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Appendix A: Additive Homomorphic Encryption Scheme from [164]

Description

The main idea of the scheme is to replace the xor (Exclusive-OR) operation typically found in stream ciphers with modular addition (+).

Assume that $0 \leq m < M$. Due to the commutative property of addition, the above scheme is additively homomorphic. In fact, if $c_1 = Enc(m_1, k_1, M)$ and $c_2 = Enc(m_2, k_2, M)$ then $c_1 + c_2 = Enc(m_1 + m_2, k_1 + k_2, M)$.

Note that if n different ciphers c_i are added, then M must be larger than $\sum_{i=1}^n m_i$, otherwise correctness is not provided. In fact if $\sum_{i=1}^n m_i$ is larger than M , decryption will result in a value m' that is smaller than M . In practice, if $p = \max(m_i)$ then M should be selected as $M = 2^{\log_2(p*n)}$.

The keystream k can be generated by using a stream cipher, such as RC4, generated from a private key.

Security Analysis

This additive homomorphic encryption scheme is very similar to a xor-based stream cipher and its security can be proven using a similar proof.

The security relies on two important features: (1) the keystream changes from one message to another and (2) all the operations are performed modulo an integer M . These two features protect the scheme from frequency analysis attacks. In fact, it can be proven that the scheme is *perfectly secure*.

Theorem 1 *The previous encryption scheme is perfectly secure.*

Preuve For plaintext space M , keystream space K , let $\mathcal{K} = |M|$, $m \in [0; M - 1]$, $c \in$

Additively Homomorphic Encryption Scheme

Encryption:

1. Represent message m as integer $m \in [0, M - 1]$ where M is a large integer
2. Let k be a randomly keystream, where $k \in [0, M - 1]$
3. Compute $c = Enc(m, k, M) = m + k \pmod{M}$.

Decryption:

1. $Dec(c, k, M) = c - k \pmod{M}$

Addition of Ciphertexts:

1. Let $c_1 = Enc(m_1, k_1, M)$ and $c_2 = Enc(m_2, k_2, M)$
2. For $k = k_1 + k_2$, $Dec(c_1 + c_2, k, M) = m_1 + m_2$

$[0; M - 1]$. Set $k^* = c - m \pmod{M}$. Then:

$$\begin{aligned}
\text{Prob}_{k \leftarrow \mathcal{K}}[Enc(k, m, M) = c] &= \text{Prob}_{k \leftarrow \mathcal{K}}[k + m = c \pmod{M}] \\
&= \text{Prob}_{k \leftarrow \mathcal{K}}[k = c - m \pmod{M}] \\
&= \text{Prob}_{k \leftarrow \mathcal{K}}[k = k^*]
\end{aligned} \tag{1}$$

If we assume that the maximum number of ciphertexts to be added is n and that each plaintext is l -bit long, we must have $M = 2^{l+\log(n)}$, i.e., $|M| = l + \log(n)$. If $c_i = (m_i + k_i)$, then the probability that $c_i \in [0, 2^l - 1]$ is twice the probability that $c_i \in [2^l; M - 1]$. More specifically, we have: $\text{Prob}_{k \leftarrow \mathcal{K}}[k = k^*] = 1/(2^l + M)$ if $c > 2^l$ and $\text{Prob}_{k \leftarrow \mathcal{K}}[k = k^*] = 2/(2^l + M)$ if $c < 2^l$.

Since these two equations hold for every $m \in \mathcal{M}$, it follows that for every $m_1, m_2 \in \mathcal{M}$ we have $\text{Prob}_{k \leftarrow \mathcal{K}}[Enc(k, m_1, M) = c] = \text{Prob}_{k \leftarrow \mathcal{K}}[Enc(k, m_2, M) = c]$ which establishes perfect security of the scheme.

