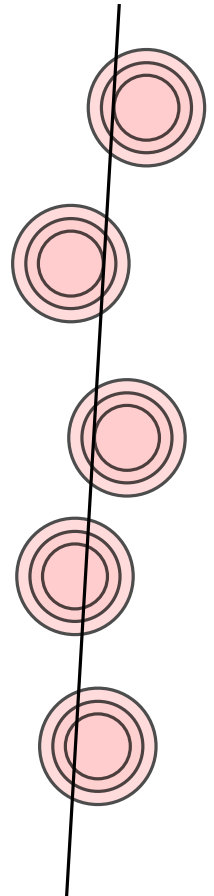
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pinning theorem and considerations on geometric permutations

$\Rightarrow \ell$  is **pinned** and **the only transversal** of  $G^* \subseteq G$  of size at most  $4d - 2$ .

$\forall B \in \mathcal{F}, G^* \cup \{B\}$  **still** has a transversal; it can only be  $\ell$ .



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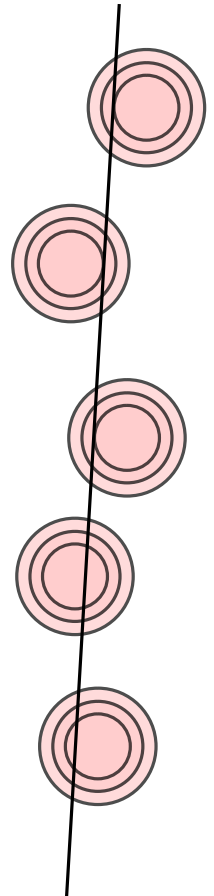
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★ Or  $G$  has **two** line transversals  $\ell_1$  and  $\ell_2$ , which it pins.

Similar (but slightly more complicated) arguments.



[DCG]

Joint work with O. Cheong, A. Holmsen and S. Petitjean

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for disjoint balls in  $\mathbb{R}^d$

$\implies$   
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Pinning theorems for other shapes  
(polytopes and ovaloids).

**Stable** pinning

- ★ tight lower bound for the pinning theorem
- ★ lower bound of  $2d - 1$  for the Helly number
- ★ relation to transversality of intersection

Application: computing a smallest enclosing cylinder (1/3)

**Smallest enclosing cylinder (SEC) problem:** given  $n$  points in  $\mathbb{R}^d$ , compute the cylinder with minimum radius that contains all the points.

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In the **Real RAM** model:

For  $d = 2$  the worst-case complexity of SEC is  $\Theta(n \log n)$ .

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In the Turing machine model:

SEC is NP-hard when the dimension  $d$  is part of the input.

[Megiddo 1990]

## Application: computing a smallest enclosing cylinder (2/3)

Let  $P$  be a set of  $n$  points in  $\mathbb{R}^d$ .

$$\text{Let } \phi : \begin{cases} 2^P & \rightarrow \mathbb{R} \times \mathbb{N} \\ Q & \mapsto (r_Q, n_Q) \text{ where } \begin{cases} r_Q = \text{radius of the SEC of } Q \\ n_Q = \# \text{enclosing cylinders of } Q \text{ of radius } r_Q \end{cases} \end{cases}$$



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**Proposition.**  $(P, \phi)$  is a LP-type problem.

The **combinatorial dimension** of  $(P, \phi)$  is the **maximum** size of a subset  $Q \subseteq P$  such that  $\forall x \in Q, \phi(Q \setminus \{x\}) \neq \phi(Q)$ .

For any LP-type problem  $(P, \phi)$  with **constant** combinatorial dimension,  $\phi(P)$  can be computed in randomized time **linear** in  $|P|$ .

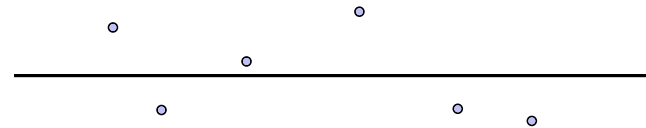
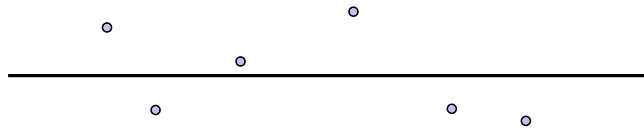
[Matoušek-Sharir-Welzl, 1992], [Seidel 1991], [Clarkson 1995]

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$Q$  is enclosed by the cylinder  
with axis  $\ell$  and radius  $r$



The line  $\ell$  is a transversal to the balls  
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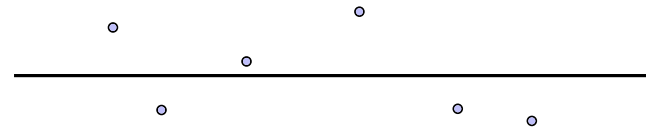
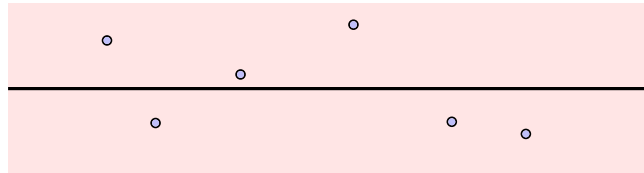


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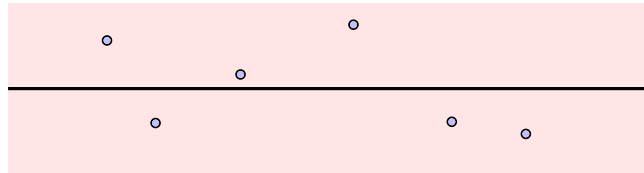


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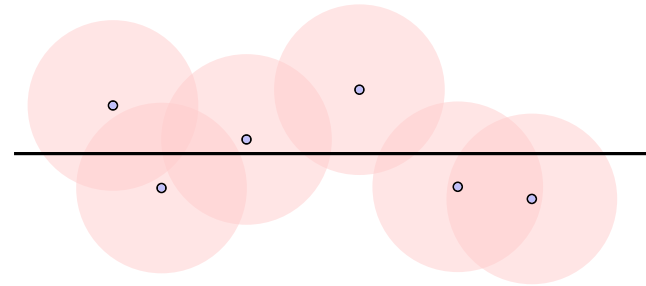


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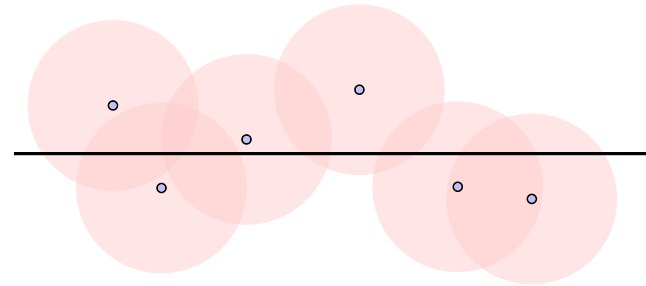
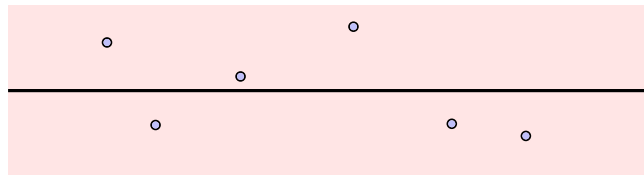


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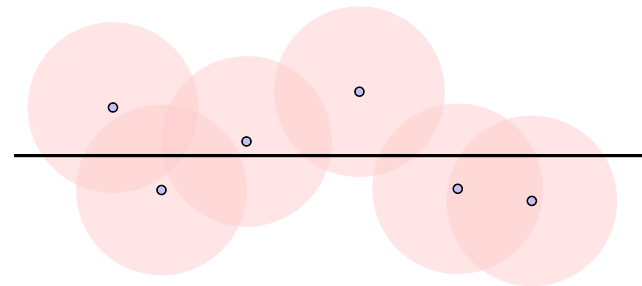
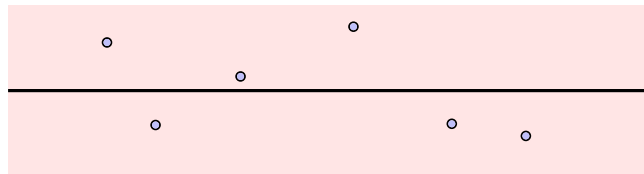
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**Corollary.** If  $P$  is sparse then  $(P, \phi)$  has combinatorial dimension at most  $4d - 1$  and the SEC of  $P$  can be computed in randomized linear time.

(in **any fixed** dimension  $d$ )

# Summary

Complete proof of Danzer's conjecture.

Algorithmic consequences.

The proof uses a combination of techniques from...

- ★ convex and euclidean geometry
- ★ topology
- ★ (classical) algebraic geometry
- ★ combinatorics

... and opens new perspectives

- ★ Topology of  $\mathcal{K}(\mathcal{F})$  for disjoint convex sets in  $\mathbb{R}^d$ .
- ★ Pinning theorems for disjoint convex sets in  $\mathbb{R}^3$ .



# Helly numbers from homological conditions

An interesting pattern...

Sets of line transversals with bounded Helly number...

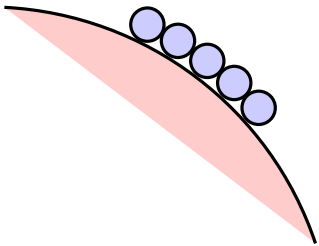
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Disjoint unit balls

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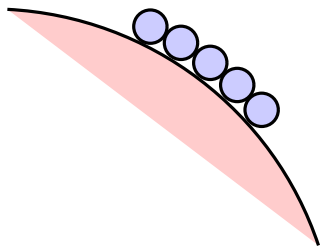
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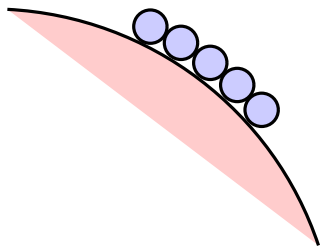
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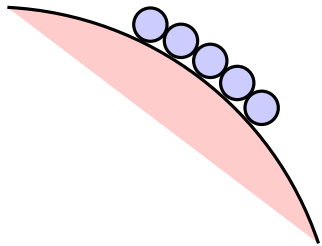
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The proofs all rely on ad hoc **geometric** arguments

Can we bring them under the same (topological) umbrella?

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**Theorem (Matoušek, 1999).** For any  $d \geq 2$  and  $r \geq 1$  **there exists** a constant  $h(d, r)$  such that the following holds: any finite family of subsets of  $\mathbb{R}^d$  such that the intersection of every subfamily has at most  $r$  connected components, each  $\lceil \frac{d}{2} \rceil$ -connected, has Helly number at most  $h(d, r)$ .

The topological condition resemble what we are looking for but...

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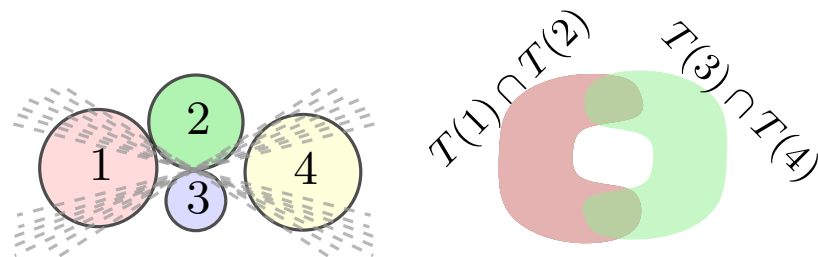
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**Theorem (Kalai-Meshulam, 2008).** Let  $G$  be an open **good cover** in  $\mathbb{R}^d$ . Any family  $\mathcal{F}$  such that the intersection of every subfamily is a disjoint union of at most  $r$  elements of  $G$  has Helly number at most  $r(d + 1)$ .

The bound look like what we'd like to have but...

- ★ Not the kind of topological conditions we have.





**Theorem (Colin de Verdière-Ginot-G, 2011).** If  $\mathcal{F}$  is a finite family of open subsets of  $\mathbb{R}^d$  such that the intersection of every subfamily has at most  $r$  connected components, each a homology cell, then  $\mathcal{F}$  has Helly number at most  $r(d + 1)$ .

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In fact, we prove a more general statement where:

★ the ambient space is any (locally connected) topological space  $\Gamma$ ,

$d$  is replaced by  $d_\Gamma$ , the minimum dimension from which **all** open sets of  $\Gamma$  have trivial homology.

★ only families of cardinality **at least**  $t$  need to intersect in at most  $r$  connected components,

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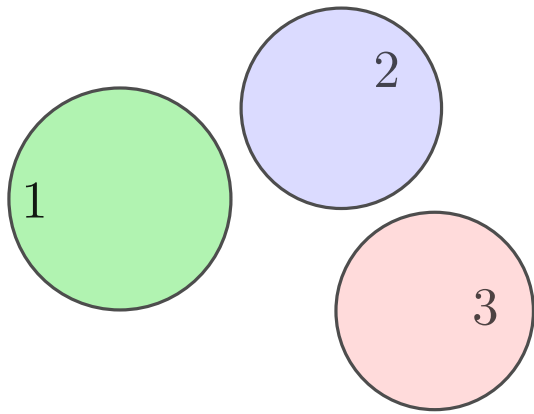
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This hammer implies Tverberg's theorem and Danzer's conjecture.

The **nerve**  $N(F)$  of a family  $\mathcal{F}$  of sets is:

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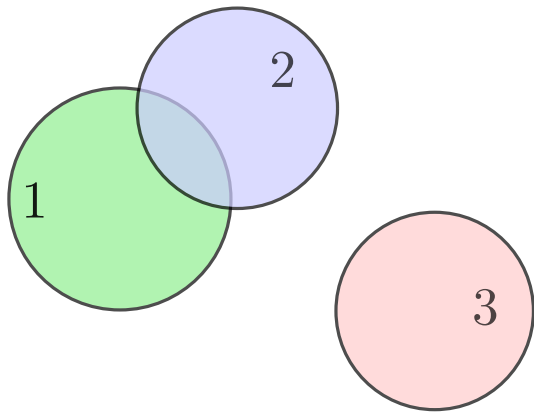
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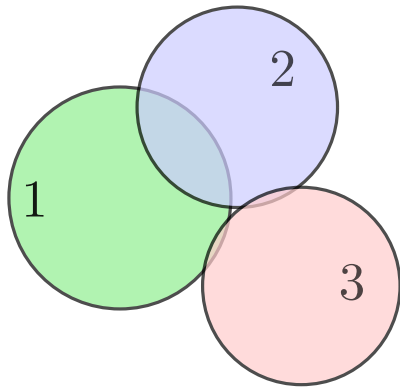
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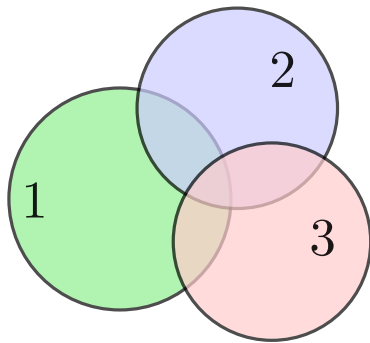
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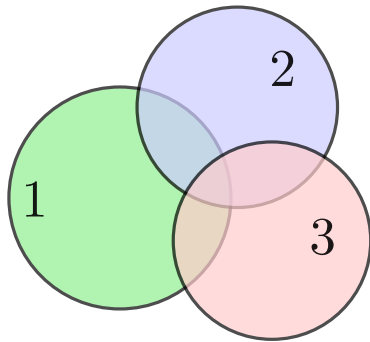
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$N(F)$

It is an abstract simplicial complex.

(= a family of finite sets closed under taking subsets).

(= a monotone hypergraph / set system).



Geometric realization of an abstract simplicial complex.

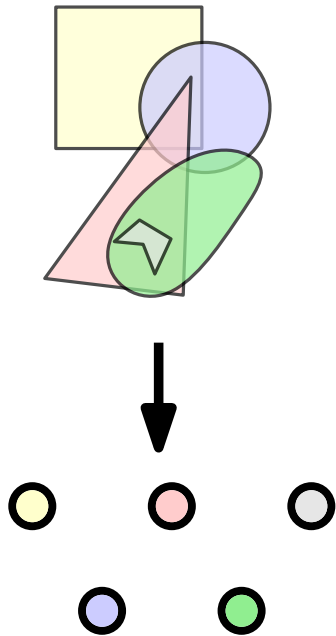
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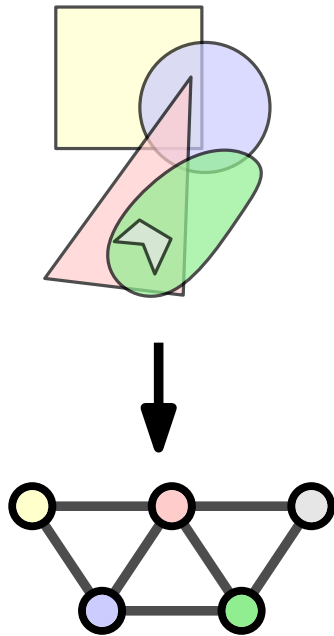
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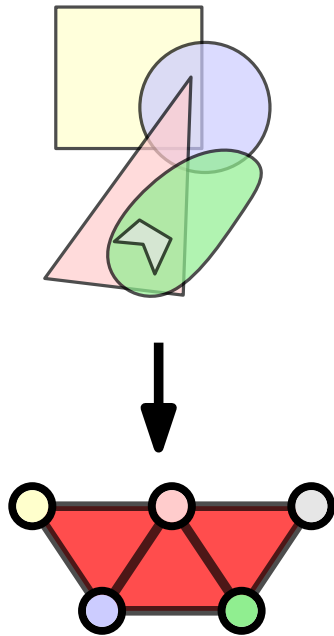
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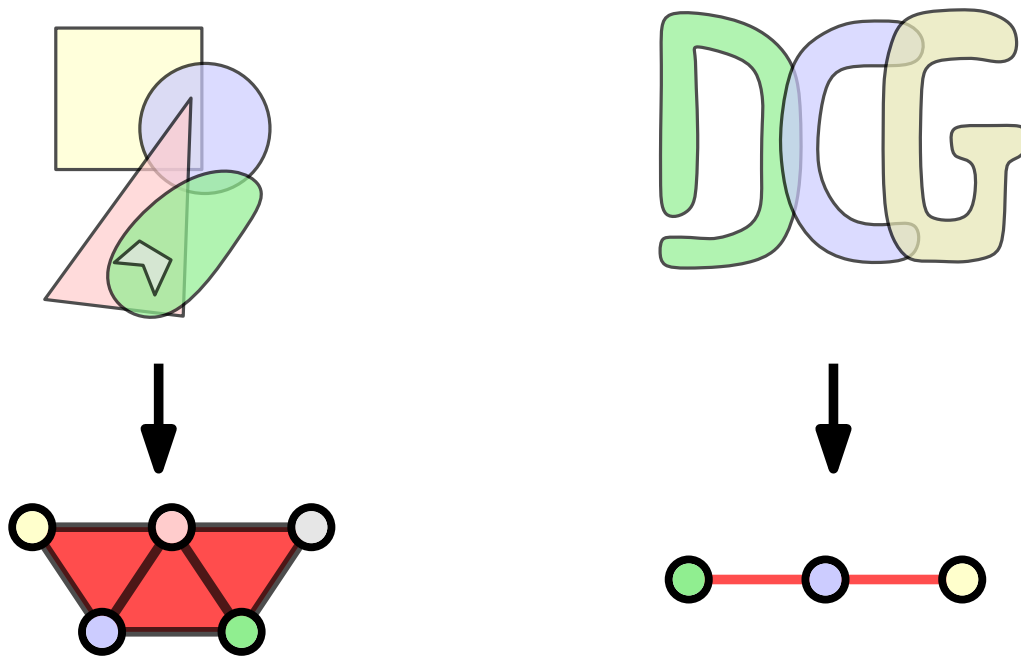
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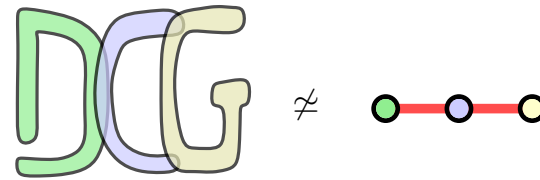
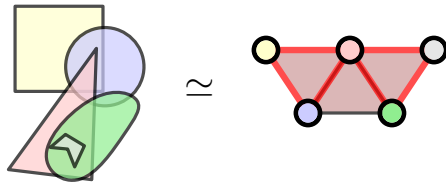
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Nerve Theorem  $\Rightarrow$  Helly's topological theorem

Let  $F$  be an **open** good cover in  $\mathbb{R}^d$ ,  
and  $G \subseteq F$  an inclusion-minimal subfamily with empty intersection.

$\bigcup_G$  must have **non-vanishing homology** in dimension  $|G| - 2$ .

Minimality  $\Rightarrow N(G) = 2^G \setminus \{G\} \simeq \mathbb{S}^{|G|-2}$

Nerve Theorem  $\Rightarrow \bigcup_G \simeq N(G)$ .

**Open** subsets of  $\mathbb{R}^d$  have vanishing homology in dimension  $\geq d$ .

$|G| - 2 < d \Rightarrow |G| \leq d + 1$ .



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We introduce a combinatorial structure that...

- ★ is close enough to a simplicial complex that Helly numbers are "within sight",
- ★ retains "enough" of the Čech complex that its homology is controlled by the union.

Let  $\mathcal{F}$  be a family of subsets of a [topological space](#).

The [multinerve](#)  $M(\mathcal{F})$  of  $\mathcal{F}$  is the [poset](#)

$$M(\mathcal{F}) = \{(G, X) \mid G \subseteq \mathcal{F}, X \text{ is a connected component of } \bigcap_G\}$$

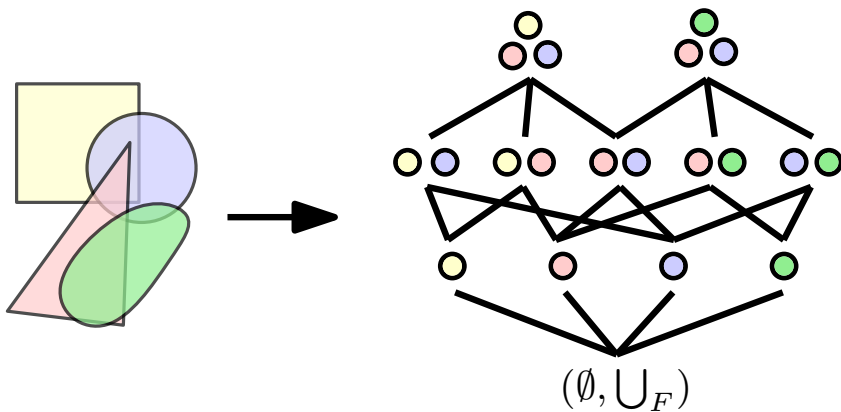
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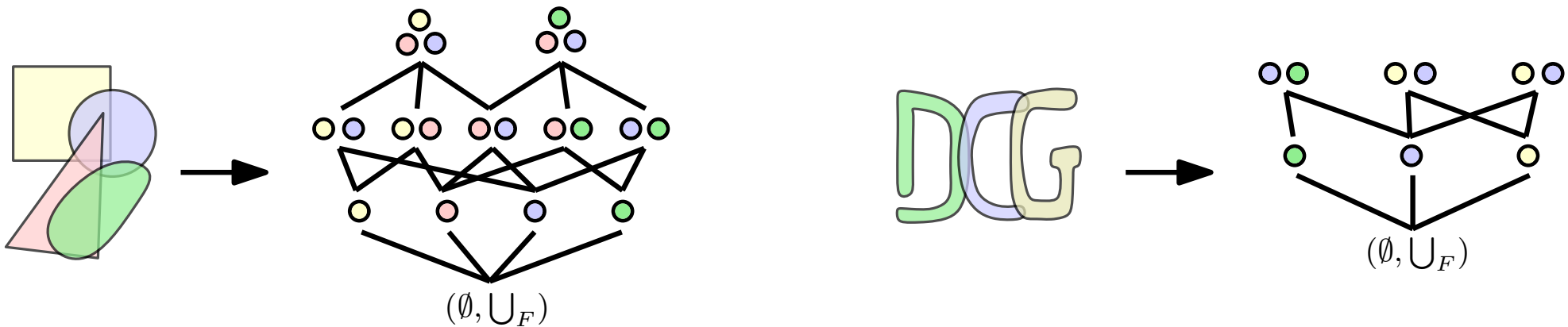


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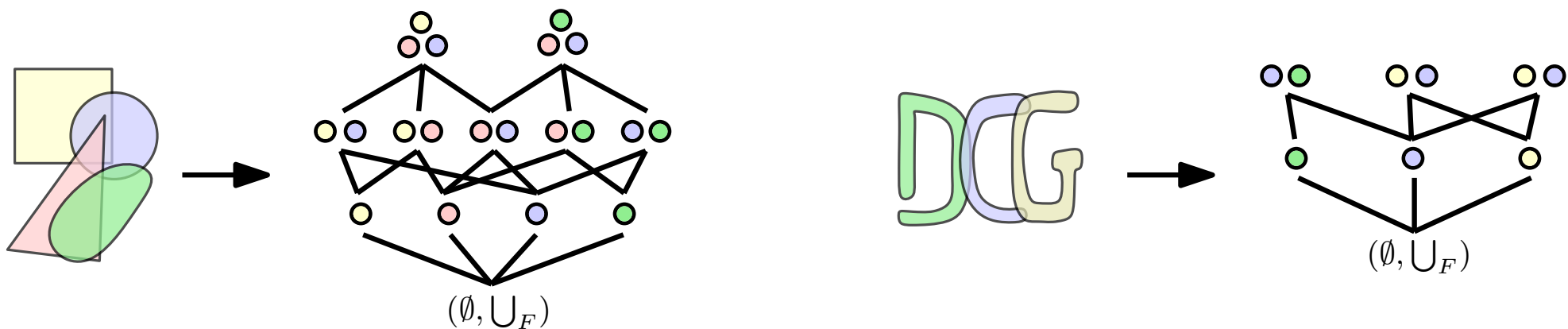


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$M(\mathcal{F})$  is a **simplicial** poset:  $\begin{cases} \text{Unique minimum element.} \\ \text{Every lower interval is isomorphic to the face lattice of a simplex.} \end{cases}$

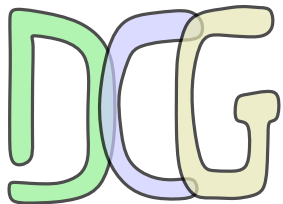
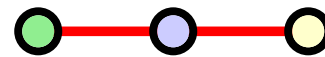
Geometric realizations, homology... extend to simplicial posets.



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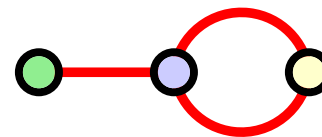
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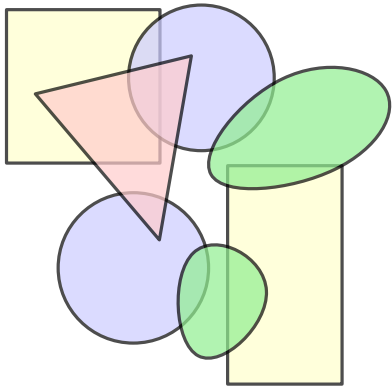


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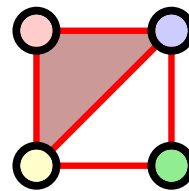
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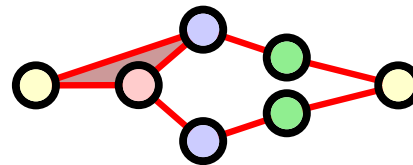


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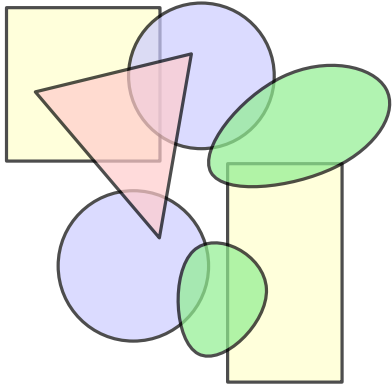


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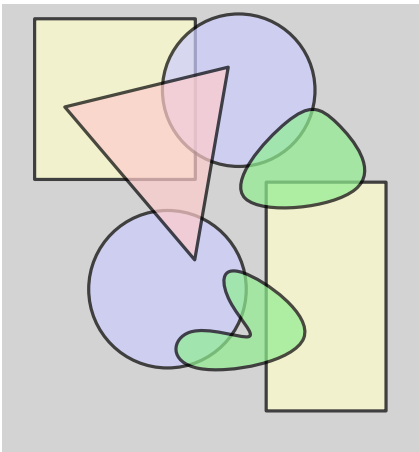
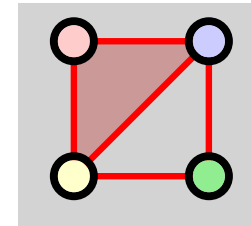
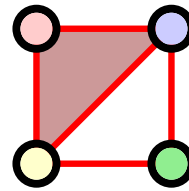
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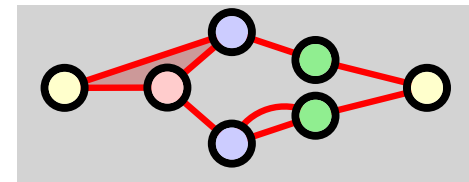
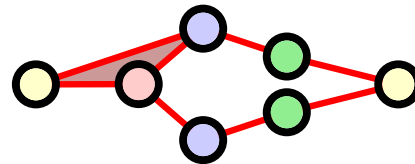


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Goal: "induced subcomplexes of  $N(\mathcal{F})$  have trivial homology in dimension  $\geq h$ "  
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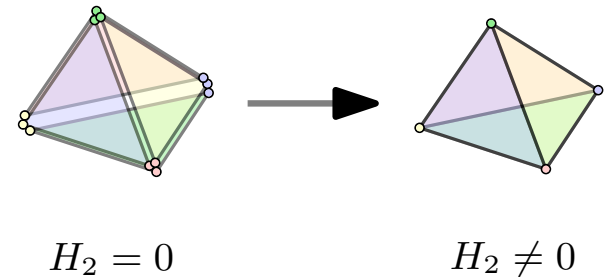
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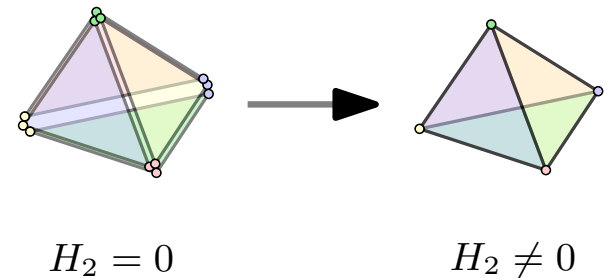
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★ the projection preserves dimension and is at most  $r$ -to-one  
 $\Rightarrow$  induced subcomplexes of  $N(\mathcal{F})$  have trivial homology in dimension  $\geq rd + r - 1$ .

Coming back to Danzer's conjecture...

If  $\mathcal{F} = \{B_1, \dots, B_n\}$  is a family of disjoint balls in  $\mathbb{R}^d$  and  $G \subseteq \mathcal{F}$  then...

$$T(B_i) \simeq \mathbb{RP}^{d-1},$$

$T(G)$  has contractible connected components if  $|G| \geq 2$ ,

The number of connected components of  $T(G)$  is at most 3 in general and at most 2 if  $|G| \geq 9$ .

The ambient space  $\Gamma$  is a compact subset of  $\mathbb{R}\mathbb{G}_{2,d+1}$  which is of dimension  $2d - 2$ .

Applying our homological Helly-type theorem we obtain:

For  $d \geq 6$  the Helly number of  $\{T(B_1), \dots, T(B_n)\}$  is at most  $2(2d - 1)$ .

(to use  $r = 2$  in our bound we need  $2d - 2 \geq 9$ , hence the condition  $d \geq 6$ ).

# Summary

Refinement of the classical nerve that enjoys a similar "Nerve Theorem".

"Combinatorial interface" to Leray's acyclic cover theorem.

Homological Helly-type theorem that

essentially generalizes those of Matoušek and Kalai-Meshulam,  
(re)proves in a unified manner Helly numbers in geometric transversal theory.

Raises questions on the combinatorics of simplicial complexes and posets.

dimension-preserving projections

# Some perspectives

## Short/medium term

### Simplify...

- ★ projection of simplicial complexes

### Generalize...

- ★ Topology of sets of  $k$ -dimensional transversals to convex sets in  $\mathbb{R}^d$ .
- ★ Tangents to convex sets and transversality

### Apply...

- ★ Efficient computation of SEC, from sparse to general point sets.
- ★ Other uses of multinerves for Hadwiger-type theorems, Stanley-Riesner ideals...

## In a more distant future

### More territory to map...

- ★ Does a "recursively" bounded sum of Betti numbers imply a bounded Helly number?
- ★ Algorithmic applications of bounded "local combinatorial dimension"?

### And some "hard nuts" to break...

- ★ Geometric permutations of disjoint convex sets in  $\mathbb{R}^d$ .
- ★ Pinning theorem for disjoint convex sets in  $\mathbb{R}^3$ ?  $\mathbb{R}^d$ ?

# Other directions

## ★ Combinatorics of geometric structures

Shatter functions, VC-dimension & excluded patterns for geometric permutations.

Topological combinatorics (inclusion-exclusion formulas...)

Random generation of combinatorial structures underlying geometric objects (order-types...).

## ★ Complexity of random geometric structures

Average-case analysis (random polytopes, Delaunay of points on a surface).

Smoothed complexity (convex hull, Delaunay triangulation...).

## Co-workers (2004-2011)

### Geometric transversals

Boris Aronov (NYU-Poly)  
Ciprian Borcea (Rider University)  
Otfried Cheong (KAIST)  
Andreas Holmsen (KAIST)  
Stefan Koenig (TU Munchen)  
Sylvain Petitjean (INRIA)  
Günter Rote (FU Berlin)

### Random geometric structures

Dominique Attali (CNRS)  
Olivier Devillers (INRIA)  
Jeff Erickson (UIUC)  
Marc Glisse (INRIA)

⋮

### Line geometry for visibility & imaging

Guillaume Batog (U. Nancy 2)  
Julien Demouth (U. Nancy 2)  
Jeong-Hwan Jang (KAIST)  
Bruno Levy (INRIA)  
Jean Ponce (ENS)

### Combinatorics of geometric structures

Otfried Cheong (KAIST)  
Éric Colin de Verdière (CNRS)  
Grégory Ginot (U. Paris 6)  
Cyril Nicaud (U. Marne la Vallée)

⋮

**And also...** Mark de Berg (TU Eindhoven), Veronique Cortier (CNRS), Hyo-Sil Kim (KAIST), Jan Kratochvil (Charles U.), Sylvain Lazard (INRIA), Mira Lee (KAIST), Hyeon-Suk Na (Soongsil U.), Yoshio Okamoto (JAIST), Chan-Su Shin (HUFS), Frank Van der Stappen (Utrecht U.), Alexander Wolff (U. Würzburg).

(student at the time of the collaboration)