

# Scheduling Handling Resources: Robotic Flowshops with Circular Layouts and a Practical Railway Problem

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# Scheduling and transportation

Scheduling: assigning *tasks* to *resources* submitted to *constraints* to optimize a *performance criteria*

Physical handling or transportation between resources: connectivity, layout, additional specific handling resources

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Two contexts:

- **Manufacturing management**

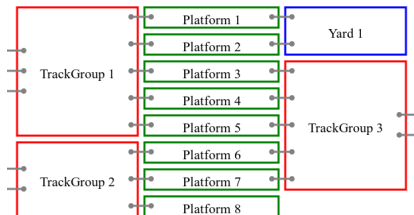


- **Railway management**



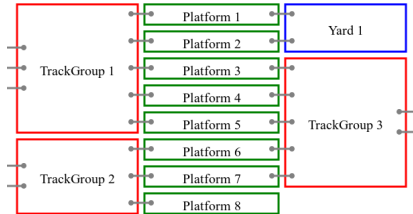
# Railway management

“Trains don't vanish”...except when they do.  
EURO/Roadef Challenge 2014, SNCF



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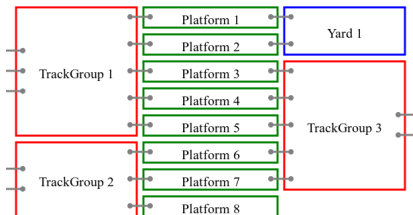


- Resource capacity
- Resource compatibility
- Conflict between trains
- ...



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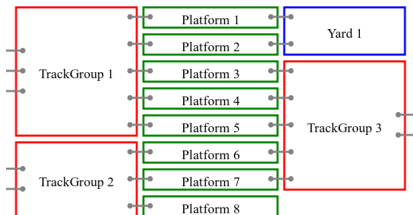
- Resource capacity
- Resource compatibility
- Conflict between trains
- ...

Optimize departure coverage and performance costs



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Horizon: 14 days

200–2000 departures/arrivals

30–90 resources

- Resource capacity
- Resource compatibility
- Conflict between trains
- ...

Optimize departure coverage and performance costs



# Railway management

Joint work with H. Joudrier [[Joudrier and T. 17](#)]

- Simplify the station's graph by grouping similar resources
- Sophisticated routing algorithm
- (very) simple assignment algorithm
- Up to 40% covered departures...

... And the 1st place in the junior category!





# Manufacturing management

Flowshop: the parts must be processed on machine  $M_1$ , then  $M_2$ ...



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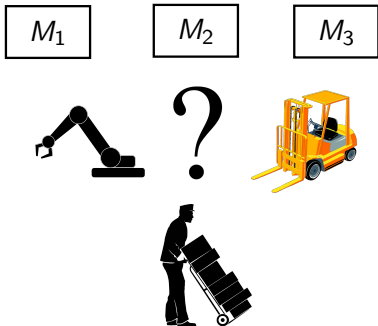


Sufficient model?



# Manufacturing management

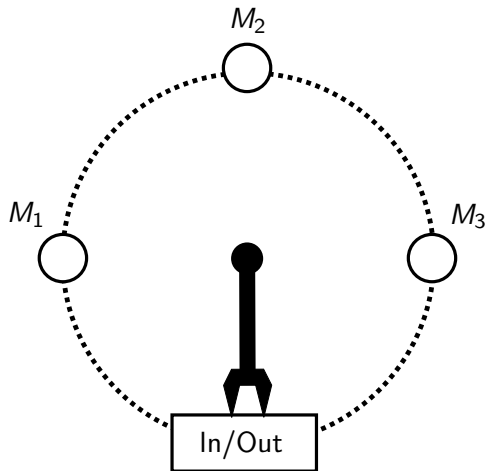
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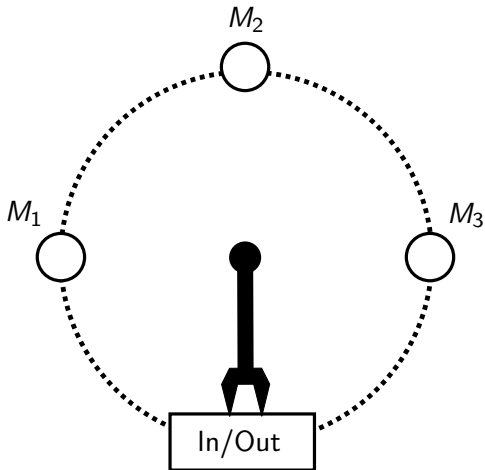
Sufficient model? Operator, tools, **handling resources**...



# Robotic cell – Example

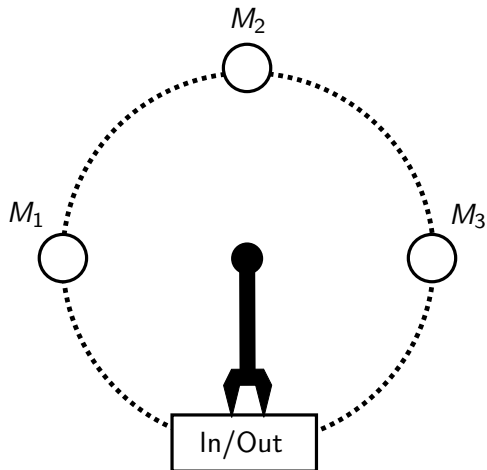


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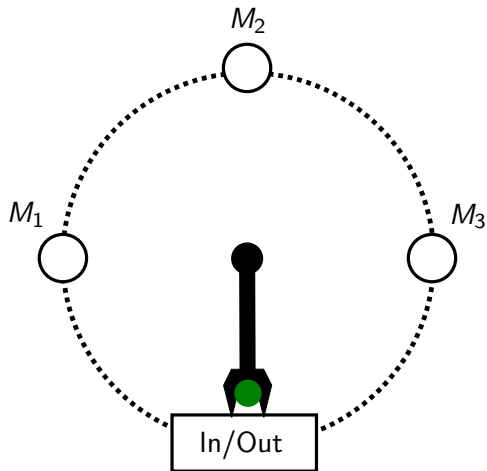
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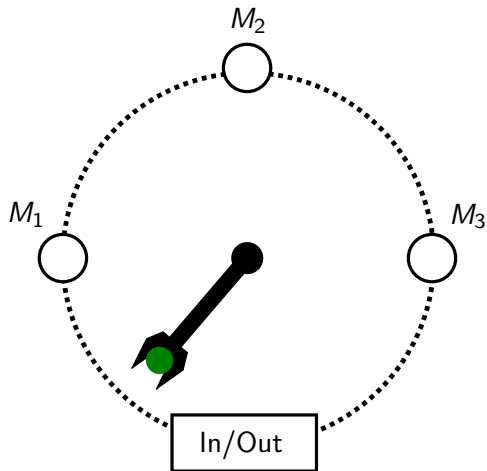
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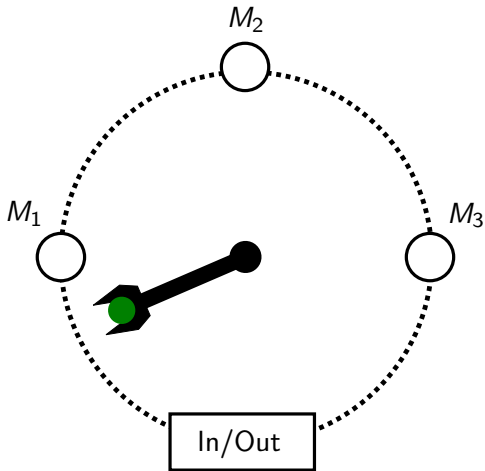
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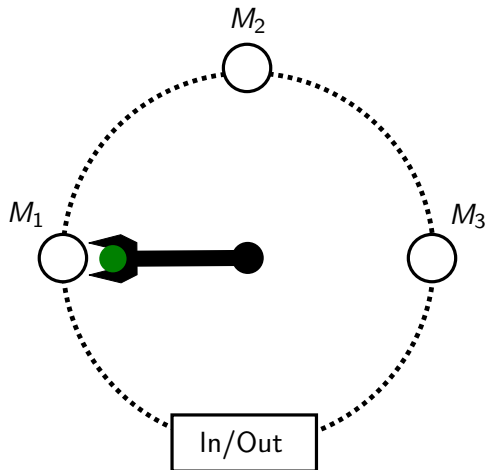


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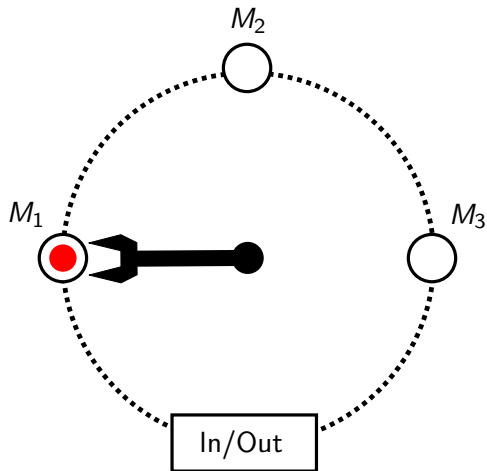
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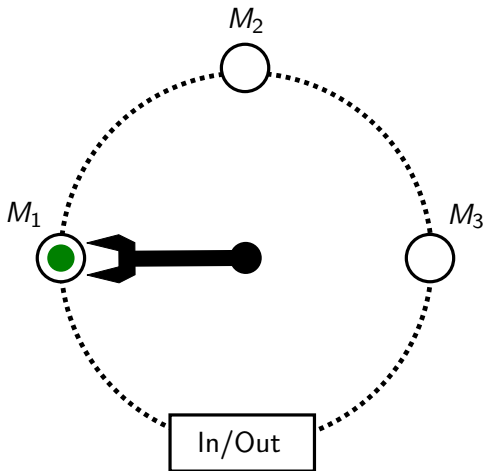
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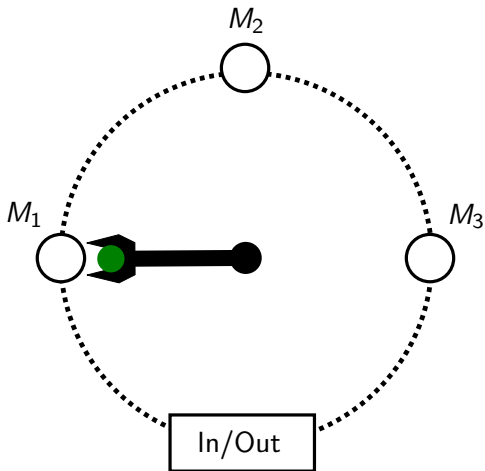
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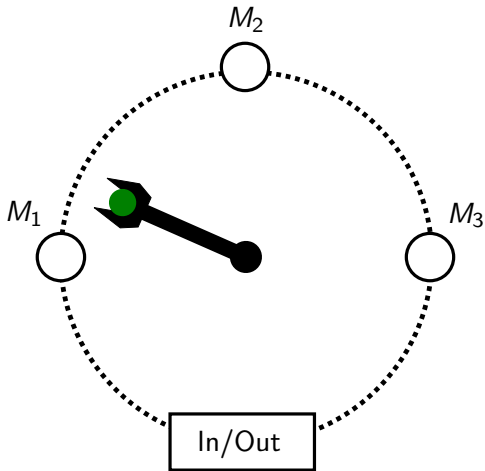
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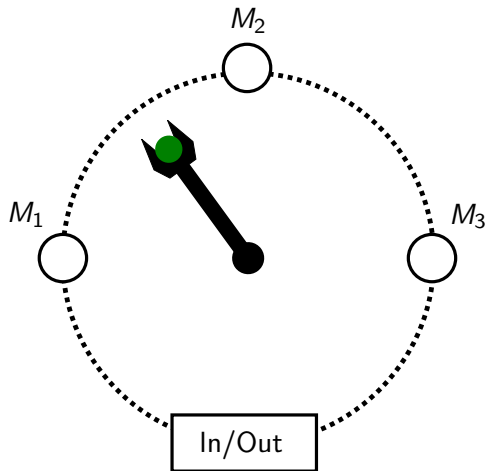
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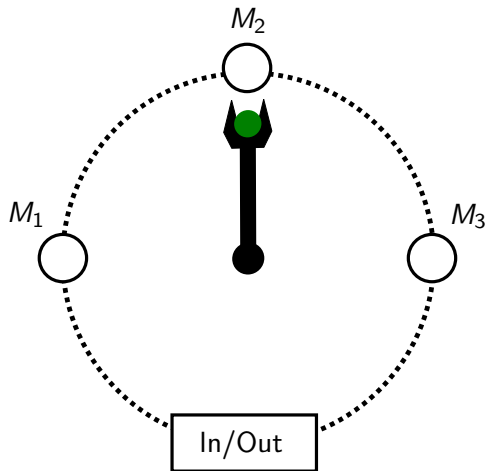
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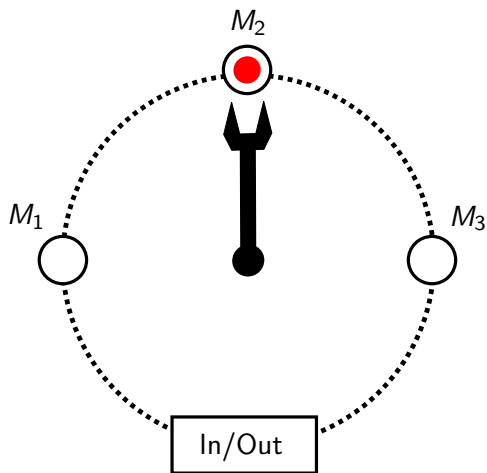
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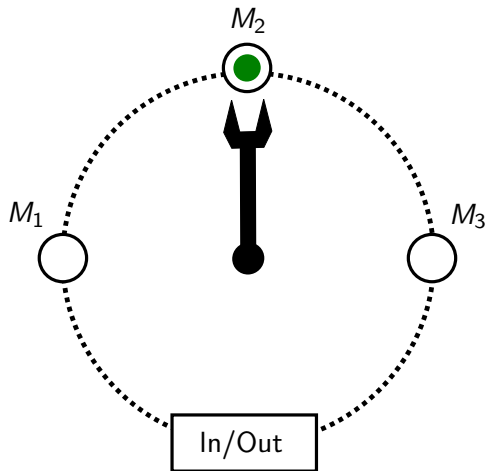


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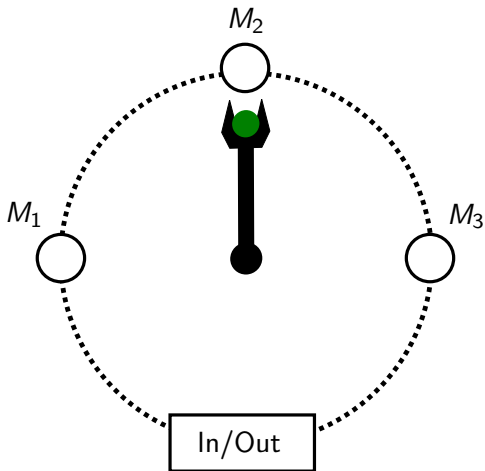
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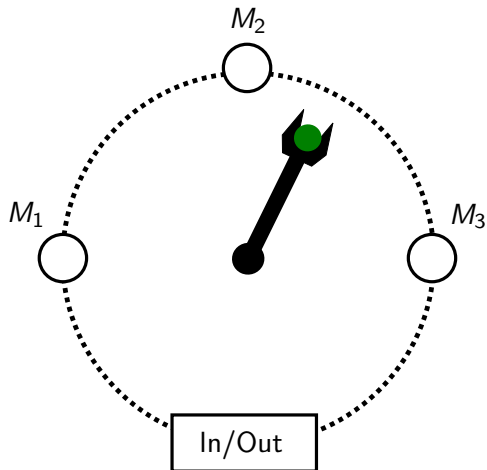
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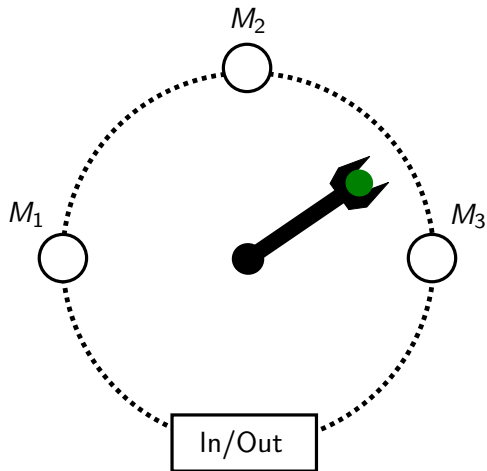
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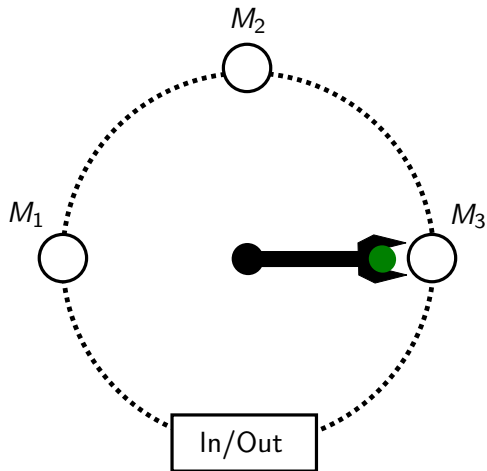
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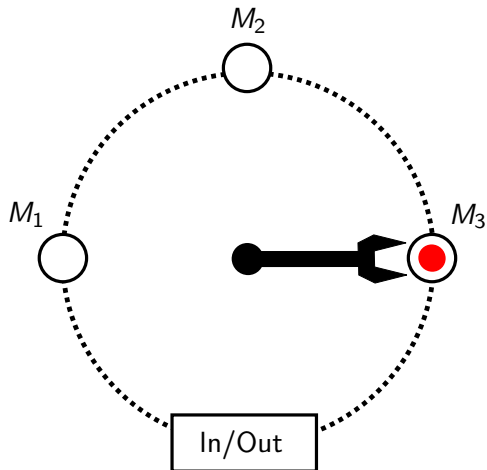
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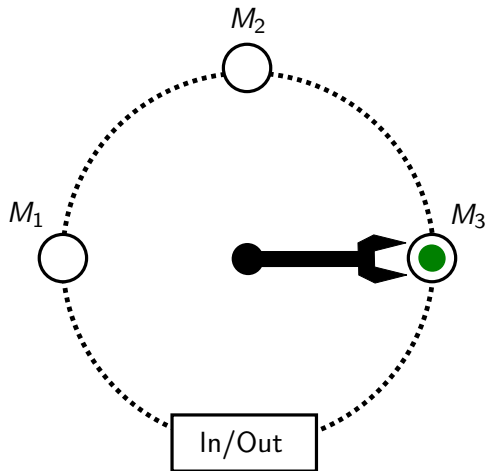
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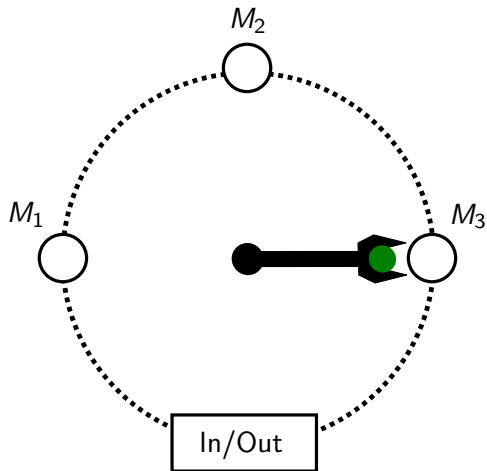
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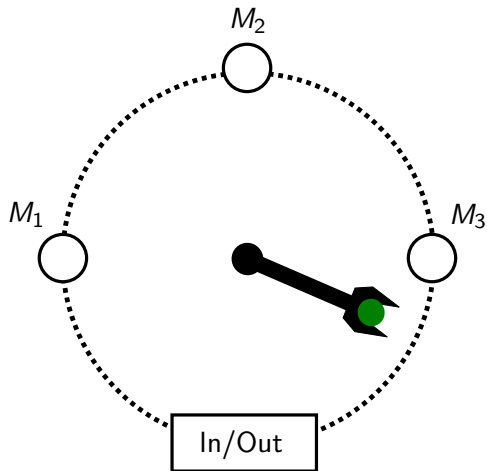


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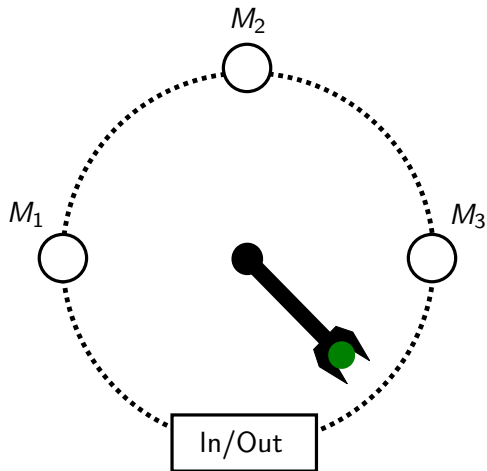
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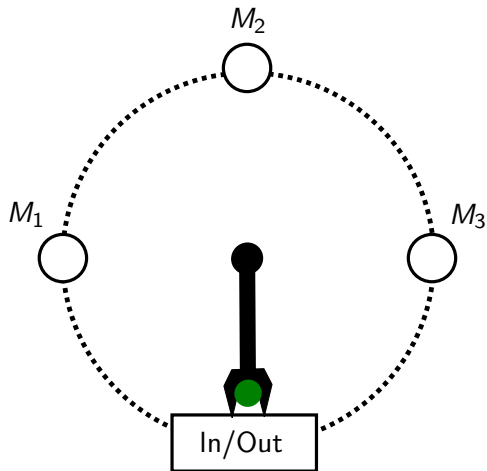
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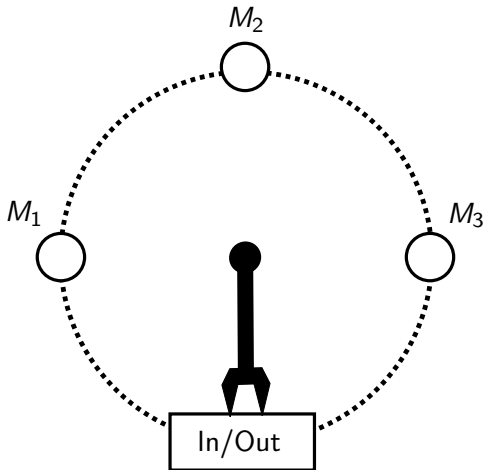
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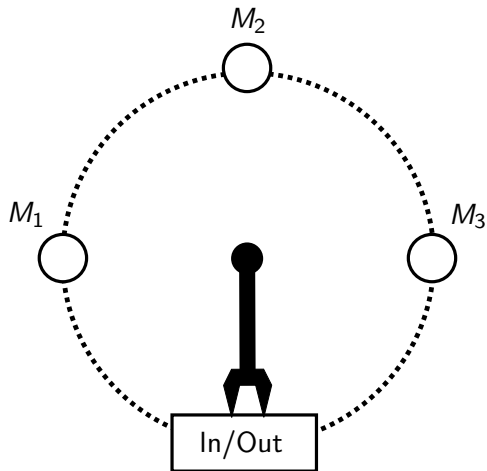
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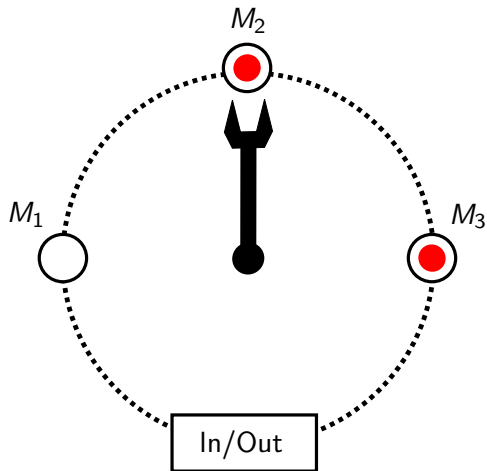
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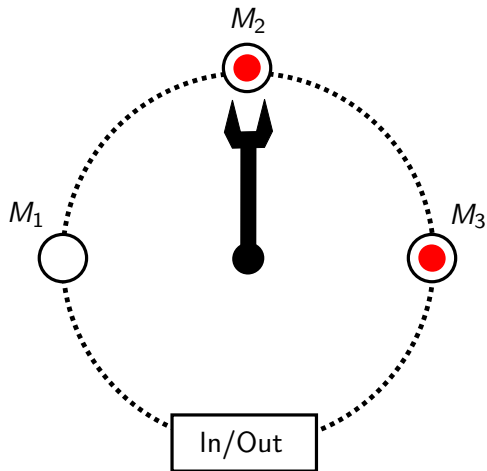
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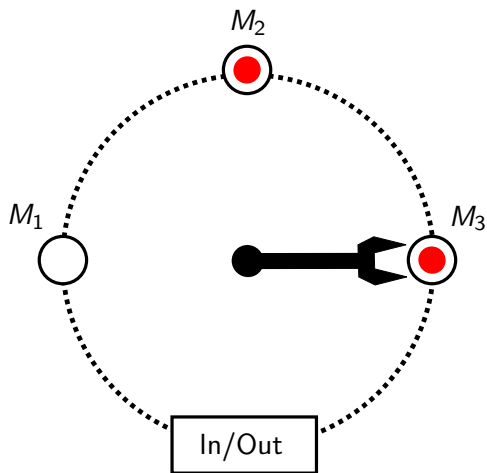
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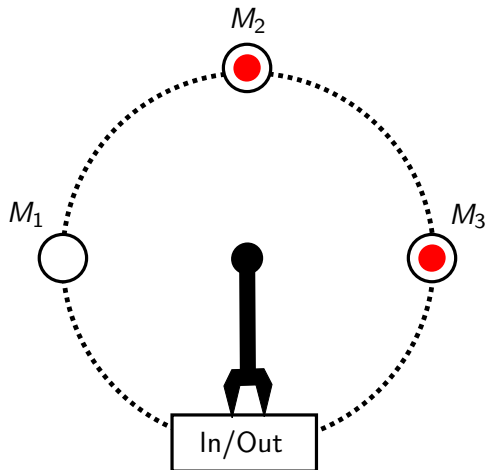


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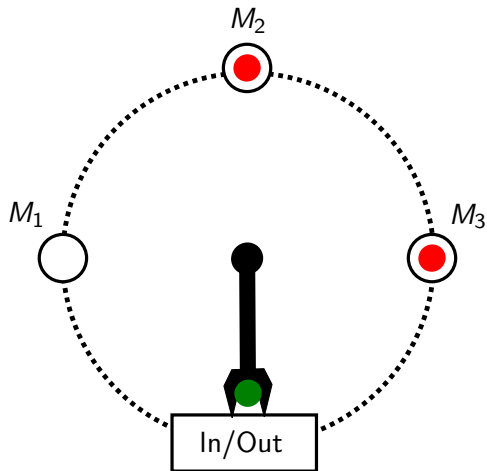
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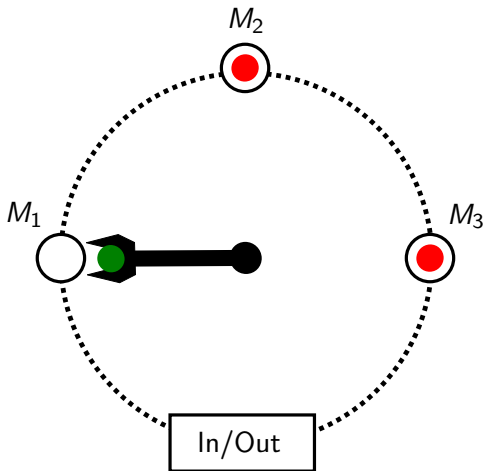
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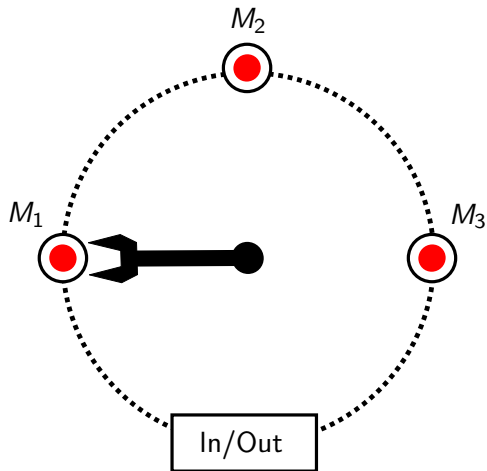
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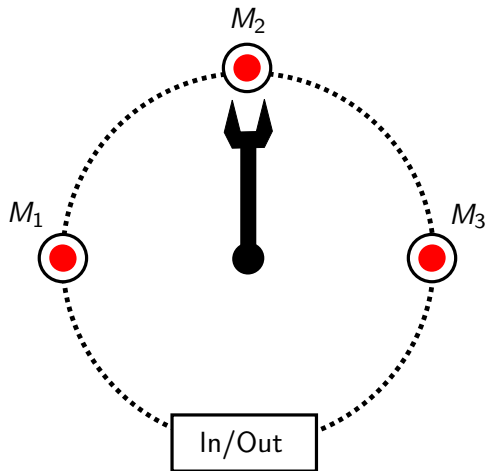
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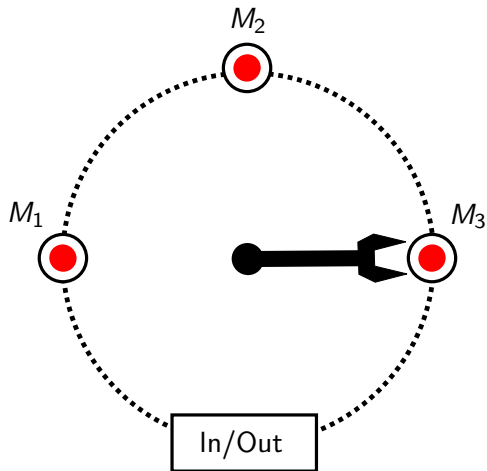
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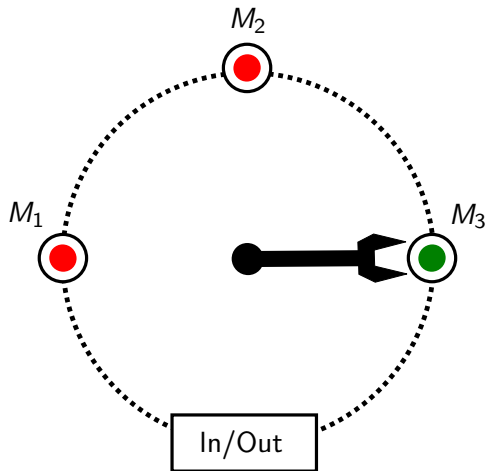
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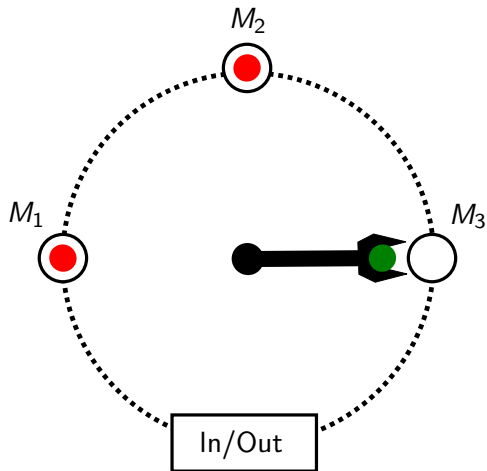
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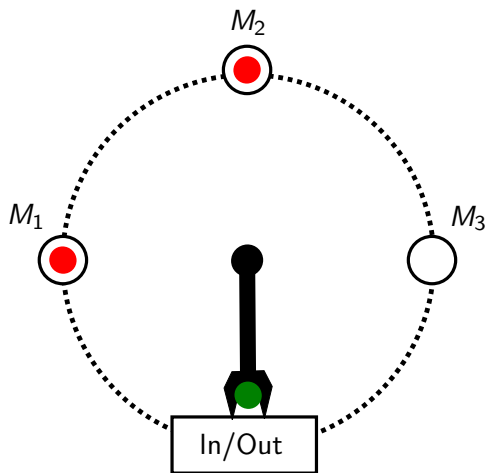


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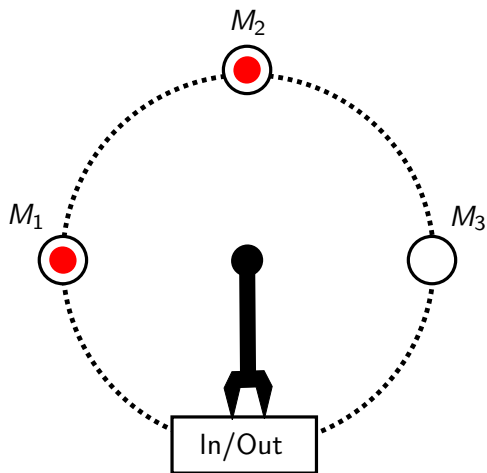
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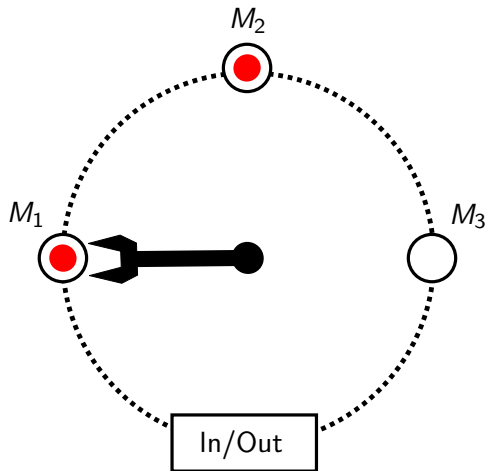
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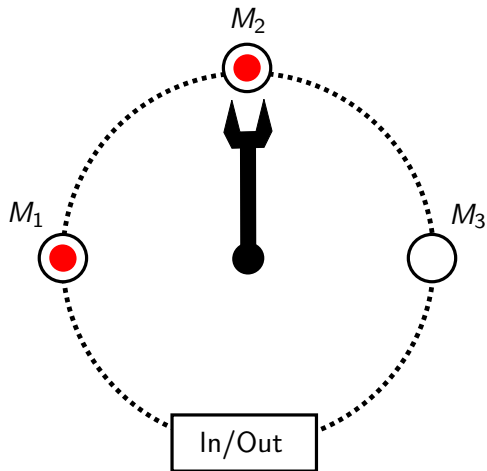
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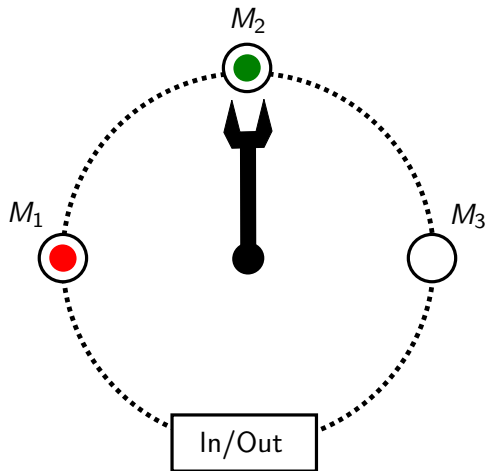
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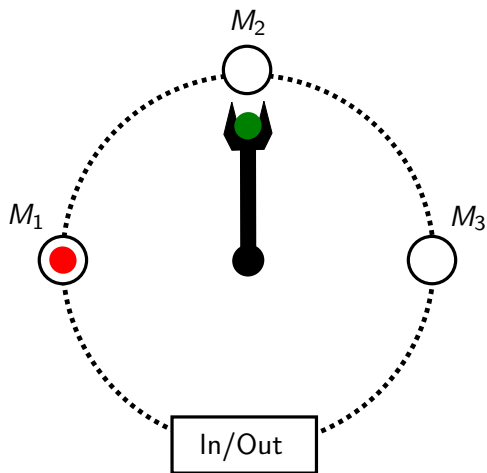
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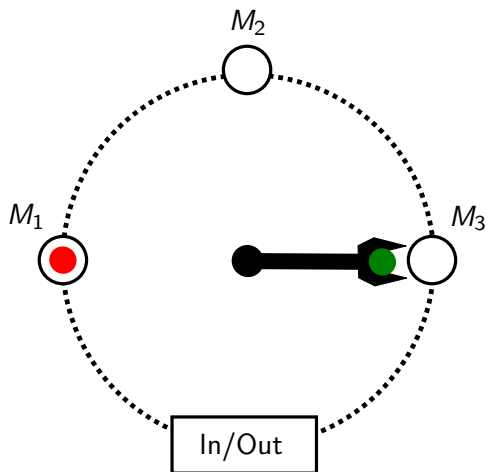
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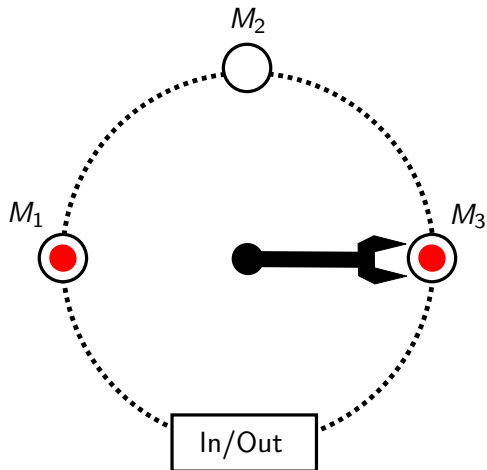
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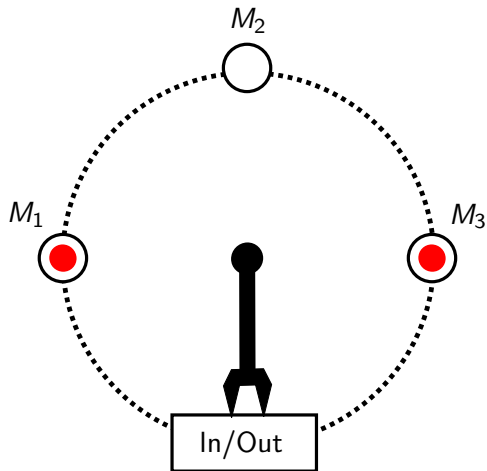


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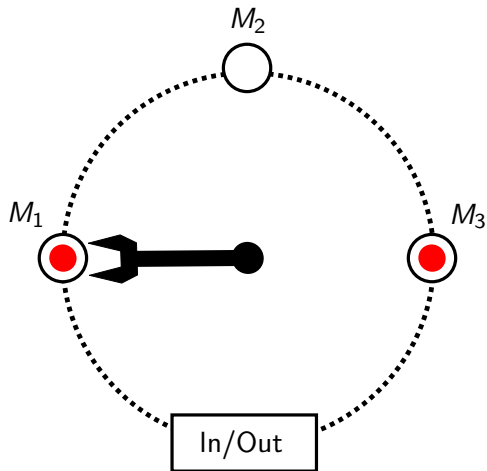
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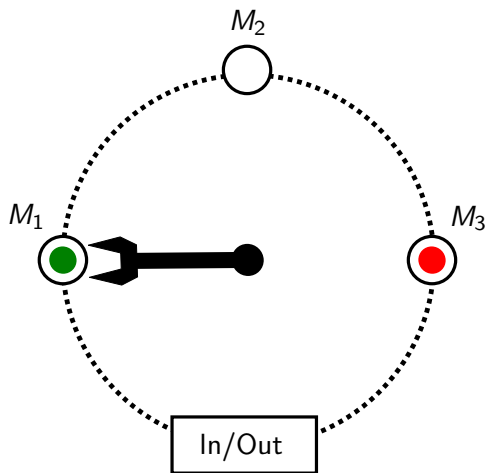
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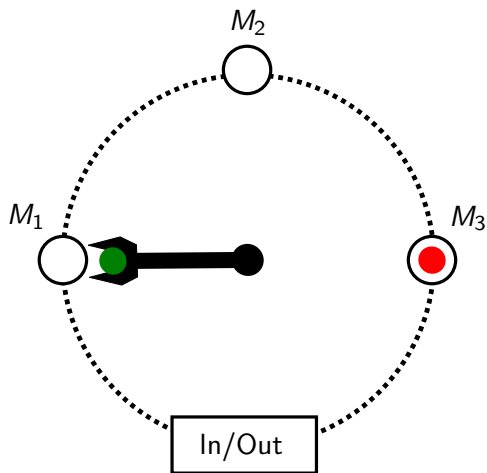
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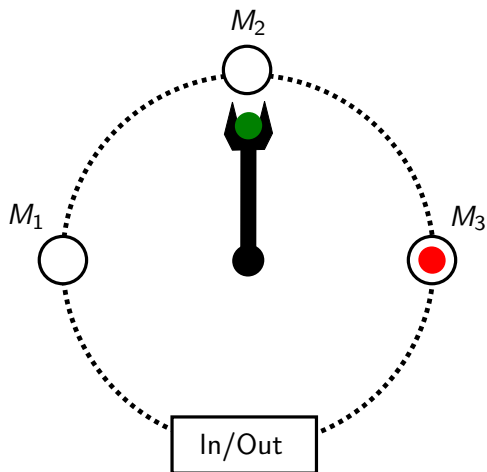
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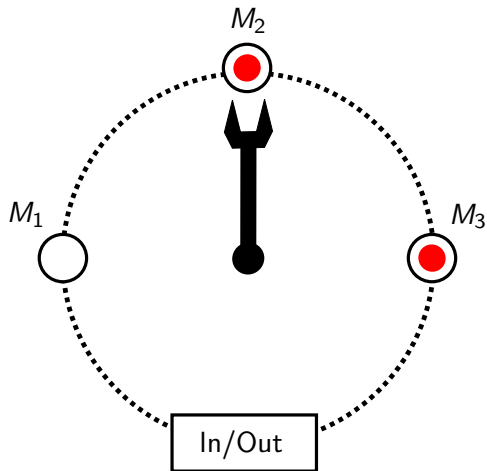
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# And in-between?

That's where the problem is

Balancing

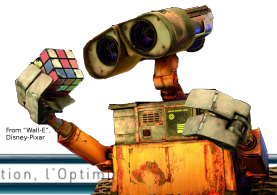
- Total travel time
- Waiting time



Given the *processing times* and *travel time between machines*, find a robot programming which optimizes the **throughput**:

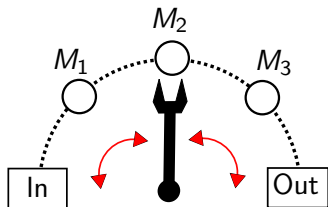
$$\text{Maximize } \frac{k}{\text{Time to produce } k \text{ parts}}$$

$$\text{Minimize } \frac{\text{Time to produce } k \text{ parts}}{k}$$

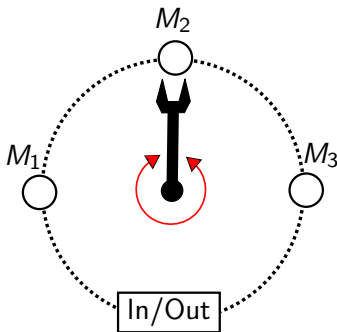


# Models: Layouts

Linear



Circular



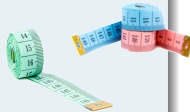
# Models: cell parameters

## Waiting policy

- No-wait [Agnetis, 00 ; Kats *et al* ; 09, Che *et al*, 12]
- **Unbounded** [Crama *et al*, 97 ; Rajapakshe *et al.*, 11]
- Time-window (HSP) [Dawande *et al*, 09; Zhou *et al*, 12]

## Travel metric

- General
- **Additive**  
| regular
- Constant



## Processing times

### Identical parts

- General
- **Balanced**



# Machine and robot capacity

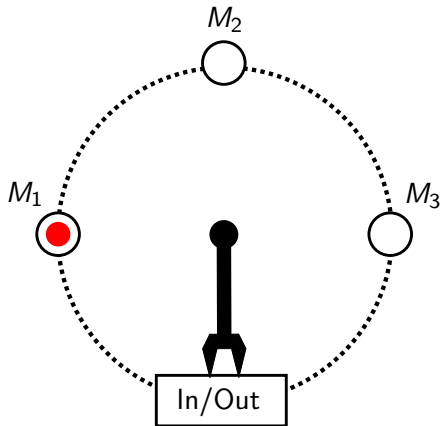
## Classical model (blocking)

Few studies for the circular layout [Rajapakshe *et al.*, 11]

## Relaxing the blocking constraint

- Dual-gripper  
[Sethi *et al.*, 01 ; Jung *et al.*, 15 ; Drobouchevitch *et al.*, 06]
- Swapping [Jolai *et al.*]
- Machine buffers [Drobouchevitch *et al.*, 06]

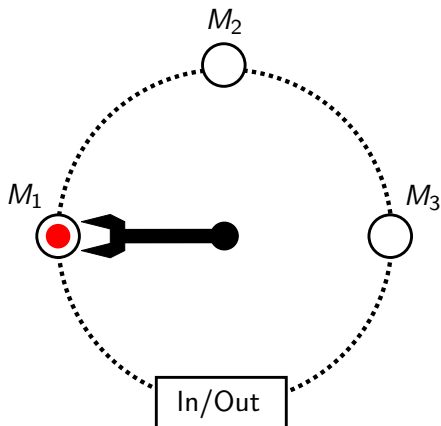
# SCOP Describing the robot moves: Activities



[Crama *et al*, 97]

Activity  $A_i$

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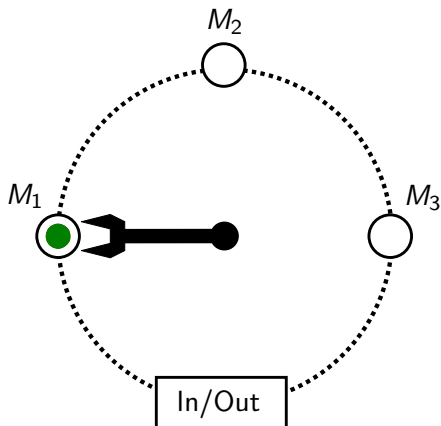


[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$

# SCOP Describing the robot moves: Activities

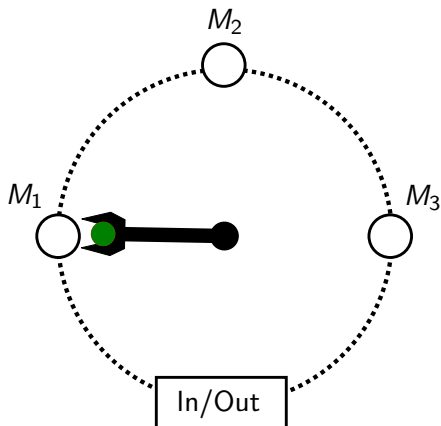


[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$
- Wait?

# SCOP Describing the robot moves: Activities



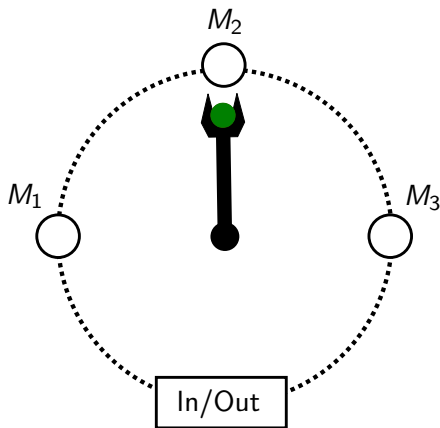
[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$
- Wait?
- Unload  $M_i$



# SCOP Describing the robot moves: Activities

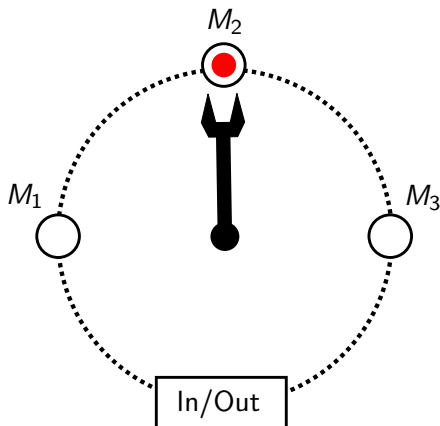


[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$
- Wait?
- Unload  $M_i$
- Go to  $M_{i+1}$

# SCOP Describing the robot moves: Activities

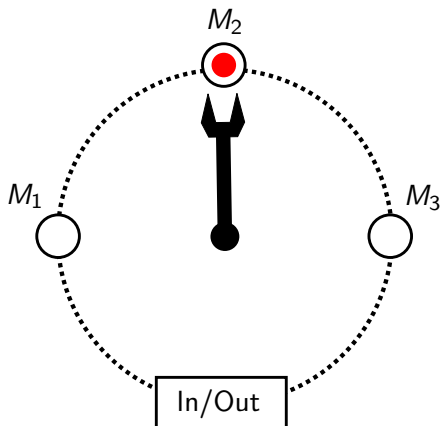


[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$
- Wait?
- Unload  $M_i$
- Go to  $M_{i+1}$
- Load  $M_{i+1}$

# SCOP Describing the robot moves: Activities



[Crama *et al*, 97]

## Activity $A_i$

- Go to  $M_i$
- Wait?
- Unload  $M_i$
- Go to  $M_{i+1}$
- Load  $M_{i+1}$

$A_0, A_1, \dots, A_m$

# Cyclic programmation

Dominance [Dawande *et al*, 05]

## Cycle

- Feasible sequence of activities
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1-cycles  $\Leftrightarrow$  permutations of the  $m$  activities  $A_1 \dots A_m$

## Objective function

max throughput  $\longleftrightarrow$  min  $\frac{\text{cycle time}}{\text{number of parts produced}}$   $\longleftrightarrow$  min  $\frac{T(C_k)}{k}$

# Best 1-cycle problem

## 1-cycles are...

- easy to describe
- easy to implement
- better known and more studied

But not necessarily optimal!



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## Finding the best 1-cycle...

- ...in linear regular cells: P [Crama *et al.*, 97]  
(balanced [Brauner *et al.*, 99])
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( $\frac{5}{3}$  approx).

- ...in circular regular balanced cells: ??

# Dominance?

1-cycle conjecture

1-cycles dominate all cycles

[Sethi *et al.*, 92]

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[Crama *et al.*, 97]

[Brauner *et al.*, 01]

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## Circular layout

???

# Dominance?

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## Linear layout

- Valid for regular cells, 2- to 3-machine [Crama *et al.*, 97]
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## Circular layout

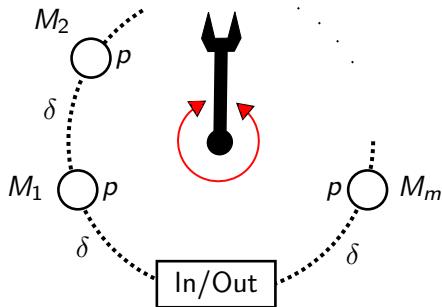
**Counter-example for regular balanced cells with 6 machines**

[T. *et al.* MIM 16]

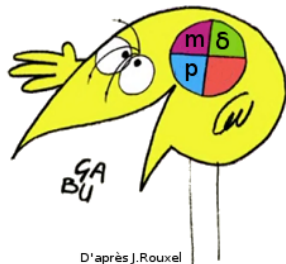


# Our problem

Circular, identical parts, **regular, balanced**:  $(m, \delta, p)$



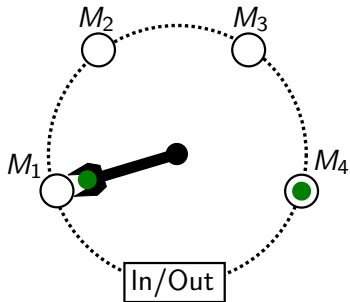
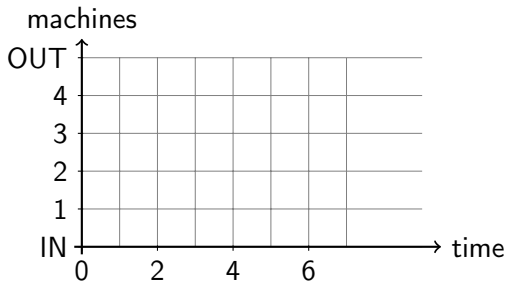
Small enough to fit in a Shadok's head!



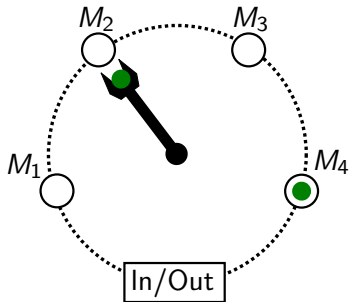
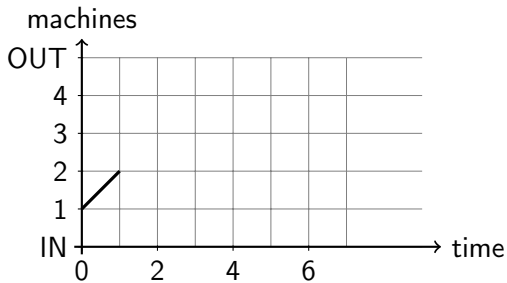
D'après J. Rouxel



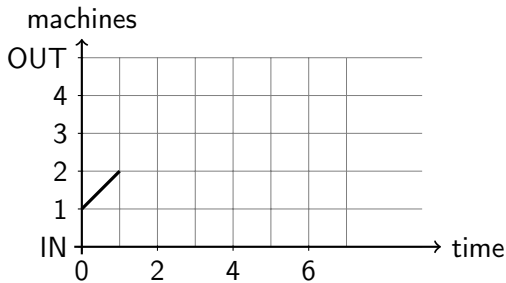
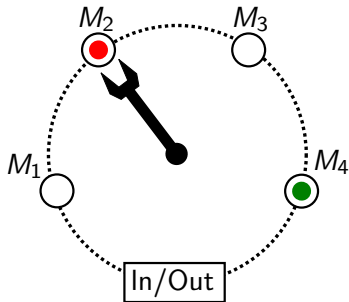
# Representing the robot moves

 $A_1$

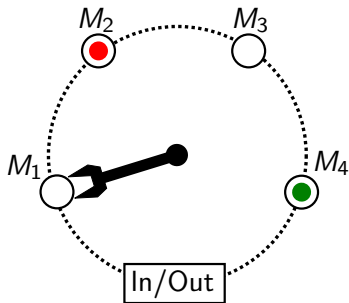
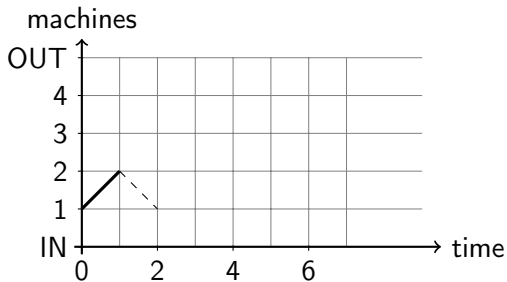
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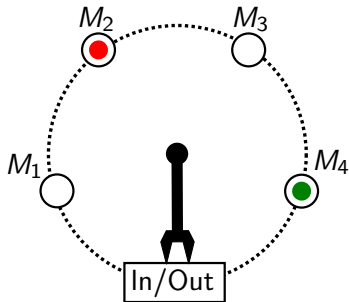
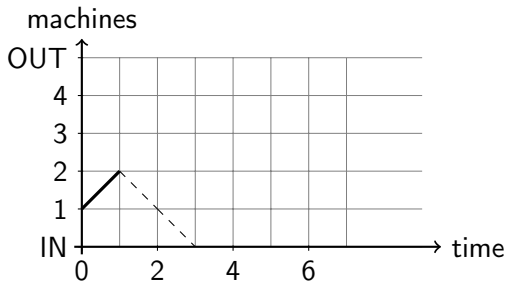
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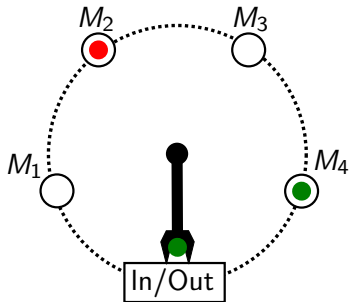
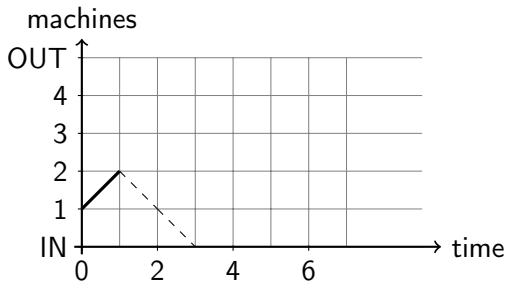
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 $A_1 A_0$

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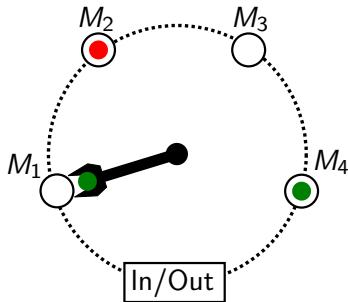
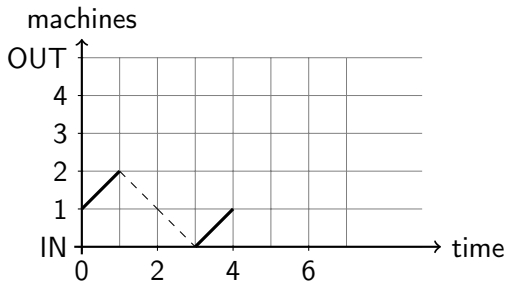
 $A_1 A_0$

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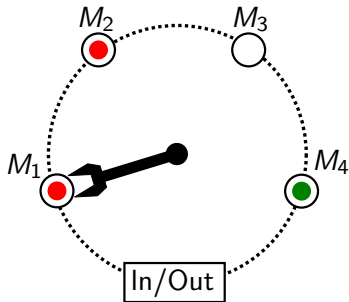
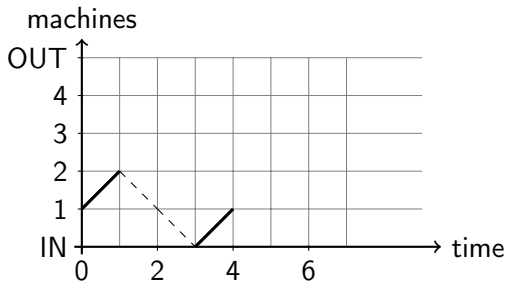
 $A_1 A_0$



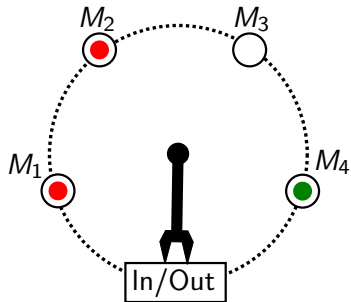
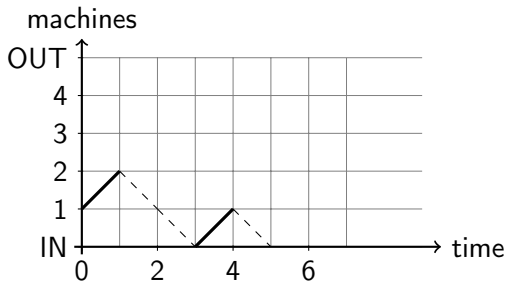
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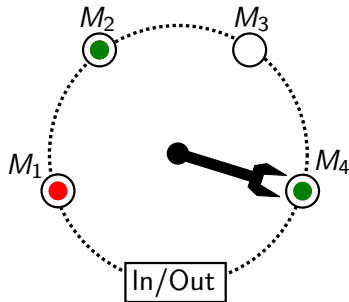
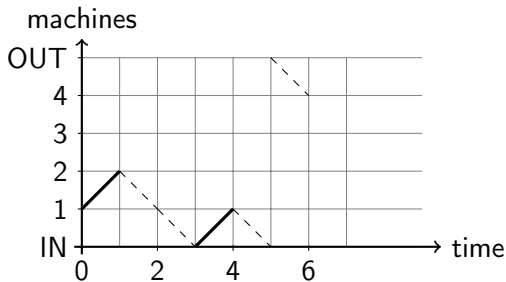
 $A_1 A_0$

# Representing the robot moves



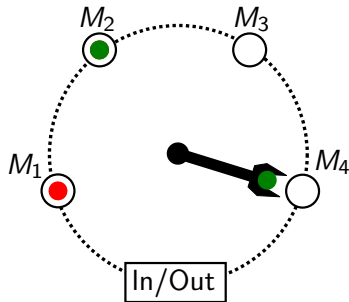
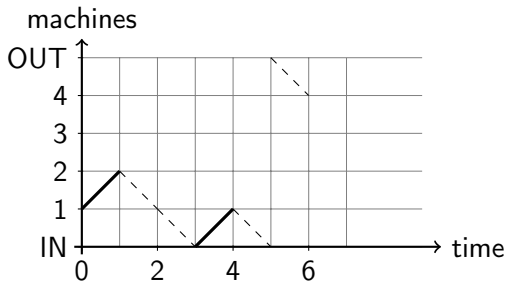
$A_1 A_0 A_4$

# Representing the robot moves



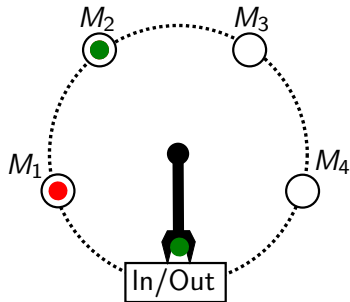
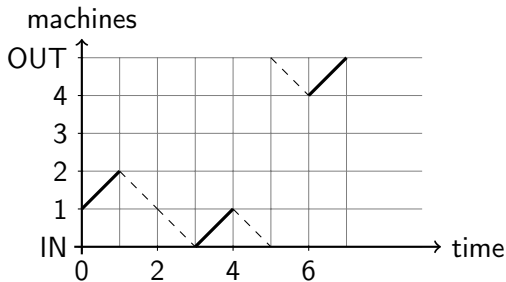
$A_1 A_0 A_4$

# Representing the robot moves



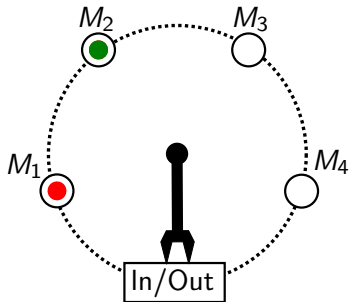
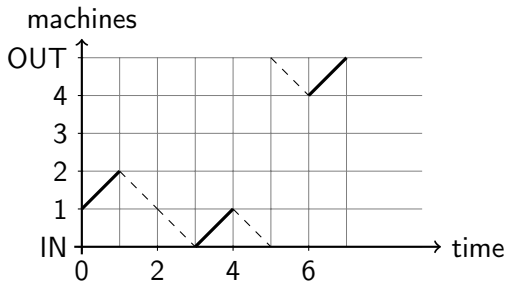
$A_1 A_0 A_4$

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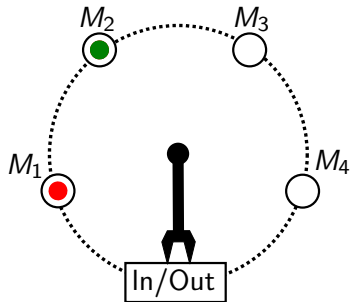
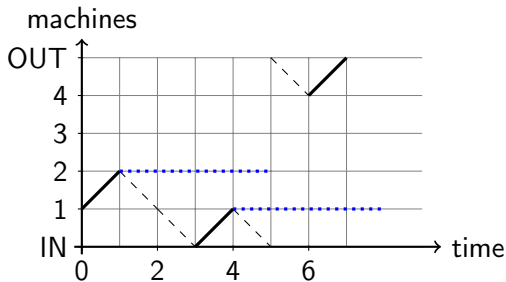


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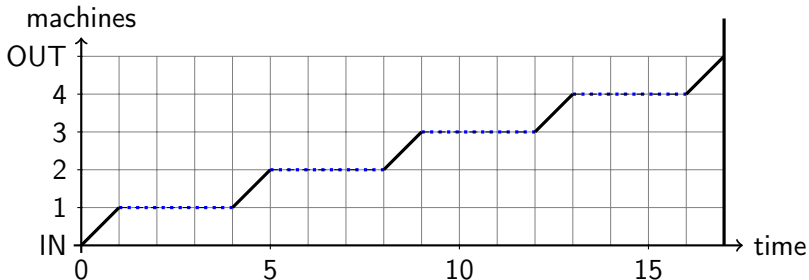

 $A_1 A_0 A_4$



## Some usual 1-cycles: Identity cycle

### Identity cycle $\pi_{id}$

$$\pi_{id}: A_0 A_1 \dots A_m$$



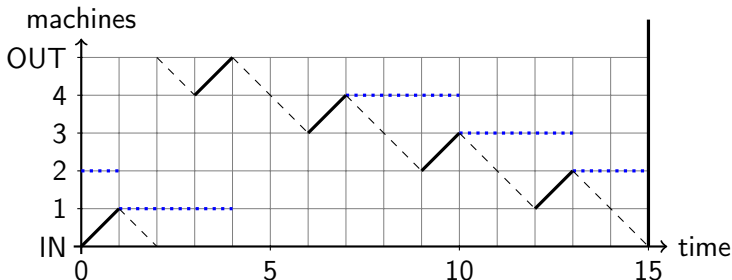
### Cycle time

$$T(\pi_{id}) = (m + 1)\delta + mp$$

# Some usual 1-cycles: Downhill cycle

## Downhill cycle $\pi_d$

$$\pi_d: A_0 A_m A_{m-1} \dots A_1$$



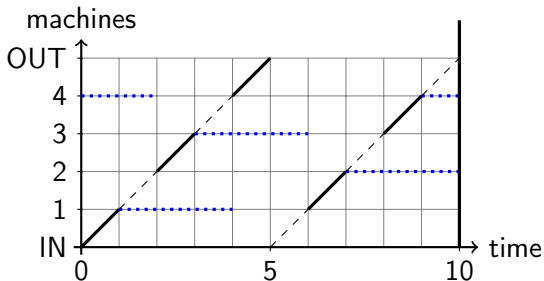
## Cycle Time

$$T(\pi_d) = 3(m+1)\delta + \max(0, p - (3m-1)\delta)$$

# Some usual 1-cycles: Odd-even cycle

## Odd-Even cycle $\pi_{oe}$

$\pi_{oe}: A_0A_2 \dots A_1A_3 \dots$



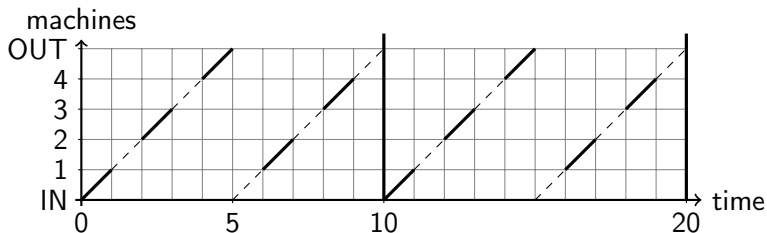
## Cycle time

$T(\pi_{oe}) = ??$

# Odd-even's cycle time

$$T(\pi_{oe}) = 2(m+1)\delta + \max\left(0, \frac{2\alpha-1}{\alpha}(p - (m+1)\delta)\right) \quad \alpha = \lfloor \frac{m+1}{2} \rfloor$$

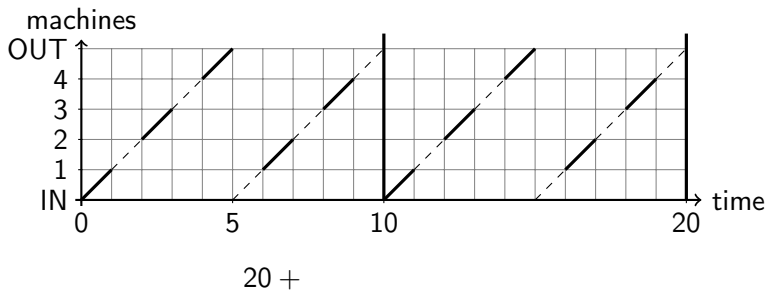
$$\leq 2p$$



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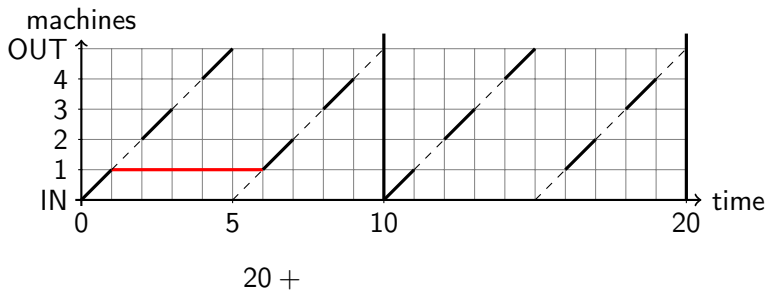
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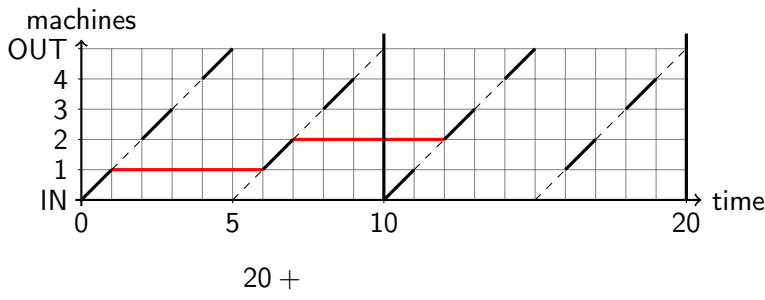
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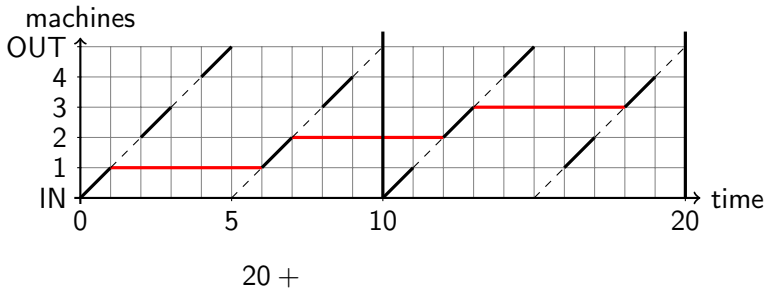
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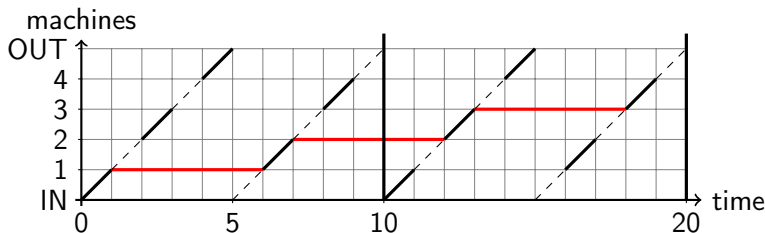




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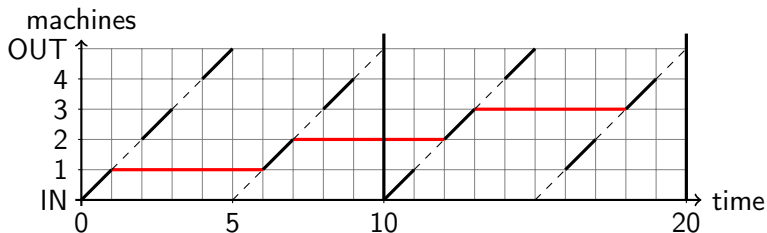


$$\underline{20 + 3 \max(0, (p - 5\delta))}$$

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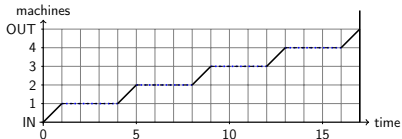
$$\leq 2p$$



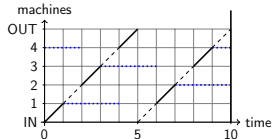
$$\frac{20 + 3 \max(0, (p - 5\delta))}{2}$$

# Finding the best 1-cycle

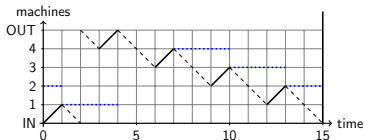
Identity



Odd-Even



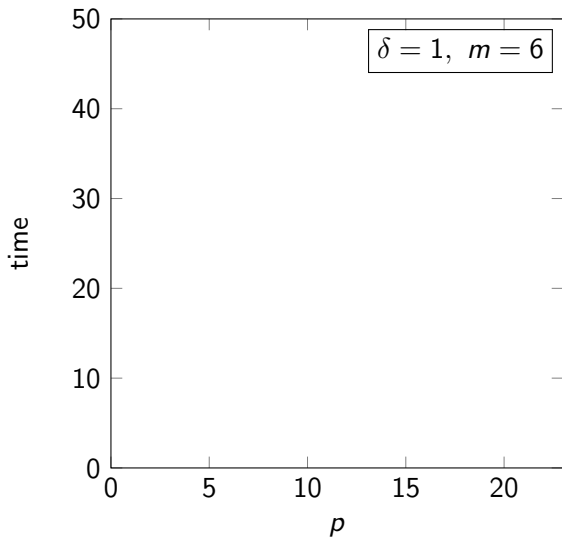
Downhill



?

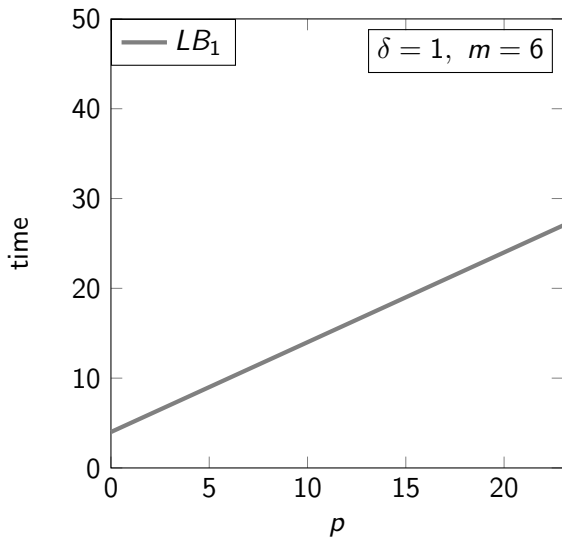
Which is the best 1-cycle?

# Best 1-cycles?



Lower bounds

# Best 1-cycles?

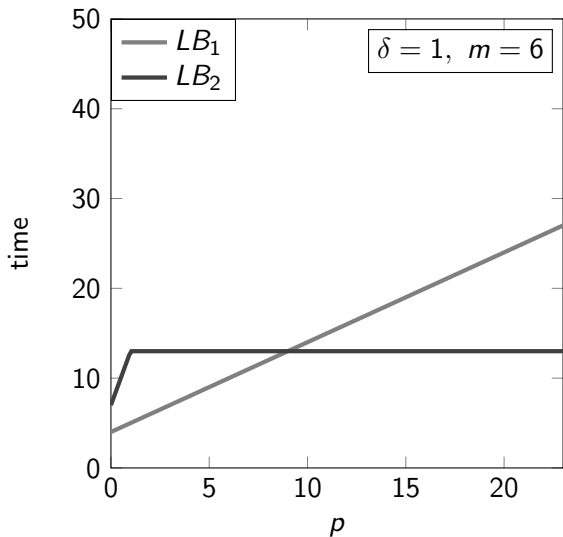


Lower bounds

$LB_1$  [Crama, 97]

$$T(\pi) \geq p + 4\delta$$

## Best 1-cycles?



Lower bounds

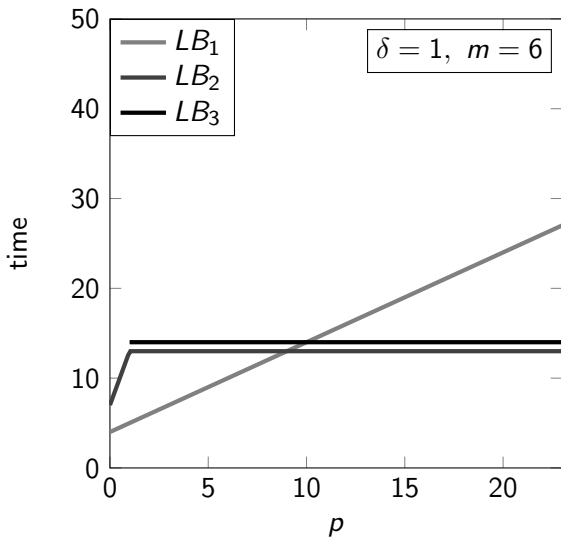
$LB_1$  [Crama, 97]

$$T(\pi) \geq p + 4\delta$$

$LB_2$  [Dawande et al, 02]

$$T(\pi) \geq (m + 1)\delta + m \min(\delta, p)$$

## Best 1-cycles?



Lower bounds

**LB<sub>1</sub>** [Crama, 97]

$$T(\pi) \geq p + 4\delta$$

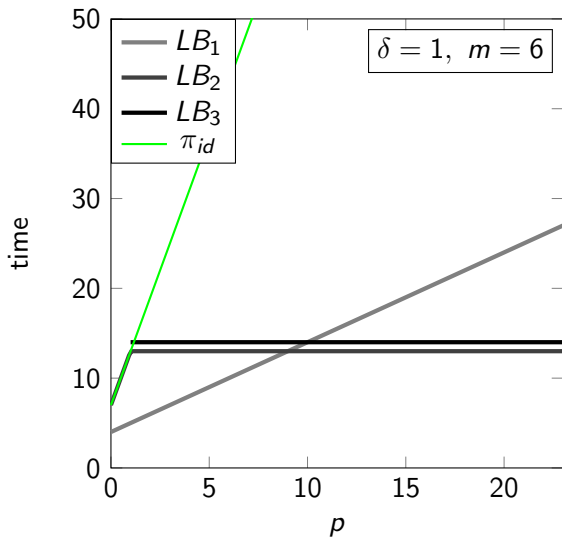
**LB<sub>2</sub>** [Dawande et al, 02]

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**LB<sub>3</sub>** (1-cycles only)

$$T(\pi) \geq 2(m + 1)\delta$$

## Best 1-cycles?



Lower bounds

 $LB_1$  [Crama, 97]

$$T(\pi) \geq p + 4\delta$$

 $LB_2$  [Dawande et al, 02]

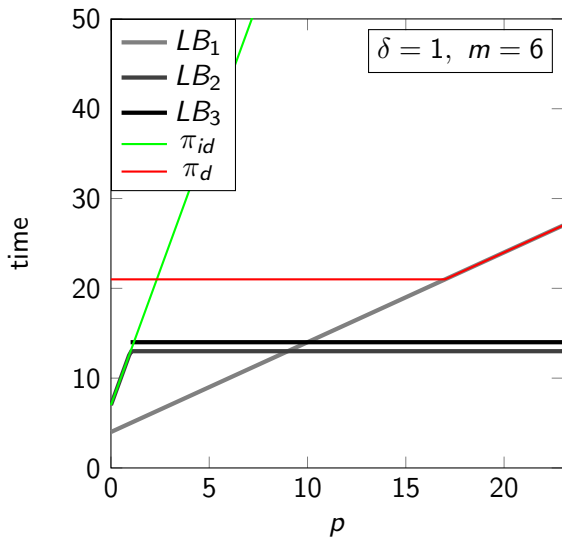
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## Best 1-cycles?



Lower bounds

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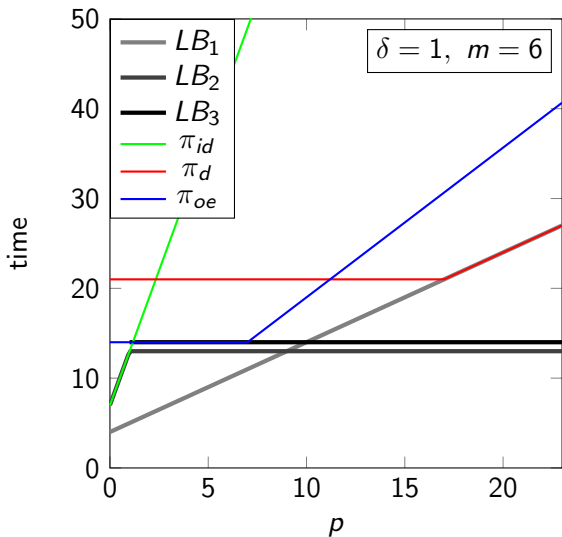
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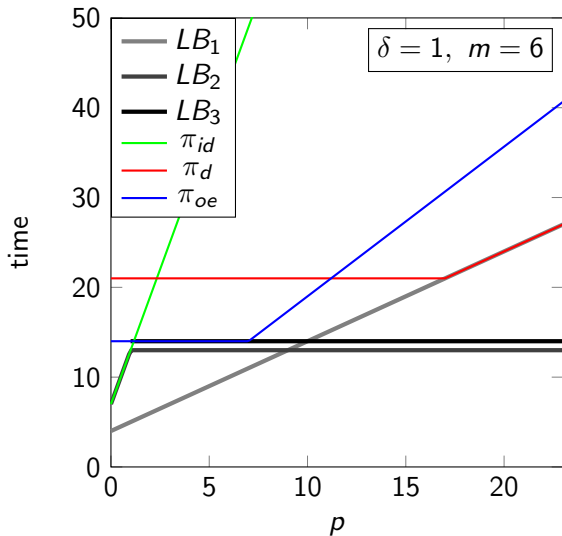
 $LB_2$  [Dawande et al, 02]

$$T(\pi) \geq (m + 1)\delta + m \min(\delta, p)$$

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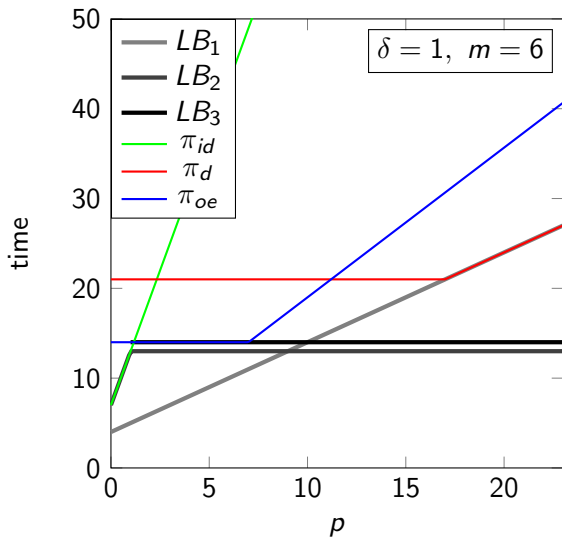
# The unknown region



## Best 1-cycle

- $p \leq \frac{m+1}{m} \delta$   
 $\pi_{id}$

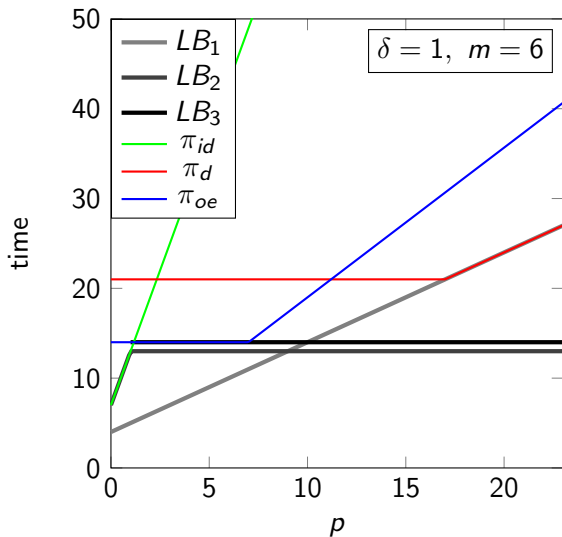
# The unknown region



## Best 1-cycle

- $p \leq \frac{m+1}{m}\delta$   
 $\pi_{id}$
- $p > \frac{m+1}{m}\delta$   
 $p \leq (m+1)\delta$   
 $\pi_{oe}$

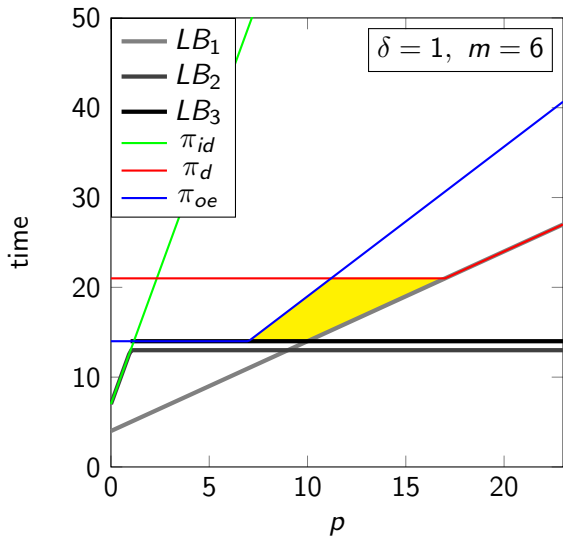
# The unknown region



## Best 1-cycle

- $p \leq \frac{m+1}{m}\delta$   
 $\pi_{id}$
- $p > \frac{m+1}{m}\delta$   
 $p \leq (m+1)\delta$   
 $\pi_{oe}$
- $p \geq (3m-1)\delta$   
 $\pi_d$

# The unknown region



## Best 1-cycle

- $p \leq \frac{m+1}{m}\delta$   
 $\pi_{id}$
- $p > \frac{m+1}{m}\delta$   
 $p \leq (m+1)\delta$   
 $\pi_{oe}$
- $p \geq (3m-1)\delta$   
 $\pi_d$

between  $(m+1)\delta$   
and  $(3m-1)\delta$ ?

## Exploring the unknown region

Best 1-cycles :  $\{\pi_{id}, \pi_{oe}, \pi_d\} + ??$

### Introducing $\pi^*$

- $\pi^*$  is a 1-cycle
- $\pi^*$  crosses the yellow area  
(meaning: for some  $(m + 1)\delta \leq p \leq (3m - 1)\delta$ ,  
 $\pi^*$  does strictly better than both  $\pi_{oe}$  and  $\pi_d$ )

# Exploring the unknown region

Best 1-cycles :  $\{\pi_{id}, \pi_{oe}, \pi_d\} + ??$

## Introducing $\pi^*$

- $\pi^*$  is a 1-cycle
- $\pi^*$  crosses the yellow area  
(meaning: for some  $(m + 1)\delta \leq p \leq (3m - 1)\delta$ ,  
 $\pi^*$  does strictly better than both  $\pi_{oe}$  and  $\pi_d$ )

Who is  $\pi^*$ ?

What can we say about it? Does it exist?



from "Up", Disney - Pixar



# Is there anybody in there?

Size of a minimum dominant set within 1-cycles:

<i>m</i>	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

proven

computed

$$\{\pi_{id}, \pi_{oe}, \pi_d\} \cup ?$$

# Is there anybody in there?

Size of a minimum dominant set within 1-cycles:

<i>m</i>	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

proven

computed

$\{\pi_{id}, \pi_{oe}, \pi_d\} \cup ?$

[Brauner, 99]

# Is there anybody in there?

Size of a minimum dominant set within 1-cycles:

<i>m</i>	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

proven

computed

$\{\pi_{id}, \pi_{oe}, \pi_d\} \cup ?$

Properties on the total travel  
time

# First properties

## Notations (last ones...)

$\Delta(\pi)$ : total travel time

$d_i(\pi)$ : travel time between the loading and unloading of machine  $M_i$

$d_{min}(\pi)$ :  $\min d_i(\pi)$  e.g.  $d_{min}(\pi_{oe}) = (m + 1)\delta$

## Lower bound $LB_4(\pi)$

$$T(\pi) \geq \underbrace{\Delta(\pi)}_{\text{Total travel time}} + \underbrace{\max(0, p - d_{min}(\pi))}_{\text{Minimum waiting time}}$$

$d_{min}$  is the minimum value of  $p$  for which waiting is necessary.

# First properties

If  $\pi^*$  exists:

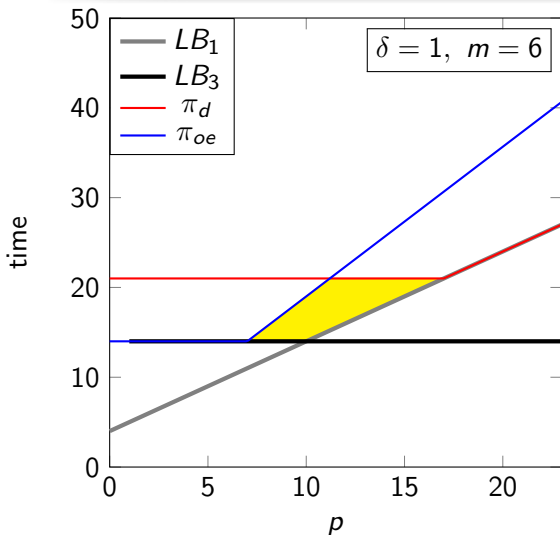
- The robot travels between 2 and 3 times the size of the cell:

$$2(m+1)\delta < \Delta(\pi^*) < 3(m+1)\delta$$

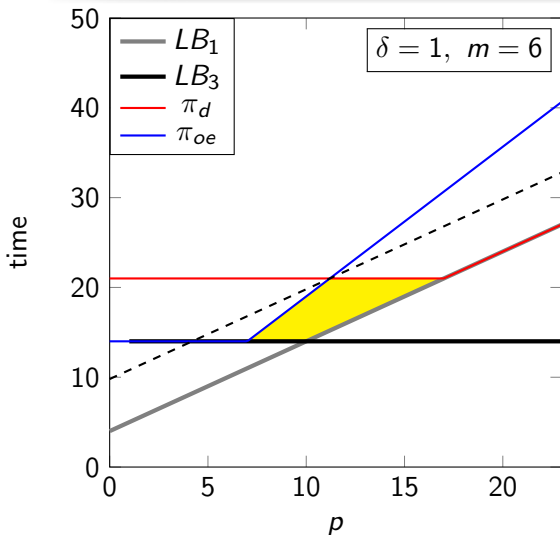
- No waiting time if  $p$  is “small” relatively to the travel time:

$$d_{min}(\pi^*) > \Delta(\pi^*) - \frac{3\alpha - 2}{2\alpha - 1}(m+1)\delta$$

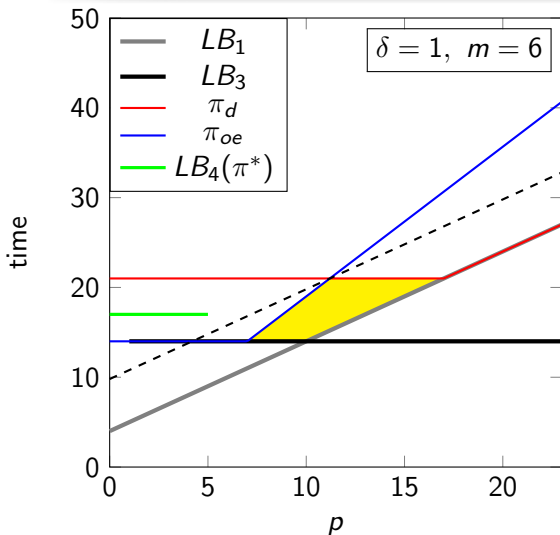
## First properties



# First properties

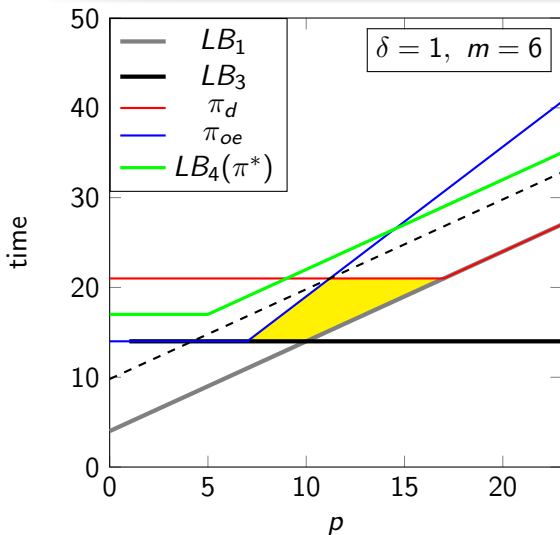


# First properties

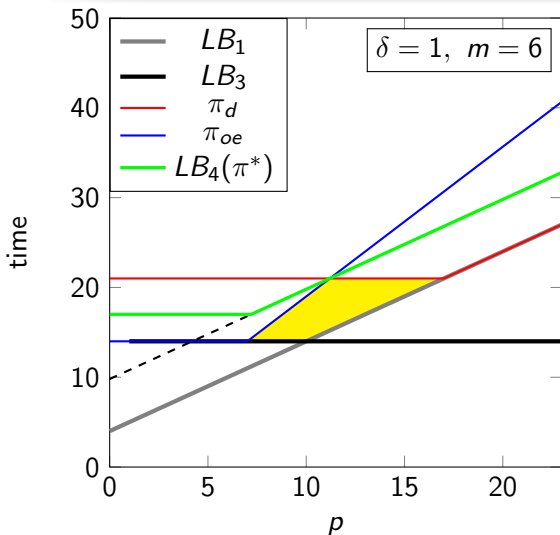




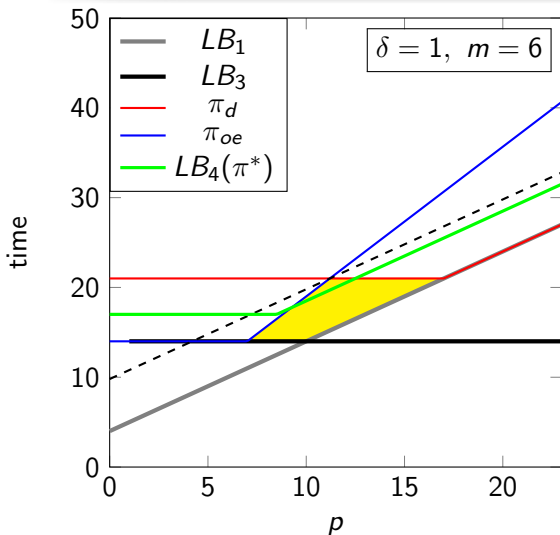
# First properties



# First properties



# First properties



And for  $m > 8$ ?

Cardinal of a minimum dominant set within 1-cycles:

$m$	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

proven

computed

And for  $m > 8$ ?

Cardinal of a minimum dominant set within 1-cycles:

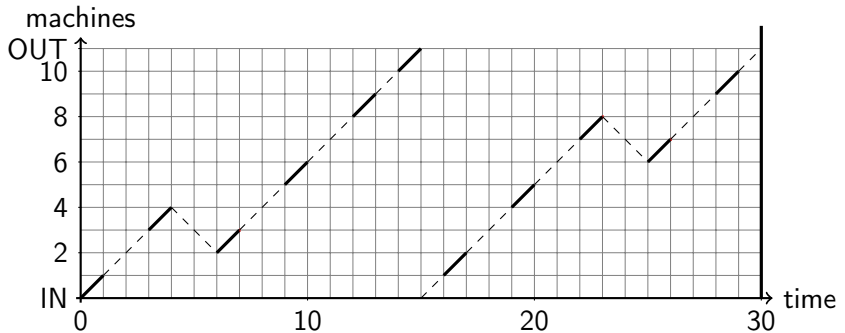
$m$	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

proven

computed

## Two wavelets ( $\pi_{2w}$ )...

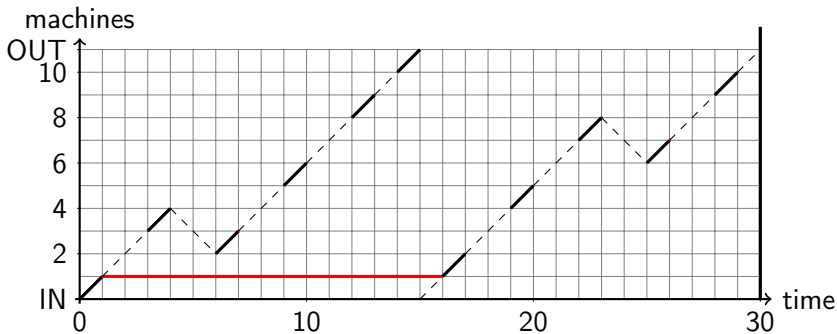
For  $m > 8$ , this guy...



...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...

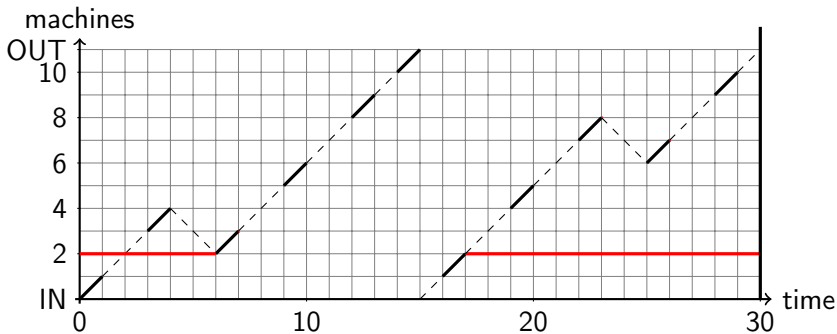
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## Two wavelets ( $\pi_{2w}$ )...

For  $m > 8$ , this guy...

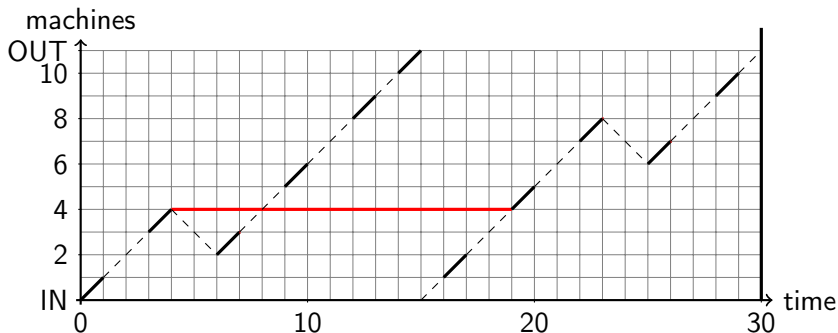


...always (slightly) crosses the yellow area



## Two wavelets ( $\pi_{2w}$ )...

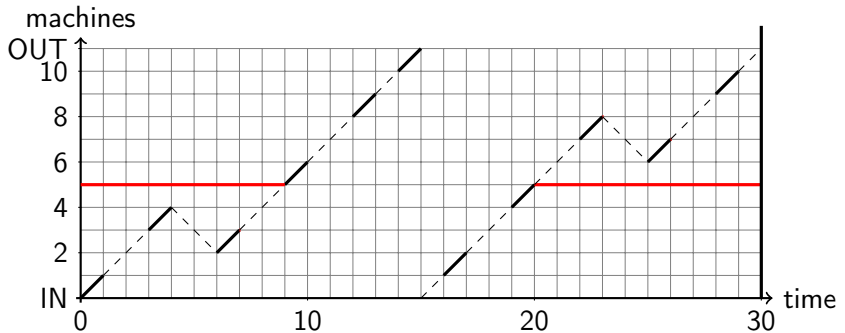
For  $m > 8$ , this guy...



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## Two wavelets ( $\pi_{2w}$ )...

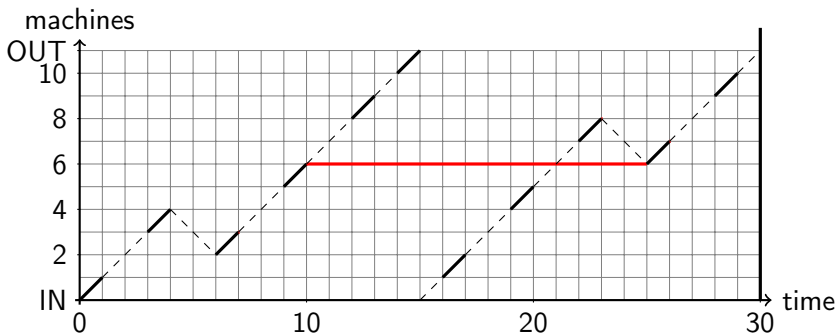
For  $m > 8$ , this guy...



...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...

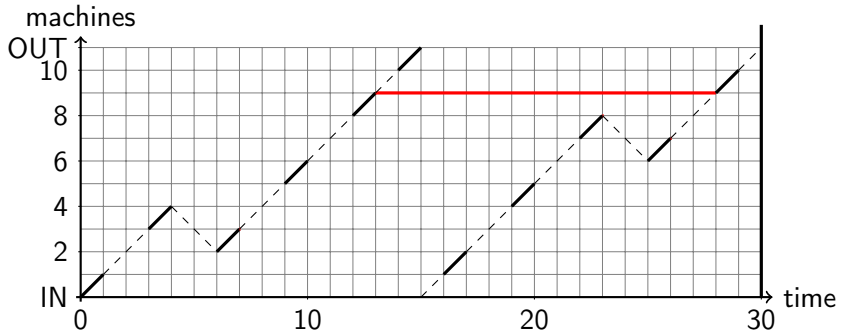
For  $m > 8$ , this guy...



...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...

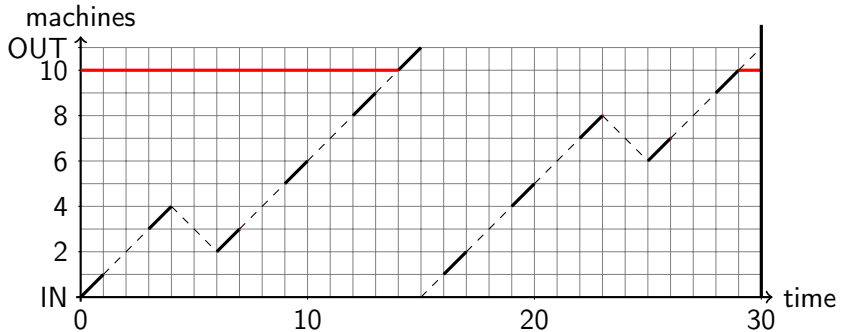
For  $m > 8$ , this guy...



...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...

For  $m > 8$ , this guy...



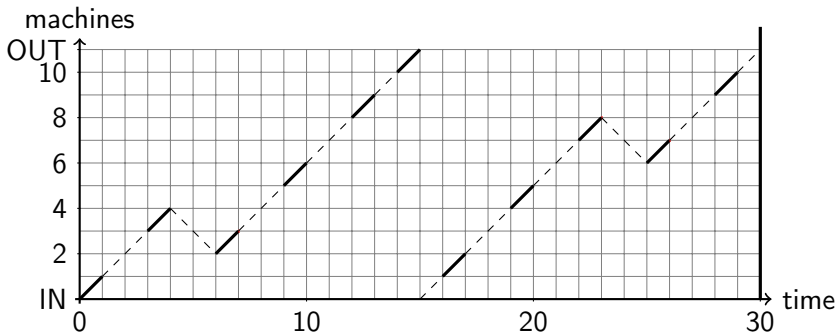
...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...

For  $m > 8$ , this guy...

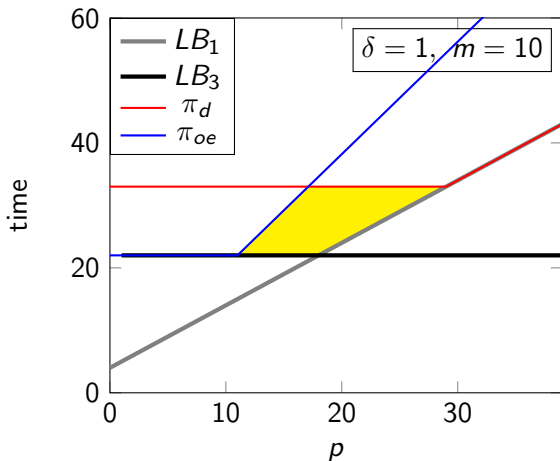
$$d_{\min}(\pi_{2w}) = (m + 5)\delta$$

(One turn and one wave)

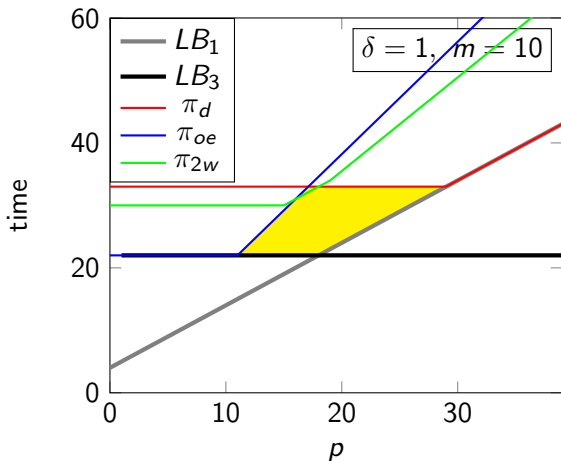


...always (slightly) crosses the yellow area

## Two wavelets ( $\pi_{2w}$ )...



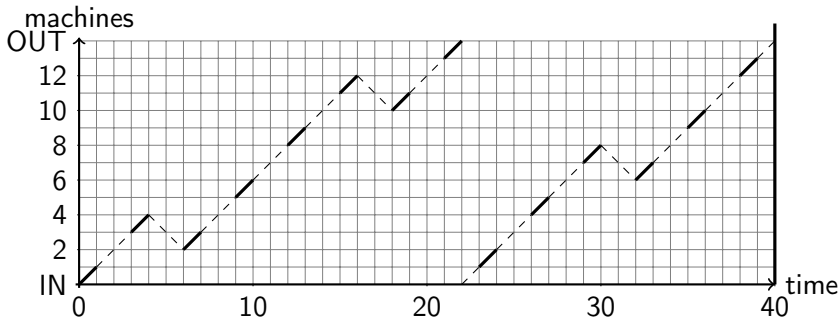
# Two wavelets ( $\pi_{2w}$ )...





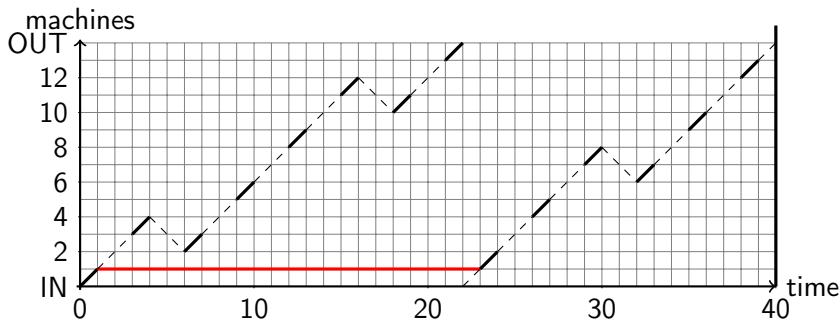
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



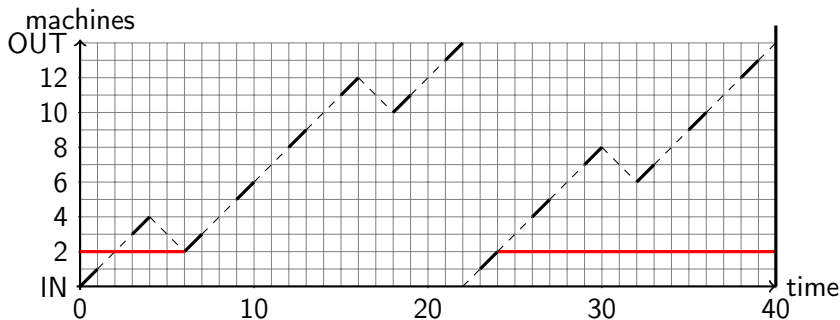
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



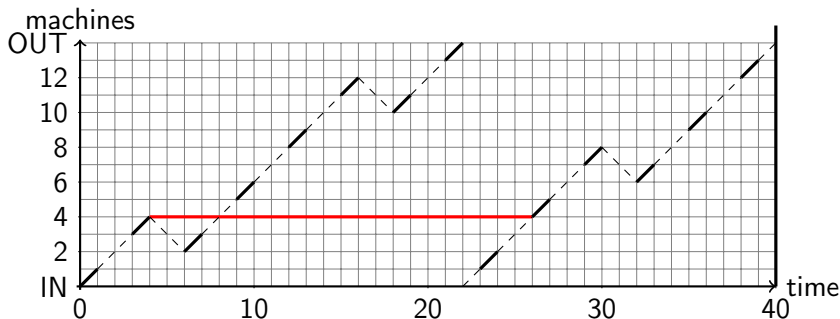
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



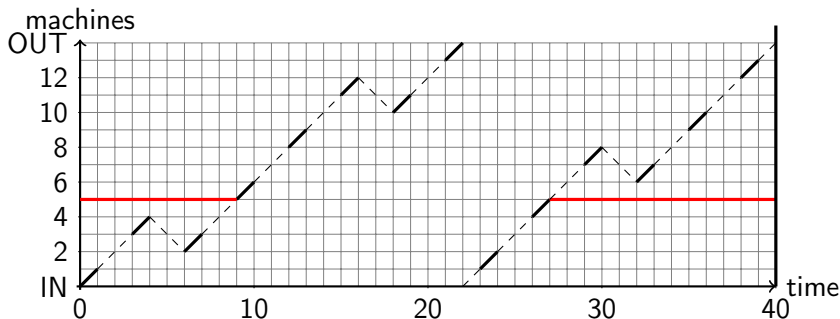
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



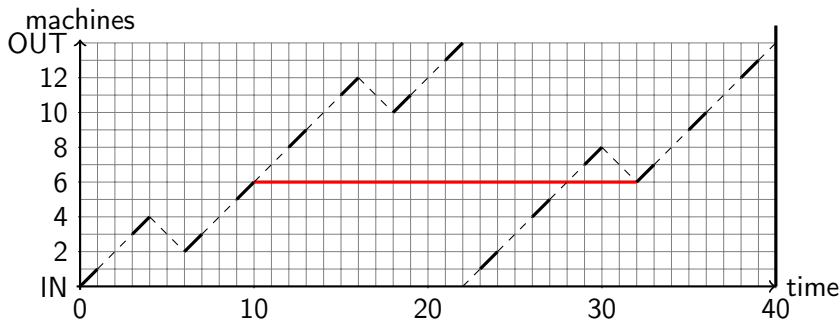
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



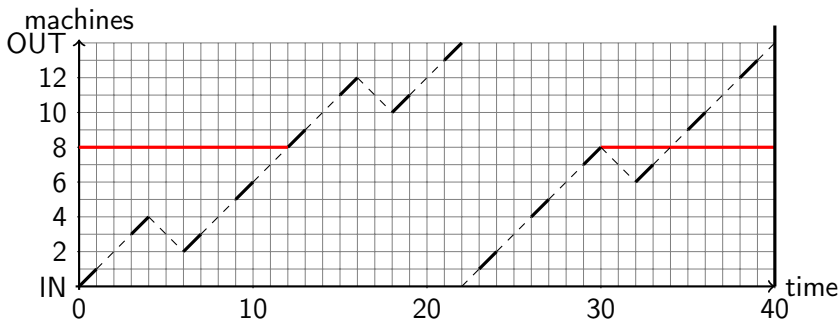
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



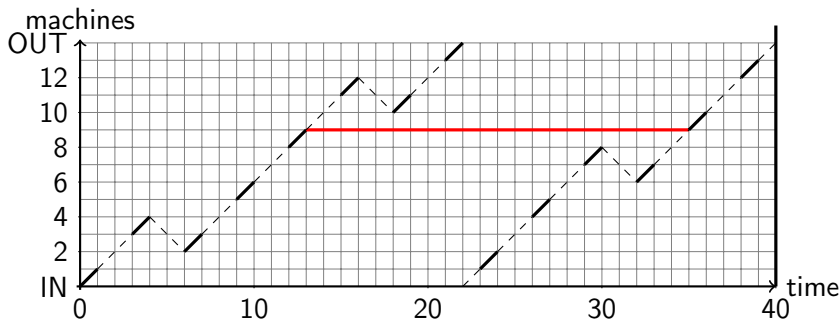
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



# Three wavelets ( $\pi_{3w}$ )...

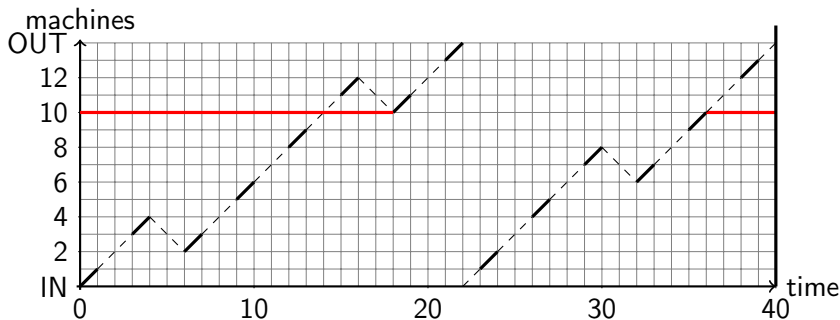
For  $m = 12$  and  $m = 13$ ...





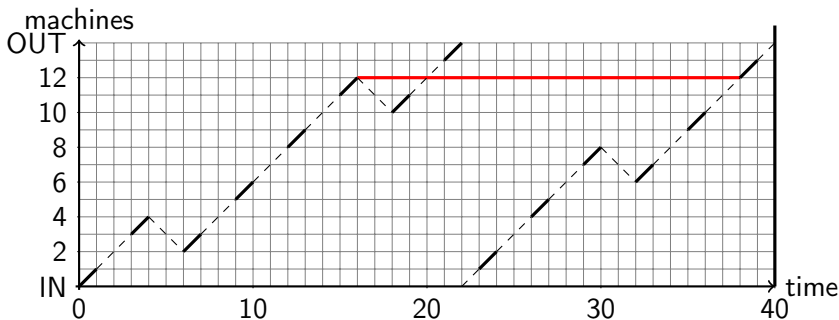
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



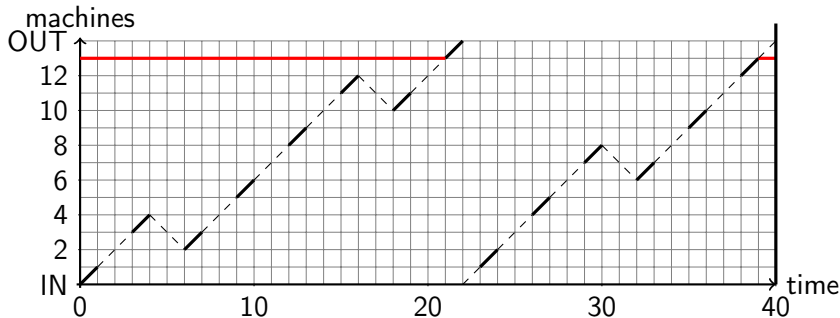
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...



# Three wavelets ( $\pi_{3w}$ )...

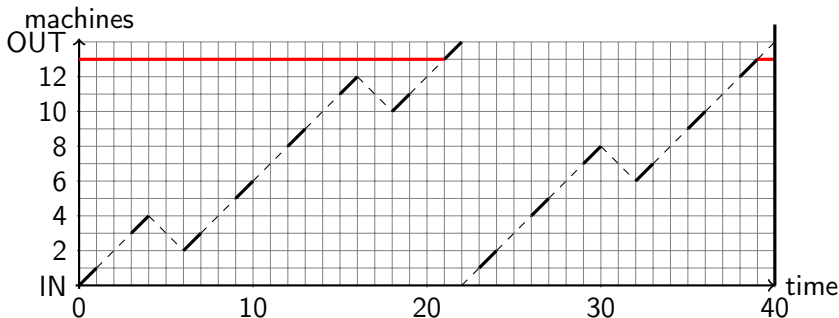
For  $m = 12$  and  $m = 13$ ...



# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...

$d_{min}(\pi_{3w}) = (m + 9)\delta$   
(One turn and two waves)



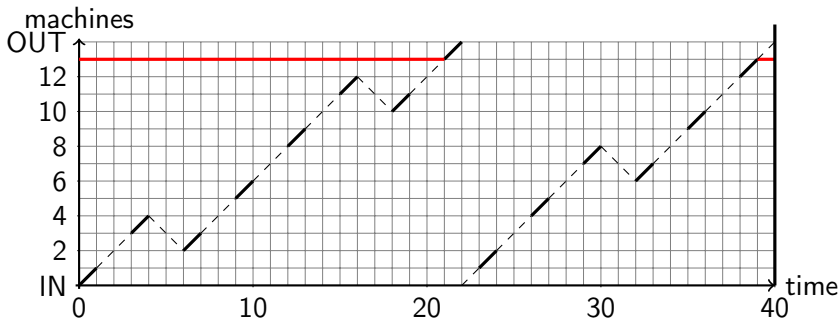
# Three wavelets ( $\pi_{3w}$ )...

For  $m = 12$  and  $m = 13$ ...

False for  $m \geq 14$ !

$$d_{min}(\pi_{3w}) = (m + 9)\delta$$

(One turn and two waves)

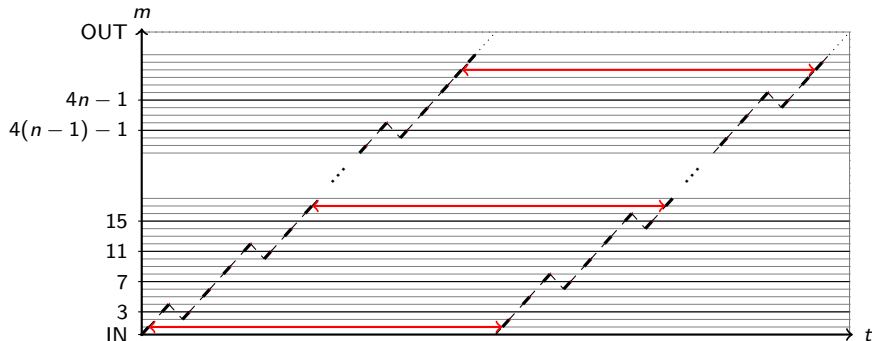


# ... $n$ wavelets ( $\pi_{nw}$ )

(here for  $n$  even)

$$d_{min}(\pi_{nw}) = (m + 2n + 1)\delta$$

(One turn and  $\frac{n}{2}$  waves)



# G-SCOP $m \geq 8$ : A proposition and a conjecture

Proposition: best cycles within  $\{\pi_{id}, \pi_d, \pi_{oe}\} \cup (\pi_{nw})_n$

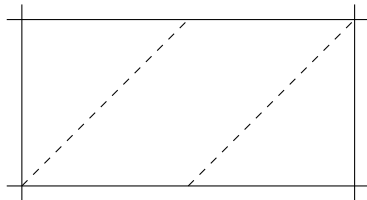
- **Three** classical 1-cycles  $\pi_{id}, \pi_d, \pi_{oe}$
- **+One**  $\pi_{nw}$  with the highest **even** value of  $n$  possible  
And/or
- **+One**  $\pi_{nw}$  with the highest **odd** value of  $n$  possible

Conjecture: best 1-cycle

These also dominate **all** 1-cycles

Proven for  $m \leq 11$

- 2 “turns”



$$T(\pi^*) < 2p$$

(otherwise, dominated by  $\pi_{oe}$ )

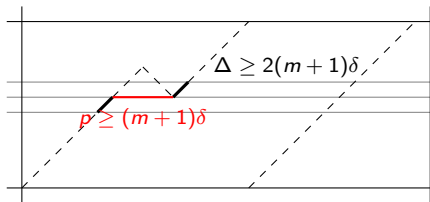
$$T(\pi^*) < 3(m+1)\delta$$

(otherwise, dominated by  $\pi_d$ )



# Proof ideas

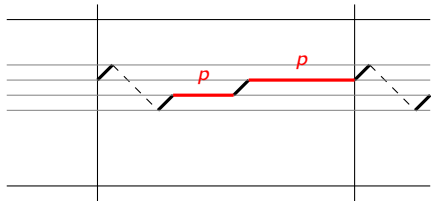
- 2 “turns”
- $A_i$  and  $A_{i+1}$  can't be on the same turn in that order



$T(\pi^*) < 2p$   
 (otherwise, dominated by  $\pi_{oe}$ )  
 $T(\pi^*) < 3(m+1)\delta$   
 (otherwise, dominated by  $\pi_d$ )

# Proof ideas

- 2 “turns”
- $A_i$  and  $A_{i+1}$  can't be on the same turn in that order
- no subsequence  $A_i A_{i-2}$   
(rules out bigger alterations up to 11 machines)



$$T(\pi^*) < 2p$$

(otherwise, dominated by  $\pi_{oe}$ )

$$T(\pi^*) < 3(m+1)\delta$$

(otherwise, dominated by  $\pi_d$ )

# Best 1 cycles for $m \leq 11$ ...

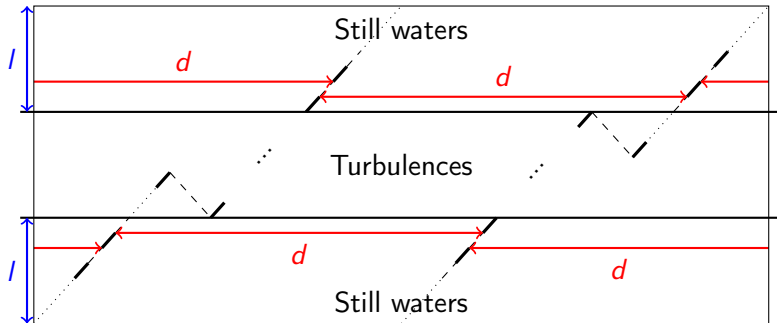
Cardinal of a minimum dominant set within 1-cycles:

$m$	3	4	5	6	7	8	9	10	11	12	13	14
<i>cardinal</i>	3	4	4	3	3	3	4	4	4	5	4	4

$$6 \leq m \leq 8: \{\pi_{id}, \pi_{oe}, \pi_d\}$$

$$9 \leq m \leq 11: \{\pi_{id}, \pi_{oe}, \pi_d, \pi_{2w}\}$$

# And for any $m$ ?

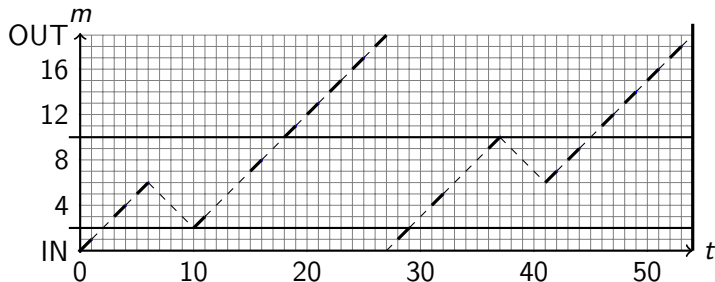


Lower bound on the waiting time:  $\nearrow l$ ,  $\searrow d$   
 (tight for  $p_{nw}$  cycles)

# And for any $m$ ?

**Not proven** Regular disposition of the waves in the “turbulence” areas: seems intuitive but...

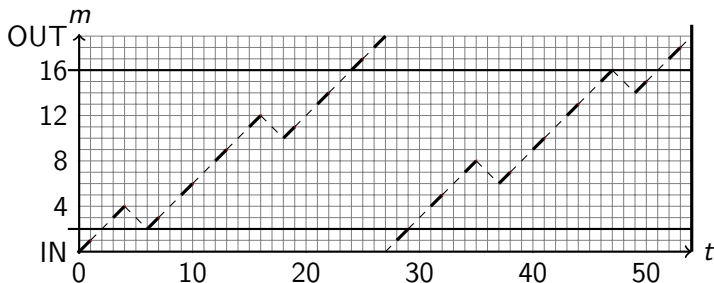
**Proven** Wavelets preferable to big waves:



## And for any $m$ ?

**Not proven** Regular disposition of the waves in the “turbulence” areas: seems intuitive but...

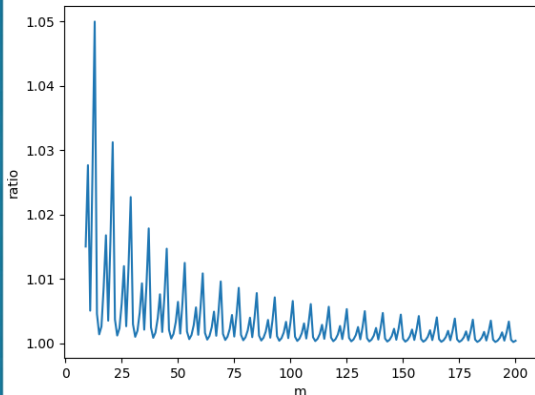
**Proven** Wavelets preferable to big waves:



## If the conjecture is valid...

If the conjecture is valid, then...

- The best 1-cycle problem in circular, regular balanced cell would be **polynomial**
- Performance ratio of the usual cycles:



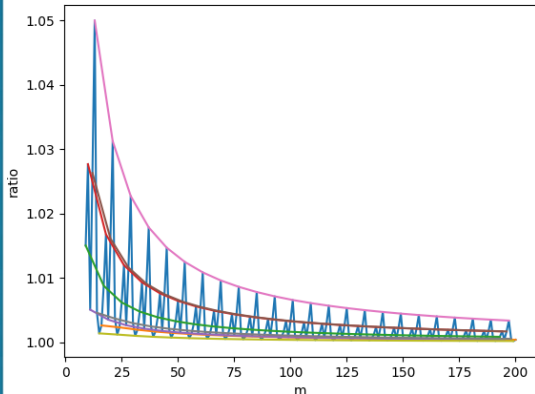
$$\frac{\min(T(\pi_{oe}), T(\pi_d))}{\min_n(T(\pi_{oe}), T(\pi_d), T(\pi_{nw}))}$$

(as a function of  $m$ )

## If the conjecture is valid...

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- The best 1-cycle problem in circular, regular balanced cell would be **polynomial**
- Performance ratio of the usual cycles:



$$\frac{\min(T(\pi_{oe}), T(\pi_d))}{\min_n(T(\pi_{oe}), T(\pi_d), T(\pi_{nw}))}$$

(as a function of  $m$ )



## Conclusions...

Throughput optimization in robotic cells with **circular layout** was less studied and poorly understood so far. We:

- proposed new tools and specific cycle structures...
- ... leading to a **conjecture** on the best 1-cycle problem.

	Best 1-cycle	1-cycle conjecture
Linear	P	$m \leq 3$ : valid $m = 4$ : false
balanced	P	$m \leq 15$ : valid
Circular	NP-hard	
balanced	$m \leq 11$ : P $m \geq 12$ : also P?	$m = 6$ : false

(Well, aside from settling the conjecture)

Other types of production constraints:

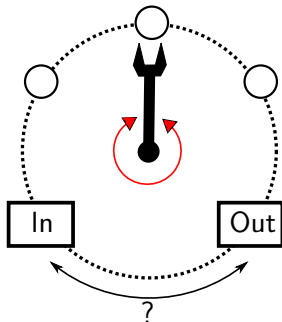
- Non-balanced case:  
Improving existing approximation...
- Proportionate flow-shop

Open questions for regular balanced cells:

- Best 1-cycle for  $m > 11$ ...
- 1-cycle conjecture for  $m \leq 5$
- Cycle function for  $m \geq 6$

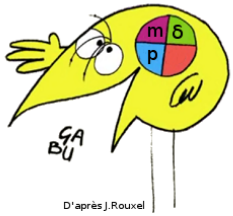
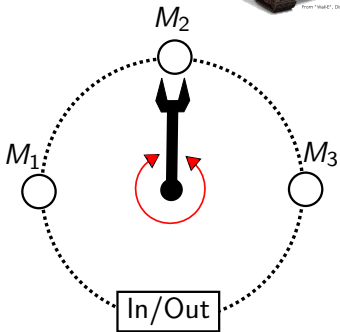
Relationships with other layouts:

- Comparisons of layouts
- Generalization of the circular layout:





Thank you!



D'après J. Rouxel



## Third party content

- This presentation features some tributes to the following works (short excerpts):
  - *Wall-E* (Pixar Animation Studios, 2008), slides 1, 4, 8 and 43;
  - *Up* (Pixar Animation Studios & Walt Disney Pictures, 2009), slides 25 and 43;
  - *Les Shadoks* (Jacques Rouxel), slides 16 and 43.
- The picture of a toy train featured on slides 1, 2, 3 and 43 is extracted from a Brio commercial.
- Other clipart images (slides 4 and 10) are either public domain or released under a CC0 license.
- The template and backgrounds belong to G-SCOP laboratory.