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Thèse de doctorat en Mathématiques

Dirigée par François Métayer et par Clemens Berger

Homology of strict ω -categories

par Léonard Guetta

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Résumé

Dans cette thèse, on compare l'homologie « classique » d'une ω -catégorie (définie comme l'homologie de son nerf de Street) avec son homologie polygraphique. Plus précisément, on prouve que les deux homologies ne coïncident pas en général et qualifions d'homologiquement cohérentes les ω -catégories particulières pour lesquelles l'homologie polygraphique coïncide effectivement avec l'homologie du nerf. Le but poursuivi est de trouver des critères abstraits et concrets permettant de détecter les ω -catégories homologiquement cohérentes. Par exemple, on démontre que toutes les (petites) catégories, que l'on considère comme des ω -catégories strictes dont toutes les cellules au-delà de la dimension 1 sont des unités, sont homologiquement cohérentes. On introduit également la notion de 2-catégorie sans bulles et on conjecture qu'une 2-catégorie cofibrante est homologiquement cohérente si et seulement si elle est sans bulles. On démontre également des résultats importants concernant les ω -catégories strictes qui sont libres sur un polygraphe, comme le fait que si $F: C \to D$ est un ω -foncteur de Conduché discret et si D est libre sur un polygraphe alors C l'est aussi. Dans son ensemble, cette thèse établit un cadre général dans lequel étudier l'homologie des ω -catégories en faisant appel à des outils d'algèbre homotopique abstraite, tels que la théorie des catégories de modèles de Quillen ou la théorie des dérivateurs de Grothendieck.

Mots-clés : Catégories supérieures, ω -catégories, homologie, théorie de l'homotopie, polygraphes.

Abstract

In this dissertation, we compare the "classical" homology of an ω -category (defined as the homology of its Street nerve) with its polygraphic homology. More precisely, we prove that both homologies generally do not coincide and call *homologically coherent* the particular strict ω -categories for which polygraphic homology and homology of the nerve do coincide. The goal pursued is to find abstract and concrete criteria to detect homologically coherent ω -categories. For example, we prove that all (small) categories, considered as strict ω -categories with unit cells above dimension 1, are homologically coherent. We also introduce the notion of *bubble-free* 2-category and conjecture that a cofibrant 2-category is homologically coherent if and only if it is bubble-free. We also prove important results concerning free strict ω -categories on polygraphs (also known as computads), such as the fact that if $F : C \to D$ is a discrete Conduché ω -functor and D is a free strict ω -category on a polygraph, then so is C. Overall, this thesis achieves to build a general framework in which to study the homology of strict ω -categories using tools of abstract homotopical algebra such as Quillen's theory of model categories or Grothendieck's theory of derivators.

Keywords : Higher categories, ω -categories, homology, homotopy theory, polygraphs.

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INTRODUCTION

The general framework in which this dissertation takes place is the *homotopy theory of strict* ω -*categories*, and, as the title suggests, its focus is on homological aspects of this theory. The goal is to study and compare two different homological invariants for strict ω -categories; that is to say, two different functors

$$\mathbf{Str}\omega\mathbf{Cat} \to \mathrm{ho}(\mathbf{Ch}_{\geq 0})$$

from the category of strict ω -categories to the homotopy category of non-negatively graded chain complexes (i.e. the localization of the category of non-negatively graded chain complexes with respect to the quasi-isomorphisms).

Before we enter into the heart of the subject, let us emphasize that, with the sole exception of the end of this introduction, all the ω -categories that we consider are strict ω -categories. Hence, we drop the adjective "strict" and simply say ω -category instead of strict ω -category and we write ω Cat instead of Str ω Cat for the category of (strict) ω -categories.

Background: ω -categories as spaces. The homotopy theory of ω -categories begins with the nerve functor introduced by Street in [Str87]

$$N_{\omega}: \omega \mathbf{Cat} \to \widehat{\Delta}$$

that associates to every ω -category C a simplicial set $N_{\omega}(C)$ called the *nerve of* C, generalizing the usual nerve of (small) categories. Using this functor, we can transfer the homotopy theory of simplicial sets to ω -categories, as it is done for example in the articles [AM14, AM18, Gag18, Ara19, AM20c, AM20a]. Following the terminology of these articles, a morphism $f: C \to D$ of ω Cat is a *Thomason equivalence* if $N_{\omega}(f)$ is a Kan–Quillen weak equivalence of simplicial sets. By definition, the nerve functor induces a functor at the level of homotopy categories

$$\overline{N_{\omega}}$$
: ho($\omega \mathbf{Cat}^{\mathrm{Th}}$) \rightarrow ho($\widehat{\Delta}$),

where $ho(\omega Cat^{Th})$ is the localization of ωCat with respect to the Thomason equivalences and $ho(\widehat{\Delta})$ is the localization of $\widehat{\Delta}$ with respect to the Kan–Quillen weak equivalences of simplicial sets. As it so happens, the functor $\overline{N_{\omega}}$ is an equivalence of categories, as proved by Gagna in [Gag18]. In other words, the homotopy theory of ω -categories induced by Thomason equivalences is the same as the homotopy theory of spaces.

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Gagna's result is in fact a generalization of the analogous result for the usual nerve of small categories, which is attributed to Quillen in [III72]. In the case of small categories, Thomason even showed the existence of a model structure whose weak equivalences are the ones induced by the nerve functor [Tho80]. The analogous result for ω Cat is conjectured but not yet established [AM14].

Two homologies for ω -categories. Keeping in mind the nerve functor of Street, a natural thing to do is to define the *k*-th homology group of an ω -category *C* as the *k*-th homology group of the nerve of *C*. In light of Gagna's result, these homology groups are just another way of looking at the homology groups of spaces. In order to explicitly avoid future confusion, we shall now use the name singular homology groups of *C* for these homology groups and the notation $H_k^{\text{Sing}}(C)$.

On the other hand, Métayer gives a definition in [Mét03] of other homology groups for ω -categories. This definition is based on the notion of ω -categories free on a polygraph (also known as ω -categories free on a computad), which are ω -categories that are obtained from the empty category by recursively freely adjoining cells. From now on, we simply say free ω -category. Métayer observed that every ω -category C admits what we call a polygraphic resolution, which means that there exists a free ω -category P and a morphism of ω Cat

$$f: P \to C$$

that satisfies properties formally resembling those of trivial fibrations of topological spaces (or of simplicial sets). Furthermore, every free ω -category P can be "abelianized" to a chain complex $\lambda(P)$ and Métayer proved that for two different polygraphic resolutions $P \to C$ and $P' \to C$ of the same ω -category C, the chain complexes $\lambda(P)$ and $\lambda(P')$ are quasi-isomorphic. Hence, we can define the *k*-th polygraphic homology group of C, denoted by $H_k^{\text{pol}}(C)$, as the *k*-th homology group of $\lambda(P)$ for any polygraphic resolution $P \to C$.

One is then led to the following question:

Do we have
$$H^{\text{pol}}_{\bullet}(C) \simeq H^{\text{Sing}}_{\bullet}(C)$$
 for every ω -category C?

A first partial answer to this question is given by Lafont and Métayer in [LM09]: for a monoid M (seen as category with one object and hence as an ω -category), we have $H_{\bullet}^{\text{pol}}(M) \simeq H_{\bullet}^{\text{Sing}}(M)$. In fact, the original motivation for polygraphic homology was the homology of monoids and is part of a program that generalizes to higher dimension the results of Squier on the rewriting theory of monoids [Gui06, Laf07, GM09, GM18]. However, interestingly enough, the general answer to the above question is *no*. A counterexample was found by Maltsiniotis and Ara. Let B be the commutative monoid ($\mathbb{N}, +$), seen as a 2-category with only one 0-cell and no non-trivial 1-cells. This 2-category is free (as an ω -category) and a quick computation shows that:

$$H_k^{\text{pol}}(B) = \begin{cases} \mathbb{Z} & \text{ if } k = 0, 2\\ 0 & \text{ otherwise.} \end{cases}$$

On the other hand, it is shown in [Ara19, Theorem 4.9 and Example 4.10] that the nerve of *B* is a $K(\mathbb{Z}, 2)$; hence, it has non-trivial homology groups in all even dimension. A question that still remains is:

question that still remains is.

(Q) Which are the ω -categories C such that $H^{\text{pol}}_{\bullet}(C) \simeq H^{\text{Sing}}_{\bullet}(C)$?

This is precisely the question around which this dissertation revolves. Nevertheless, the reader will also find several new notions and results within this document that, although primarily motivated by the above question, are of interest in the theory of ω -categories and whose *raisons d'être* go beyond the above considerations.

Another formulation of the problem. One of the achievements of the present work is a more abstract reformulation of the question of comparison of singular and polygraphic homology of ω -categories.

In order to do so, recall first that by a variation of the Dold–Kan equivalence (see for example [Bou90]), the category of abelian group objects in ω Cat is equivalent to the category of non-negatively graded chain complexes

$$Ab(\omega Cat) \simeq Ch_{>0}.$$

Hence, we have a forgetful functor $Ch_{\geq 0} \simeq Ab(\omega Cat) \rightarrow \omega Cat$, which has a left adjoint

$$\lambda: \omega \mathbf{Cat} \to \mathbf{Ch}_{>0}.$$

Moreover, for a *free* ω -category C, the chain complex $\lambda(C)$ is exactly the one obtained by the "abelianization" process considered in Métayer's definition of polygraphic homology.

Now, the category ω Cat admits a model structure, known as the *folk model structure* [LMW10], whose weak equivalences are the *equivalences of* ω -categories (a generalization of the usual notion of equivalence of categories) and whose cofibrant objects are exactly the free ω -categories [Mét08]. Polygraphic resolutions are then nothing but cofibrant replacements in this model category. As the definition of polygraphic homology groups strongly suggests, the functor λ is left Quillen with respect to this model structure. In particular, it admits a left derived functor

$$\mathbb{L}\lambda^{\text{folk}} : \text{ho}(\omega \mathbf{Cat}^{\text{folk}}) \to \text{ho}(\mathbf{Ch}_{\geq 0})$$

and we tautologically have that $H_k^{\text{pol}}(C) = H_k(\mathbb{L}\lambda^{\text{folk}}(C))$ for every ω -category C and every $k \ge 0$. From now on, we set

$$\mathbb{H}^{\mathrm{pol}}(C) := \mathbb{L}\lambda^{\mathrm{folk}}(C).$$

This way of understanding polygraphic homology as a left derived functor has been around in the folklore for some time and I claim absolutely no originality for it.

On the other hand, λ is also left derivable when ω Cat is equipped with Thomason equivalences, yielding a left derived functor

$$\mathbb{L}\lambda^{\mathrm{Th}}$$
: ho($\omega \mathbf{Cat}^{\mathrm{Th}}$) \rightarrow ho($\mathbf{Ch}_{>0}$).

This left derived functor being such that $H_k^{\text{Sing}}(C) = H_k(\mathbb{L}\lambda^{\text{Th}}(C))$ for every ω -category C and every $k \ge 0$. Contrary to the "folk" case, this result is new and first appears within this document (at least to my knowledge). Note that since, as mentioned earlier, the existence of a Thomason-like model structure on ω Cat is still conjectural, usual tools from Quillen's theory of model categories were unavailable to prove the left derivability of λ and the difficulty was to find a workaround solution.

From now on, we set

$$\mathbb{H}^{\mathrm{Sing}}(C) := \mathbb{L}\lambda^{\mathrm{Th}}(C).$$

Finally, it can be shown that every equivalence of ω -categories is a Thomason equivalence. Hence, the identity functor of ω Cat induces a functor \mathcal{J} at the level of homotopy categories

$$\mathcal{J}: ho(\omega \mathbf{Cat}^{\mathrm{folk}}) \to ho(\omega \mathbf{Cat}^{\mathrm{Th}}),$$

and altogether we have a triangle

$$\begin{array}{ccc} \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{folk}}) & \xrightarrow{\mathcal{J}} & \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{Th}}) \\ & & \downarrow^{\mathbb{L}\lambda^{\operatorname{folk}}} & \downarrow^{\mathbb{L}\lambda^{\operatorname{Th}}} \\ & & \operatorname{ho}(\mathbf{Ch}_{\geq 0}). \end{array}$$

This triangle is *not* commutative (even up to isomorphism), since this would imply that the singular and polygraphic homology groups coincide for every ω -category. However, since both functors $\mathbb{L}\lambda^{\text{folk}}$ and $\mathbb{L}\lambda^{\text{Th}}$ are left derived functors of the same functor λ , the existence of a natural transformation $\pi : \mathbb{L}\lambda^{\text{Th}} \circ \mathcal{J} \Rightarrow \mathbb{L}\lambda^{\text{folk}}$ follows by universal property. Since \mathcal{J} is the identity on objects, for every ω -category C, this natural transformation yields a map

$$\pi_C : \mathbb{H}^{\operatorname{Sing}}(C) \to \mathbb{H}^{\operatorname{pol}}(C),$$

which we refer to as the *canonical comparison map*. Let us say that C is *homologically* coherent if π_C is an isomorphism (which means exactly that for every $k \ge 0$, the induced map $H_k^{\text{Sing}}(C) \to H_k^{\text{pol}}(C)$ is an isomorphism). The question of study then becomes:

(Q') Which ω -categories are homologically coherent ?

Note that, in theory, question $(\mathbf{Q'})$ is more precise than question (\mathbf{Q}) since we impose which morphism has to be an isomorphism in the comparison of homology groups.

However, for all the concrete examples that we shall meet in practice, it is always question (\mathbf{Q}') that will be answered.

As will be explained in this thesis, a formal consequence of the above is that polygraphic homology is *not* invariant under Thomason equivalence. This means that there exists at least one Thomason equivalence $f: C \to D$ such that the induced map

 $\mathbb{H}^{\mathrm{pol}}(C) \to \mathbb{H}^{\mathrm{pol}}(D)$

is *not* an isomorphism. In other words, if we think of ω Cat as a model of homotopy types (via the localization by the Thomason equivalences), then the polygraphic homology is *not* a well-defined invariant. Another point of view would be to consider that the polygraphic homology is an intrinsic invariant of ω -categories (and not up to Thomason equivalence) and in that way is finer than singular homology. This is not the point of view adopted here, and the reason will be motivated at the end of this introduction. The slogan to retain is:

Polygraphic homology is a way of computing singular homology groups of a homologically coherent ω -category.

The point is that given a *free* ω -category P (which is thus its own polygraphic resolution), the chain complex $\lambda(P)$ is much "smaller" than the chain complex associated to the nerve of P and hence the polygraphic homology groups of P are much easier to compute than its singular homology groups. The situation is comparable to using cellular homology for computing singular homology of a CW-complex. The difference is that in this last case, such a thing is always possible while in the case of ω -categories, one must ensure that the (free) ω -category is homologically coherent.

Finding homologically coherent ω -categories. One of the main results presented in this dissertation is:

Every (small) category C is homologically coherent.

In order for this result to make sense, one has to consider categories as ω -categories with only unit cells above dimension 1. Beware that this does not make the result trivial because given a polygraphic resolution $P \rightarrow C$ of a small category C, the ω -category P need *not* have only unit cells above dimension 1.

As such, this result is only a small generalization of Lafont and Métayer's result concerning monoids (although this new result, even restricted to monoids, is more precise because it means that the *canonical comparison map* is an isomorphism). But the true novelty lies in the proof which is more conceptual that the one of Lafont and Métayer. It requires the development of several new concepts and results which in the end combine together smoothly to yield the desired result. This dissertation has been written so that all the elements needed to prove this result are spread over several chapters; a

more condensed version of it is the object of the article [Gue21]. Among the new notions developed along the way, that of discrete Conduché ω -functor is probably the most significant. An ω -functor $f : C \to D$ is a *discrete Conduché* ω -functor when for every cell x of C, if f(x) can be written as

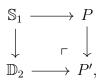
$$f(x) = y' * y'',$$

then there exists a unique pair (x', x'') of cells of C that are k-composable and such that

$$f(x') = y', f(x'') = y'' \text{ and } x = x' \underset{k}{*} x''.$$

The main result that we prove concerning discrete Conduché ω -functors is that for a discrete Conduché ω -functor $f : C \to D$, if the ω -category D is free, then C is also free. The proof of this result is long and tedious, though conceptually not extremely hard, and first appears in the paper [Gue20], which is dedicated to it.

After having settled the case of (1-)categories, it is natural to move on to 2-categories. Contrary to the case of (1-)categories, not all 2-categories are homologically coherent and the situation seems to be much harder to understand. As a simplification, one can focus on 2-categories which are free (as ω -categories). This is what is done in this dissertation. With this extra hypothesis, the problem of characterization of homologically coherent free 2-categories may be reduced to the following question: given a cocartesian square of the form



where P is a free 2-category, when is it *homotopy cocartesian* with respect to the Thomason equivalences? As a consequence, a substantial part of the work presented here consists in developing tools to detect homotopy cocartesian squares of 2-categories with respect to the Thomason equivalences. While it appears that these tools do not allow to completely answer the above question, they still make it possible to detect such homotopy cocartesian squares in many concrete situations. In fact, a whole section of the thesis is dedicated to giving examples of (free) 2-categories and computing the homotopy type of their nerve using these tools. Among all these examples, a particular class of well-behaved 2-categories, which I have coined "bubble-free 2-category, let us call *bubble* a non-trivial 2-cell whose source and target are units on a 0-cell (necessarily the same). A *bubble-free 2-category* is then nothing but a 2-category that has no bubbles. The archetypal example of a 2-category that is *not* bubble-free is the 2-category. As already said, this 2-category is not homologically coherent and this does not seem to

be a coincidence. It is indeed remarkable that of all the many examples of 2-categories studied in this work, the only ones that are not homologically coherent are exactly the ones that are *not* bubble-free. This leads to the conjecture below, which stands as a conclusion of the thesis.

(Conjecture) A free 2-category is homologically coherent if and only if it is bubble-free.

The big picture. Let us end this introduction with another point of view on the comparison of singular and polygraphic homologies. This point of view is highly conjectural and is not addressed at all in the rest of the dissertation. It should be thought of as a guideline for future work.

In the same way that (strict) 2-categories are particular cases of bicategories, strict ω -categories are in fact particular cases of what are usually called *weak* ω -categories. Such mathematical objects have been defined, for example, by Batanin using globular operads [Bat98] or by Maltsiniotis following ideas of Grothendieck [Mal10]. Similarly to the fact that the theory of quasi-categories (which is a homotopical model for the theory of weak ω -categories whose cells are invertible above dimension 1) may be expressed using the same language as the theory of usual categories, it is generally believed that all "intrinsic" notions (in a precise sense to be defined) of the theory of strict ω -categories have weak counterparts. For example, it is believed that there should be a folk model structure on the category of weak ω -categories and that there should be a good notion of free weak ω -category. In fact, this last notion should be defined as weak ω -categories that are recursively obtained from the empty ω -category by freely adjoining cells, which is the formal analogue of the strict version but in the weak context. The important point here is that a free strict ω -category is *never* free as a weak ω -category (except for the empty ω -category). Moreover, there are good candidates for the polygraphic homology of weak ω -categories obtained by mimicking the definition in the strict case. But in general the polygraphic homology of a strict ω -category need not be the same as its "weak polygraphic homology". Indeed, since free strict ω -categories are not free as weak ω -categories, taking a "weak polygraphic resolution" of a strict ω -category is not the same as taking a polygraphic resolution. In fact, when trying to compute the weak polygraphic homology of B, it would seem that it gives the homology groups of a $K(\mathbb{Z}, 2)$, which is what we would have expected of its polygraphic homology in the first place. From this observation, it is tempting to make the following conjecture:

The weak polygraphic homology of a strict ω -category coincides with its singular homology.

In other words, we conjecture that the fact that polygraphic and singular homologies of strict ω -categories do not coincide is a defect due to working in too narrow a setting. The "good" definition of polygraphic homology ought to be the weak one.

We can go even further and conjecture the same thing for weak ω -categories. In order to do so, we need a definition of singular homology for weak ω -categories. This is conjecturally done as follows. To every weak ω -category C, one can associate a weak ω -groupoid L(C) by formally inverting all the cells of C. Then, if we believe in Grothendieck's conjecture (see [Gro83] and [Mal10, Section 2]), the category of weak ω -groupoids equipped with the weak equivalences of weak ω -groupoids (see [Mal10, Paragraph 2.2]) is a model for the homotopy theory of spaces. In particular, every weak ω -groupoid has homology groups and we can define the singular homology groups of a weak ω -category C as the homology groups of L(C).

Organization of the thesis. In the first chapter, we review some aspects of the theory of ω -categories. In particular, we study with great care free ω -categories, which are at the heart of the present work. It is the only chapter of the thesis that does not contain any reference to homotopy theory whatsoever. It is also there that we introduce the notion of discrete Conduché ω -functor and study their relation with free ω -categories. The culminating point of the chapter is Theorem 1.6.18, which states that given a discrete Conduché ω -functor $F : C \to D$, if D is free, then so is C. The proof of this theorem is long and technical and is broke down into several distinct parts.

The second chapter is devoted to recalling some tools of homotopical algebra. More precisely, basic aspects of the theory of homotopy colimits using the formalism of Grothendieck's derivators are quickly presented. Note that this chapter does *not* contain any original result and can be skipped at first reading. It is only intended to give the reader a summary of useful results on homotopy colimits that are used in the rest of the dissertation.

In the third chapter, we delve into the homotopy theory of ω -categories. It is there that we define the different notions of weak equivalences for ω -categories and compare them. The two most significant new results to be found in this chapter are probably Proposition 3.6.2, which states that every equivalence of ω -categories is a Thomason equivalence, and Theorem 3.7.4, which states that equivalences of ω -categories satisfy a property reminiscent of Quillen's Theorem A [Qui73, Theorem A] and its ω -categorical generalization by Ara and Maltsiniotis [AM18, AM20c].

In the fourth chapter, we define the polygraphic and singular homologies of ω -categories and properly formulate the problem of their comparison. Up to Section 4.3 included, all the results were known prior to this thesis (at least in the folklore), but starting from Section 4.4 all the results are original. Three fundamental results of this chapter are: Theorem 4.4.5, which states that singular homology is obtained as a derived functor of an abelianization function, Proposition 4.5.10, which gives an abstract criterion to detect homologically coherent ω -categories, and Proposition 4.6.23, which states that low-dimensional singular and polygraphic homology groups always coincide.

The fifth chapter is mainly geared towards the fundamental Theorem 5.3.14, which states that every category is homologically coherent. To prove this theorem, we first

focus on a particular class of ω -categories, which we call *contractible* ω -categories, and show that every contractible ω -category is homologically coherent (Proposition 5.1.2).

Finally, the sixth and last chapter of the thesis revolves around the homology of free 2-categories. The goal pursued is to try to understand which free 2-categories are homologically coherent. In order to do so, we give a criterion to detect homotopy co-cartesian square with respect to Thomason equivalences (Proposition 6.3.14) based on the homotopy theory of bisimplicial sets. Then, we apply this criterion and some other *ad hoc* techniques to compute many examples of homotopy type of free 2-categories. The conclusion of the chapter is Conjecture 6.6.5, which states that a free 2-category is homologically coherent if and only if it is bubble-free.

INTRODUCTION (FRANÇAIS)

INTRODUCTION (FRANÇAIS)

Cette thèse a pour cadre général la *théorie de l'homotopie des* ω -catégories strictes, et, comme son titre le suggère, ce sont les aspects homologiques de cette théorie qui sont traîtés. Le but est d'étudier et de comparer deux invariants homologiques différents associés aux ω -catégories strictes; c'est-à-dire, deux foncteurs différents

 $\operatorname{Str}\omega\operatorname{Cat}\to\operatorname{ho}(\operatorname{Ch}_{\geq 0})$

de la catégorie des ω -catégories strictes vers la catégorie homotopique des complexes de chaînes en degré positif (i.e. la localisation de la catégorie des complexes de chaînes en degré positif relativement aux quasi-isomorphismes).

Avant d'entrer dans le vif du sujet, précisons sans plus tarder qu'à l'unique exception de la toute fin de cette introduction, toutes les ω -catégories que nous considérerons sont des ω -catégories strictes. C'est pourquoi nous omettrons l'adjectif « strict » et parlerons simplement de ω -catégorie plutôt que de ω -catégorie stricte. De même, nous noterons ω Cat plutôt que Str ω Cat la catégorie des ω -catégories (strictes).

Le contexte : les ω -categories en tant qu'espaces. L'étude de la théorie de l'homotopie des ω -catégories commence avec le foncteur nerf introduit par Street [Str87]

$$N_{\omega}: \omega \mathbf{Cat} \to \widehat{\Delta}$$

qui associe à toute ω -catégorie C un ensemble simplicial $N_{\omega}(C)$, appelé le *nerf de* C, généralisant le nerf usuel des (petites) catégories. En utilisant ce foncteur, il est possible de transférer la théorie de l'homotopie des ensembles simpliciaux aux ω -catégories, comme cela est fait par exemple dans les articles [AM14, AM18, Gag18, Ara19, AM20c, AM20a]. Suivant la terminologie de ces derniers, un morphisme $f : C \to D$ de la catégorie ω Cat est une équivalence de Thomason si $N_{\omega}(f)$ est une équivalence faible de Kan–Quillen d'ensembles simpliciaux. Par définition, le foncteur nerf induit un foncteur au niveau des catégories homotopiques

$$\overline{N_{\omega}}: \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) \to \mathrm{ho}(\widehat{\Delta}),$$

où $ho(\omega Cat^{Th})$ est la localisation de ωCat relativement aux équivalences de Thomason et $ho(\widehat{\Delta})$ est la localisation de $\widehat{\Delta}$ relativement aux équivalences faibles de Kan–Quillen

d'ensembles simpliciaux. Comme l'a démontré Gagna [Gag18], ce dernier foncteur est en fait une équivalence de catégories. Autrement dit, la théorie de l'homotopie des ω -catégories induite par les équivalences de Thomason est la même que la théorie de l'homotopie des espaces. Le résultat de Gagna est une généralisation du résultat analogue pour le nerf usuel des petites catégories, attribué à Quillen dans [III72]. Dans le cas des petites catégories, Thomason a même démontré l'existence d'une structure de catégorie de modèles dont les équivalences faibles sont celles induites par le nerf [Tho80]. Le résultat analogue pour ω Cat est toujours une conjecture [AM14].

Deux homologies pour les ω -catégories. Armé du foncteur nerf de Street, il est naturel de définir le *k-ème groupe d'homologie d'une* ω -catégorie *C* comme étant le *k*-ème groupe d'homologie du nerf de *C*. À la lumière du résultat de Gagna, ces groupes d'homologies sont simplement les groupes d'homologie des espaces vus sous un autre angle. Afin d'éviter de potentielles confusions à venir, nous appellerons désormais ces groupes d'homologie les groupes d'homologie singulière de *C* et nous utiliserons la notation $H_k^{\text{Sing}}(C)$.

D'autres groupes d'homologie pour les ω -catégories ont aussi été définis par Métayer dans [Mét03]. La définition de ceux-ci repose sur la notion de ω -catégorie libre sur un polygraphe (aussi connue sous le nom de computade), c'est-à-dire de ω -catégorie obtenue de manière récursive à partir de la catégorie vide en attachant librement des cellules. Désormais, nous dirons simplement ω -catégorie libre. Il a été observé par Métayer que toute ω -catégorie C admet une résolution polygraphique, c'est-à-dire qu'il existe une ω -catégorie libre P et un morphisme de ω Cat

$$f: P \to C,$$

satisfaisant des propriétés analogues à celles des fibrations triviales d'espaces topologiques (ou d'ensemples simpliciaux). De plus, toute ω -catégorie libre P peut être « abélianisée » en un complexe de chaînes $\lambda(P)$ et il a été démontré par Métayer que pour deux résolutions polygraphiques $P \to C$ et $P' \to C$ d'une même ω -catégorie, les complexes de chaînes $\lambda(P)$ et $\lambda(P')$ sont quasi-isomorphes. Ainsi, on peut définir le k-ème groupe d'homologie polygraphique de C, noté $H_k^{\text{pol}}(C)$, comme étant le k-ème groupe d'homologie de $\lambda(P)$ pour n'importe quelle résolution polygraphique $P \to C$.

Ces considérations invitent à se poser la question suivante :

A-t-on
$$H^{\text{pol}}_{\bullet}(C) \simeq H^{\text{Sing}}_{\bullet}(C)$$
 pour toute ω -catégorie C?

Une première réponse partielle a été donnée par Lafont et Métayer [LM09] : pour tout monoïde M (vu comme une ω -catégorie à un seul objet et dont toutes les cellules de dimension supérieure à 1 sont des unités), on a $H^{\text{pol}}_{\bullet}(M) \simeq H^{\text{Sing}}_{\bullet}(M)$. Mentionnons au passage que l'homologie polygraphique a été conçue originellement pour étudier l'homologie des monoïdes et fait partie d'un programme dont le but est de généraliser en dimension supérieure les travaux de Squier sur la théorie de la réécriture des monoïdes [Gui06, Laf07, GM09, GM18]. Malgré le résultat de Lafont et Métayer, la réponse générale à la question précédente est *non*. Un contre-exemple a été découvert par Maltsiniotis et Ara. Soit *B* le monoïde commutatif (\mathbb{N} , +), vu comme une 2-catégorie avec une seule 0-cellule et pas de 1-cellule non-triviale. Cette 2-catégorie est libre (en tant que ω -catégorie) et un calcul rapide montre que :

$$H_k^{\rm pol}(B) = \begin{cases} \mathbb{Z} & \text{ pour } k = 0, 2\\ 0 & \text{ sinon.} \end{cases}$$

D'autre part, Ara a démontré [Ara19, Theorem 4.9 et Example 4.10] que le nerf de B est un $K(\mathbb{Z}, 2)$, qui a donc des groupes d'homologie non triviaux en toute dimension paire.

La question devient ainsi :

(Q) Quelles sont les
$$\omega$$
-catégories C pour lesquelles $H^{\text{pol}}_{\bullet}(C) \simeq H^{\text{Sing}}_{\bullet}(C)$?

C'est précisément à cette question que tente de répondre cette thèse. Néanmoins, le lecteur trouvera dans ce document plusieurs notions nouvelles et résultats qui, bien qu'originellement motivés par la question ci-dessus, sont intrinsèquement intéressants pour la théorie des ω -catégories et dont la portée dépasse les considérations homologiques précédentes.

Une autre formulation du problème. Un des accomplissements du travail présenté ici est l'établissement d'un cadre conceptuel qui permet une reformulation plus abstraite et plus satisfaisante de la question de comparaison de l'homologie polygraphique et de l'homologie singulière des ω -catégories.

Pour cela, rappelons tout d'abord qu'une variation de l'équivalence de Dold-Kan (voir par exemple [Bou90]) permet d'affirmer que la catégorie des objets en groupes abéliens dans la catégorie ω Cat est équivalente à la catégorie des complexes de chaînes en degré positif

$$\mathbf{Ab}(\omega \mathbf{Cat}) \simeq \mathbf{Ch}_{\geq 0}.$$

Ainsi, on a un foncteur d'oubli $Ch_{\geq 0} \simeq Ab(\omega Cat) \rightarrow \omega Cat$, qui admet un adjoint à gauche

$$\lambda: \omega \mathbf{Cat} \to \mathbf{Ch}_{>0}.$$

En outre, pour une ω -catégorie *libre* C, le complexe de chaînes $\lambda(C)$ est exactement celui obtenu par le processus d'« abélianisation » que Métayer utilise dans sa définition d'homologie polygraphique.

Par ailleurs, la catégorie ω Cat peut être munie d'une structure de catégorie de modèles, communément appelée *la structure de catégorie de modèles folk* [LMW10], dont les équivalences faibles sont les *équivalences de* ω -catégories (notion généralisant celle d'équivalence de catégories) et dont les objets cofibrants sont les ω -catégories libres [Mét08]. Les résolutions polygraphiques ne sont alors rien d'autre que des remplacements cofibrants pour cette structure de catégorie de modèles. Comme la définition des groupes d'homologie polygraphique le laissait deviner, le foncteur λ est Quillen à gauche relativement à cette structure de catégorie de modèles. En particulier, ce foncteur admet un foncteur dérivé à gauche

$$\mathbb{L}\lambda^{\text{folk}}: \text{ho}(\omega \mathbf{Cat}^{\text{folk}}) \to \text{ho}(\mathbf{Ch}_{\geq 0})$$

et on a tautologiquement $H_k^{\text{pol}}(C) = H_k(\mathbb{L}\lambda^{\text{folk}}(C))$ pour toute ω -catégorie C et pour tout $k \ge 0$. Désormais, on posera même

$$\mathbb{H}^{\mathrm{pol}}(C) := \mathbb{L}\lambda^{\mathrm{folk}}(C).$$

Je précise que cette façon de comprendre l'homologie polygraphique comme foncteur dérivé à gauche fait partie du folklore depuis un certain temps et je ne prétends à aucune originalité concernant ce point précis.

D'autre part, le foncteur λ est aussi dérivable à gauche quand ω Cat est munie des équivalences de Thomason, ce qui permet d'obtenir un foncteur dérivé à gauche

$$\mathbb{L}\lambda^{\mathrm{Th}}$$
: ho($\omega \mathbf{Cat}^{\mathrm{Th}}$) \rightarrow ho($\mathbf{Ch}_{>0}$).

En outre, ce foncteur est tel que $H_k^{\text{Sing}}(C) = H_k(\mathbb{L}\lambda^{\text{Th}}(C))$ pour toute ω -catégorie C et pour tout $k \ge 0$. Contrairement au cas « folk », ce résultat est complètement nouveau et apparaît pour la première fois dans ce manuscrit (à ma connaissance, du moins). Notons également que, puisque l'existence d'une structure de catégorie de modèles « à la Thomason » sur ω Cat est toujours conjecturale, les outils habituels de la théorie de Quillen des catégories de modèles sont inutilisables pour démontrer que λ est dérivable à gauche. La difficulté fut de trouver un moyen de contourner ce problème.

Désormais, on posera

$$\mathbb{H}^{\mathrm{Sing}}(C) := \mathbb{L}\lambda^{\mathrm{Th}}(C).$$

Finalement, on peut montrer que toute équivalence de ω -catégories est une équivalence de Thomason. Ainsi, le foncteur identité de ω Cat induit formellement un foncteur \mathcal{J} au niveau des catégories homotopiques

$$\mathcal{J}: \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{folk}}) \to \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}),$$

et on obtient donc un triangle

$$\begin{array}{ccc} \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{folk}}) & \xrightarrow{\mathcal{J}} & \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{Th}}) \\ & & & \downarrow_{\mathbb{L}\lambda^{\operatorname{folk}}} & & \downarrow_{\mathbb{L}\lambda^{\operatorname{Th}}} \\ & & & \operatorname{ho}(\mathbf{Ch}_{\geq 0}). \end{array}$$

Ce triangle n'est *pas* commutatif (même à un isomorphisme près), car cela impliquerait que les groupes d'homologie singulière et les groupes d'homologie polygraphique coïncident pour toute ω -catégorie. Néanmoins, puisque les foncteurs $\mathbb{L}\lambda^{\text{folk}}$ et $\mathbb{L}\lambda^{\text{Th}}$ sont tous les deux des foncteurs dérivés à gauche du même foncteur, l'existence d'une transformation naturelle $\pi : \mathbb{L}\lambda^{\text{Th}} \circ \mathcal{J} \Rightarrow \mathbb{L}\lambda^{\text{folk}}$ découle formellement par propriété universelle. De plus, \mathcal{J} étant l'identité sur les objets, cette transformation naturelle fournit pour toute ω -catégorie C un morphisme

$$\pi_C : \mathbb{H}^{\mathrm{Sing}}(C) \to \mathbb{H}^{\mathrm{pol}}(C),$$

que nous appellerons le *morphisme de comparaison canonique*. Une ω -catégorie C est qualifiée d'*homologiquement cohérente* si π_C est un isomorphisme (ce qui signifie exactement que le morphisme induit $H_k^{\text{Sing}}(C) \to H_k^{\text{pol}}(C)$ est un isomorphisme pour tout $k \ge 0$). La question devient alors :

(Q') Quelles ω -catégories sont homologiquement cohérentes ?

Notons au passage que la question (\mathbf{Q}') est théoriquement plus précise que la question (\mathbf{Q}) . Cependant, dans tous les exemples concrets que nous rencontrerons, c'est toujours à la question (\mathbf{Q}') que nous répondrons.

Comme il sera expliqué dans la thèse, de cette reformulation en terme de foncteurs dérivés, il est également possible de déduire formellement que l'homologie polygraphique n'est *pas* invariante relativement aux équivalences de Thomason. Cela signifie qu'il existe au moins une équivalence de Thomason $f: C \to D$ telle que le morphisme induit en homologie polygraphique

$$\mathbb{H}^{\mathrm{pol}}(C) \to \mathbb{H}^{\mathrm{pol}}(D)$$

n'est *pas* un isomorphisme. En d'autres termes, si nous voyons les ω -catégories comme modèles des types d'homotopie (via la localisation de ω Cat relativement aux équivalences de Thomason), alors l'homologie polygraphique n'est *pas* un invariant bien défini. Un autre point de vue possible serait de considérer que l'homologie polygraphique est un invariant intrinsèque des ω -catégories (et non pas à équivalence de Thomason près) et, de cette façon, est un invariant plus fin que l'homologie singulière. Ce n'est pas le point de vue adopté dans cette thèse et ce choix sera motivé à la fin de cette introduction. Le slogan à retenir est :

L'homologie polygraphique est un moyen de calculer les groupes d'homologie singulière des ω -catégories homologiquement cohérentes.

L'idée étant que pour une ω -catégorie libre P (qui est donc sa propre résolution polygraphique), le complexe de chaînes $\lambda(P)$ est beaucoup moins « gros » que le complexe de chaînes associé au nerf de P, et ainsi les groupes d'homologie polygraphique de P sont beaucoup plus faciles à calculer que les groupes d'homologie singulière. La situation est comparable à l'utilisation de l'homologie cellulaire afin de calculer les groupes d'homologie singulière d'un CW-complexe. La différence étant que dans ce dernier cas il est toujours possible de procéder ainsi, alors que dans le cas des ω -catégories, on doit d'abord s'assurer que la ω -catégorie (libre) en question est bien homologiquement cohérente.

Détecter les ω -catégories qui sont homologiquement cohérentes. Un des principaux résultats de cette thèse est le suivant :

Toute (petite) catégorie est homologiquement cohérente.

Afin de donner du sens à cette assertion, il faut considérer les catégories comme des ω -catégories dont les cellules au delà de la dimension 1 sont des unités. Le résultat cidessus n'est pas pour autant trivial car pour une résolution polygraphique $P \rightarrow C$ d'une petite catégorie C, la ω -catégorie P, elle, n'a pas forcément que des cellules unités au delà de la dimension 1.

En tant que tel, ce résultat est seulement une petite généralisation du résultat de Lafont et Métayer sur les monoïdes (bien qu'il soit plus précis, même restreint aux monoïdes, car il dit que c'est le morphisme de comparaison canonique qui est un isomorphisme). Mais la véritable nouveauté du résultat en est sa démonstration qui est plus conceptuelle que celle de Lafont et Métayer. Elle repose sur l'introduction de nouvelles notions et le développement de nouveaux résultats; le tout s'assemblant élégamment pour finalement produire le résultat voulu. Cette thèse a été écrite de telle façon que tous les élements nécessaires à la démonstration du résultat précédent sont répartis sur plusieurs chapitres; une version plus condensée de celle-ci se trouve dans l'article [Gue21]. Parmi les nouvelles notions développées, la plus significative est probablement celle de ω -foncteur de Conduché discret. Un ω -foncteur $f : C \to D$ est un ω -foncteur de Conduché discret quand pour toute cellule x de C, si f(x) peut être décomposé en

$$f(x) = y' \underset{k}{*} y'',$$

alors il existe une unique paire (x', x'') de cellules de C qui sont k-composables et telles que

$$f(x') = y', f(x'') = y'' \text{ and } x = x' * x''.$$

Le résultat principal démontré concernant cette notion est que pour tout ω -foncteur de Conduché discret $f: C \to D$, si D est libre alors C est aussi libre. La démonstration est longue et fastidieuse, bien que relativement facile d'un point de vue conceptuel, et apparaît pour la première fois dans le papier [Gue20], qui lui est dédié.

Une fois le cas de l'homologie des (1-)catégories complètement résolu, il est naturel de s'intéresser aux 2-catégories. Contrairement au cas des (1-)catégories, les 2-catégories ne sont pas toutes homologiquement cohérentes et la situation est beaucoup plus

compliquée. En premier lieu, on peut se restreindre aux 2-catégories qui sont libres (en tant que ω -catégories) et c'est ce qui est fait dans cette thèse. Avec cette hypothèse simplificatrice, le problème de caractérisation des 2-catégories libres homologiquement cohérentes peut être réduit à la question suivante : soit un carré cocartésien de la forme

$$\begin{array}{c} \mathbb{S}_1 \longrightarrow P \\ \downarrow & & \downarrow \\ \mathbb{D}_2 \longrightarrow P', \end{array}$$

où P est une 2-catégorie libre, quand est-il homotopiquement cocartésien relativement aux équivalences de Thomason? En conséquence, une partie substantielle du travail présenté ici consiste à développer des outils permettant de reconnaître les carrés homotopiquement cocartésiens de 2-catégories relativement aux équivalences de Thomason. Bien que les outils qui seront présentés ne permettent pas de répondre entièrement à la question ci-dessus, ils permettent tout de même de détecter de tels carrés cocartésiens dans beaucoup de situations concrètes. Il y a même une section entière de la thèse qui consiste uniquement en une liste d'exemples détaillés de calculs du type d'homotopie de 2-catégories libres en utilisant ces outils. De ces exemples se dégage très clairement une classe particulière de 2-catégories, que j'ai nommées les « 2-catégories sans bulles » et qui sont caractérisées comme suit. Pour une 2-catégorie, appelons bulle une 2-cellule non-triviale dont la source et le but sont des unités sur une 0-cellule (nécessairement la même). Une 2-catégorie sans bulles est simplement une 2-catégorie qui n'a aucune bulle. L'archétype de la 2-catégorie qui n'est pas sans bulles est la 2-catégorie B que nous avons déjà rencontrée plus haut (c'est-à-dire le monoïde commutatif $(\mathbb{N}, +)$ vu comme une 2-catégorie). Comme dit précédemment, cette 2-catégorie n'est pas homologiquement cohérente et cela ne semble pas être une coïncidence. Il est tout à fait remarquable que de toutes les 2-catégories étudiées dans cette thèse, les seules qui ne sont pas homologiquement cohérentes sont exactement celles qui ne sont pas sans bulles. Cela conduit à la conjecture ci-dessous, qui est le point d'orgue de la thèse.

(**Conjecture**) Une 2-catégorie libre est homologie cohérente si et seulement si elle est sans bulles.

Une vue d'ensemble. Terminons cette introduction par un autre point de vue sur la comparaison des homologies polygraphique et singulière. Précisons immédiatement que ce point de vue est hautement conjectural et n'est pas du tout abordé dans le reste de la thèse. Il s'agit plus d'un guide pour des travaux futurs qu'autre chose.

De la même façon que les 2-catégories (strictes) sont des cas particuliers de bicatégories, les ω -catégories strictes sont en réalité des cas particuliers de ce qui est communément appelé des ω -catégories faibles. Ces objets mathématiques ont été définis, par exemple, par Batanin en utilisant le formalisme des opérades globulaires [Bat98] ou

par Maltsiniotis en suivant des idées de Grothendieck [Mal10]. Tout comme la théorie des quasi-catégories (qui est un modèle homotopique pour la théorie des ω -catégories faibles dont les cellules sont toutes inversibles au delà de la dimension 1) s'exprime avec le même langage que la théorie des catégories usuelle, il est attendu que toutes les notions « intrinsèques » (dans un sens précis à définir) de la théorie des ω -catégories strictes ont des analogues faibles. Par exemple, il est attendu qu'il y ait une structure de catégorie de modèles folk sur la catégorie des ω -catégories faibles et qu'il y ait une bonne notion de ω -catégorie faible libre. En fait, ces dernières seraient certainement définies comme les ω -catégories faibles qui sont récursivement obtenues à partir de la ω -catégorie vide en attachant librement des cellules, ce qui est l'analogue formelle du cas strict. Le point clé ici est qu'une ω -catégorie stricte libre n'est jamais libre en tant que ω -catégorie faible (excepté la ω -catégorie vide). Par ailleurs, il existe de bons candidats pour l'homologie polygraphique des ω -catégories faibles qui sont obtenus par mimétisme de la définition du cas strict. Mais il n'y aucune raison en général que l'homologie polygraphique d'une ω -catégorie stricte soit la même que son « homologie polygraphique faible ». En effet, puisque les ω -catégories strictes libres ne sont pas libres en tant que ω -catégories faibles, prendre une « résolution polygraphique faible » d'une ω -catégorie libre ne revient pas à prendre une résolution polygraphique. De fait, lorsqu'on essaye de calculer l'homologie polygraphique faible de B, il semblerait que cela donne les groupes d'homologie d'un $K(\mathbb{Z}, 2)$, ce qui aurait été attendu de l'homologie polygraphique au départ. De cette observation, il est tentant de faire la conjecture suivante :

L'homologie polygraphique faible d'une ω -categorie stricte coïncide avec son homologie singulière.

En d'autres termes, nous conjecturons que le fait que l'homologie polygraphique et l'homologie singulière d'une ω -catégorie stricte ne coïncident pas est un défaut dû à un cadre de travail trop étroit. La « bonne » définition de l'homologie polygraphique devrait être la faible.

Nous pourrions même aller plus loin et conjecturer la même chose pour les ω -catégories faibles. Pour cela, il est nécessaire de disposer d'une définition de l'homologie singulière des ω -catégories faibles. Conjecturellement, on procède de la manière suivante. À toute ω -catégorie faible C, on peut associer un ω -groupoïde faible L(C)en inversant formellement toutes les cellules de C. Puis, si on en croit la conjecture de Grothendieck (voir [Gro83] et [Mal10, Section 2]), la catégorie des ω -groupoïdes faibles munie des équivalences de ω -groupoïdes faibles (voir [Mal10, Paragraph 2.2]) est un modèle de la théorie homotopique des espaces. En particulier, chaque ω -groupoïde a des groupes d'homologie et on peut définir les groupes d'homologie singulière d'une ω -catégorie faible C comme les groupes d'homologie de L(C). **Organisation de la thèse.** Dans le premier chapitre, nous passerons en revue quelques aspects de la théorie des ω -catégories. En particulier, nous étudierons avec grand soin les ω -catégories libres, qui sont au cœur de cette thèse. C'est le seul chapitre de la thèse qui ne contient aucune référence à la théorie de l'homotopie. C'est également dans ce chapitre que nous introduirons la notion de ω -foncteur de Conduché discret et que nous étudierons leur relation avec les ω -catégories libres. Le point culminant du chapitre étant le théorème 1.6.18, qui dit que pour un ω -foncteur de Conduché discret $F : C \to D$, si D est libre, alors C l'est aussi. La démonstration de ce théorème est longue et technique et est décomposée en plusieurs parties distinctes.

Le second chapitre a pour but de rappeler quelques outils d'algèbre homotopique. En particulier, les aspects élémentaires de la théorie des colimites homotopiques en utilisant le formalisme de Grothendieck des dérivateurs y sont rapidement présentés. Notons au passage que ce chapitre ne contient *aucun* résultat original et peut être omis en première lecture. Son unique raison d'être est de donner au lecteur un catalogue de résultats concernant les colimites homotopiques qui seront utilisés par la suite.

Dans le troisième chapitre, nous aborderons enfin la théorie de l'homotopie des ω -catégories. C'est là que seront définies et comparées entre elles les différentes notions d'équivalences faibles pour les ω -catégories. Les deux résultats les plus significatifs de ce chapitre sont probablement la proposition 3.6.2, qui dit que toute équivalence de ω -catégorie est une équivalence de Thomason, et le théorème 3.7.4, qui dit que les équivalences de ω -catégories satisfont une propriété réminiscente du théorème A de Quillen [Qui73, Theorem A] et sa généralisation ω -catégorique par Ara et Maltsiniotis [AM18, AM20c].

Dans le quatrième chapitre, nous définirons les homologies polygraphique et singulière des ω -catégories et fomulerons précisément le problème de leur comparaison. Jusqu'à la section 4.3 incluse, tous les résultats étaient connus avant cette thèse (au moins dans le folklore), mais à partir de la section 4.4 tous les résultats sont orignaux. Trois résultats fondamentaux de ce chapitre sont les suivants : le théorème 4.4.5, qui dit que l'homologie singulière s'obtient comme le foncteur dérivé d'un foncteur d'abélianisation, la proposition 4.5.10, qui donne un critère abstrait pour détecter les ω -catégories homologiquement cohérentes, et la proposition 4.6.23, qui dit que les groupes d'homologies polygraphique et singulière coïncident toujours en basse dimension.

Le cinquième chapitre a pour but de démontrer le théorème fondamental 5.3.14, qui dit que toute catégorie est homologiquement cohérente. Pour cela, nous nous intéresserons en premier lieu à une classe particulière de ω -catégories, dites *contractiles*, et nous montrerons que toute ω -catégorie contractile est homologiquement cohérente (Proposition 5.1.2).

Enfin, le sixième et dernier chapitre de la thèse s'intéresse à l'homologie des 2-catégories libres. Le but est d'essayer de comprendre quelles sont les 2-catégories libres qui sont homologiquement cohérentes. Pour cela, un critère pour détecter les carrés homotopiquement cocartésiens relativement aux équivalences de Thomason y est donné (Proposition 6.3.14). Ce critère est fondé sur la théorie de l'homotopie des ensembles bisimpliciaux. Ensuite, nous appliquerons ce critère ainsi que d'autres techniques *ad hoc* au calcul du type d'homotopie d'un grand nombre de 2-catégories libres. La conclusion du chapitre est la conjecture 6.6.5, qui énonce qu'une 2-catégorie libre est homologiquement cohérente si et seulement si elle est sans bulles.

CHAPTER 1

YOGA OF ω -CATEGORIES

1.1 ω -GRAPHS, ω -MAGMAS AND ω -CATEGORIES

1.1.1. An ω -graph X consists of an infinite sequence of sets $(X_n)_{n \in \mathbb{N}}$ together with maps

$$X_n \xleftarrow{\mathrm{t}}_{\mathrm{s}} X_{n+1}$$

for every $n \in \mathbb{N}$, subject to the *globular identities*:

$$\begin{cases} s \circ s = s \circ t, \\ t \circ t = t \circ s. \end{cases}$$

Elements of X_n are called *n*-cells or *n*-arrows or cells of dimension *n*. For n = 0, elements of X_0 are also called *objects*. For *x* an *n*-cell with n > 0, s(x) is the source of *x* and t(x) is the target of *x*. We use the notation

$$x: a \to b$$

to say that a is the source of x and b is the target of x.

More generally, for $0 \le k < n$, we define maps $s_k : X_n \to X_k$ and $t_k : X_n \to X_k$ as

$$\mathbf{s}_k = \underbrace{\mathbf{s} \circ \cdots \circ \mathbf{s}}_{n-k \text{ times}}$$

and

$$\mathbf{t}_k = \underbrace{\mathbf{t} \circ \cdots \circ \mathbf{t}}_{n-k \text{ times}}.$$

For an *n*-cell x, the k-cells $s_k(x)$ and $t_k(x)$ are respectively the k-source and the k-target of x.

Two *n*-cells x and y are *parallel* if

n = 0

or

$$n > 0$$
 and $s(x) = s(y)$ and $t(x) = t(y)$.

Let $0 \le k < n$. Two *n*-cells *x* and *y* are *k*-composable if

$$\mathbf{s}_k(x) = \mathbf{t}_k(y).$$

Note that the expression "x and y are k-composable" is *not* symmetric in x and y and we *should* rather speak of a "k-composable pair (x, y)", although we won't always do it. The set of pairs of k-composable n-cells is denoted by $X_n \underset{X_k}{\times} X_n$, and is characterized as the following fibred product

$$\begin{array}{cccc} X_n \times X_n & \longrightarrow & X_n \\ & & \downarrow & & \downarrow^{\mathbf{t}_k} \\ & & & \downarrow & & \downarrow^{\mathbf{t}_k} \\ & X_n & \xrightarrow{\mathbf{s}_k} & X_k. \end{array}$$

A morphism of ω -graphs $f : X \to Y$ is a sequence $(f_n : X_n \to Y_n)_{n \in \mathbb{N}}$ of maps that is compatible with source and target, i.e. for every *n*-cell x of X with n > 0, we have

$$f_{n-1}(s(x)) = s(f_n(x))$$
 and $f_{n-1}(t(x)) = t(f_n(x))$.

For an *n*-cell x of X, we often write f(x) instead of $f_n(x)$.

The category of ω -graphs and morphisms of ω -graphs is denoted by ω Grph.

1.1.2. An ω -magma consists of an ω -graph X together with maps

$$1_{(-)}: X_n \to X_{n+1}$$
$$x \mapsto 1_x$$

for every $n \ge 0$, and maps

$$(-) *_{k} (-) : X_{n} \underset{X_{k}}{\times} X_{n} \to X_{n}$$
$$(x, y) \mapsto x *_{k} y$$

for all $0 \le k < n$, subject to the following axioms:

(a) For every $n \ge 0$ and every *n*-cell *x*,

$$\mathbf{s}(1_x) = \mathbf{s}(1_x) = x.$$

(b) For all $0 \le k < n$ and all k-composable n-cells x and y,

$$\mathbf{s}(x \underset{k}{*} y) = \begin{cases} \mathbf{s}(y) & \text{ when } k = n-1, \\ \mathbf{s}(x) \underset{k}{*} \mathbf{s}(y) & \text{ otherwise,} \end{cases}$$

and

$$\mathbf{t}(x * y) = \begin{cases} \mathbf{t}(x) & \text{when } k = n - 1, \\ \mathbf{t}(x) * \mathbf{t}(y) & \text{otherwise.} \end{cases}$$

We will use the same letter to denote an ω -magma and its underlying ω -graph.

For an *n*-cell x, the (n + 1)-cell 1_x is referred to as the *unit on* x. More generally, for all $0 \le k < n$, we define maps $1^{(n)}_{(-)} : C_k \to C_n$ as

$$1_{(-)}^{(n)} := \underbrace{1_{(-)} \circ \cdots \circ 1_{(-)}}_{n-k \text{ times}} : C_k \to C_n.$$

For a k-cell x and n > k, the n-cell $1_x^{(n)}$ is referred to as the n-dimensional unit on x. For consistency, we also set

$$l_r^{(k)} := x$$

for every k-cell x. An n-cell that is a unit on a strictly lower dimensional cell is sometimes referred to as a *trivial* n-cell.

For two k-composable n-cells x and y, the n-cell $x \underset{k}{*} y$ is referred to as the k-composition of x and y.

More generally, we extend the notion of k-composition for cells of different dimension in the following way. Let x be an n-cell, y be an m-cell with $m \neq n$ and $k < \min\{m, n\}$. The cells x and y are k-composable if $s_k(x) = t_k(y)$, in which case we define the cell x * y of dimension $\max\{m, n\}$ as

$$x \underset{k}{*} y := \begin{cases} 1_x^n \underset{k}{*} y & \text{if } m < n \\ x \underset{k}{*} 1_y^m & \text{if } m > n \end{cases}$$

We also follow the convention that if n < m, then $*_n$ has priority over $*_m$. This means that

$$x * y * z = (x * y) * z$$
 and $x * y * z = x * (y * z)$

whenever these equations make sense.

A morphism of ω -magmas $f : X \to Y$ is a morphism of underlying ω -graphs that is compatible with units and compositions, i.e. for every *n*-cell *x*, we have

$$f(1_x) = 1_{f(x)},$$

and for every k-composable n-cells x and y, we have

$$f(x * y) = f(x) * f(y).$$

We write ω Mag for the category of ω -magmas and morphisms of ω -magmas.

1.1.3. An ω -category is an ω -magma X that satisfies the following axioms:

Units: for all k < n, for every *n*-cell *x*, we have

$$1_{t_k(x)}^{(n)} * x = x = x * 1_{s_k(x)}^{(n)},$$

Functoriality of units: for all k < n and for all k-composable n-cells x and y, we have

$$1_{\substack{x \neq y \\ k}} = 1_x \underset{k}{*} 1_y,$$

Associativity: for all k < n, for all *n*-cells x, y and z such that x and y are k-composable, and y and z are k-composable, we have

$$(x * y) * z = x * (y * z),$$

Exchange rule: for all $k, l, n \in \mathbb{N}$ with k < l < n, for all *n*-cells x, x', y and y' such that

- x and y are *l*-composable, x' and y' are *l*-composable,

- x and x' are k-composable, y and y' are k-composable,

we have

$$(x * x') * (y * y') = (x * y) * (x' * y').$$

We will use the same letter to denote an ω -category and its underlying ω -magma. A *morphism of* ω -categories (or ω -functor) $f : X \to Y$ is simply a morphism of the underlying ω -magmas. We denote by ω Cat the category of ω -categories and morphisms of ω -categories. This category is locally presentable.

1.1.4. For $n \in \mathbb{N}$, the notions of *n*-graph, *n*-magma and *n*-category are defined as truncated version of ω -graph, ω -magma and ω -category in an obvious way. For example, a 0-category is a set and a 1-category is nothing but a usual (small) category. The category of *n*-categories and morphisms of *n*-categories (or *n*-functors) is denoted by *n*Cat. When n = 0 and n = 1, we almost always use the notation Set and Cat instead of 0Cat and 1Cat.

For every $n \ge 0$, there is a canonical functor

$$\tau_{\leq n}^s: \omega \mathbf{Cat} \to n\mathbf{Cat}$$

that simply discards all the cells of dimension strictly higher than n. This functor has a left adjoint

$$\iota_n: n\mathbf{Cat} \to \omega\mathbf{Cat},$$

where for an *n*-category C, the ω -category $\iota_n(C)$ has the same *k*-cells as C for $k \leq n$ and only unit cells in dimension strictly higher than n. This functor itself has a left adjoint

$$\tau_{\leq n}^i: \omega \mathbf{Cat} \to n\mathbf{Cat},$$

where for an ω -category C, the *n*-category $\tau_{\leq n}^i(C)$ has the same *k*-cells as C for k < nand whose set of *n*-cells is the quotient of C_n under the equivalence relation \sim generated by

 $x \sim y$ if there exists $z \in C_{n+1}$ of the form $z : x \to y$.

The functor $\tau_{< n}^s$ also have a right adjoint

$$\kappa_n : n\mathbf{Cat} \to \omega\mathbf{Cat},$$

where for an *n*-category C, the ω -category $\kappa_n(C)$ has the same *k*-cells as C for $k \leq n$ and has exactly one *k*-cell $x \to y$ for every pair of parallel (k-1)-cells (x, y) for k > n.

The sequence of adjunctions

$$\tau_{< n}^i \dashv \iota_n \dashv \tau_{< n}^s \dashv \kappa_n$$

is maximal in that κ_n doesn't have a right adjoint and $\tau_{< n}^i$ doesn't have a left adjoint.

The functors $\tau_{\leq n}^s$ and $\tau_{\leq n}^i$ are respectively referred to as the *stupid truncation functor* and the *intelligent truncation functor*.

The functor ι_n is fully faithful and preserves both limits and colimits; in regards to these properties, we often identify nCat with the essential image of ι_n , which is the full subcategory of ω Cat spanned by the ω -categories whose k-cells for k > n are all units.

1.1.5. For $n \ge 0$, we define the *n*-skeleton functor $sk_n : \omega Cat \rightarrow \omega Cat$ as

$$\operatorname{sk}_n := \iota_n \circ \tau^s_{\leq n}.$$

This functor preserves both limits and colimits. For an ω -category C, $\operatorname{sk}_n(C)$ is the sub- ω -category of C generated by the k-cells of C with $k \leq n$ in an obvious sense. It is also convenient to define $\operatorname{sk}_{-1}(C)$ to be the empty ω -category

$$\operatorname{sk}_{-1}(C) = \emptyset$$

for every ω -category C. Note that the functor $\mathrm{sk}_{-1} : \omega \mathbf{Cat} \to \omega \mathbf{Cat}$ preserves colimits but does not preserve limits. The inclusion induces a canonical filtration

$$\emptyset = \mathrm{sk}_{-1}(C) \hookrightarrow \mathrm{sk}_0(C) \hookrightarrow \mathrm{sk}_1(C) \hookrightarrow \cdots \hookrightarrow \mathrm{sk}_n(C) \hookrightarrow \cdots$$

and we leave the proof of the following lemma as an easy exercise for the reader.

Lemma 1.1.6. Let C be an ω -category. The colimit of the canonical filtration

$$\mathrm{sk}_{-1}(C) \hookrightarrow \mathrm{sk}_{0}(C) \hookrightarrow \mathrm{sk}_{1}(C) \hookrightarrow \cdots \hookrightarrow \mathrm{sk}_{n}(C) \hookrightarrow \cdots$$

is C and for $n \ge 0$ the universal arrow $\operatorname{sk}_n(C) \to C$ is given by the co-unit of the adjunction $\tau_{\le n}^s \dashv \iota_n$.

1.1.7. For $n \in \mathbb{N}$, the *n*-globe \mathbb{D}_n is the *n*-category that has:

- exactly one non-trivial *n*-cell, which we refer to as the *principal n*-cell of \mathbb{D}_n , and which we denote by e_n ,
- exactly two non-trivial k-cells for every k < n; these k-cells being parallel and given by the k-source and the k-target of e_n .

This completely describes the *n*-category \mathbb{D}_n as no non-trivial composition can occur. Here are pictures in low dimension:

$$\mathbb{D}_{0} = \bullet,$$
$$\mathbb{D}_{1} = \bullet \longrightarrow \bullet,$$
$$\mathbb{D}_{2} = \bullet \bigoplus^{*} \bullet,$$
$$\mathbb{D}_{3} = \bullet \bigoplus^{*} \bullet.$$

For every ω -category C, the map

$$\operatorname{Hom}_{\omega \operatorname{Cat}}(\mathbb{D}_n, C) \to C_n$$
$$F \mapsto F(e_n)$$

is a bijection natural in C. In other words, the n-globe represents the functor

$$\omega \mathbf{Cat} \to \mathbf{Set}$$
$$C \mapsto C_n.$$

For an n-cell x of C, we denote by

$$\langle x \rangle : \mathbb{D}_n \to C$$

the canonically associated ω -functor.

1.1.8. For $n \in \mathbb{N}$, the *n*-sphere \mathbb{S}_n is the *n*-category that has exactly two parallel non-trivial *k*-cells for every $k \leq n$. In other words, we have

$$\mathbb{S}_n = \mathrm{sk}_n(\mathbb{D}_{n+1}),$$

and in particular, we have a canonical inclusion functor

$$i_{n+1}: \mathbb{S}_n \to \mathbb{D}_{n+1}.$$

It is also customary to define \mathbb{S}_{-1} to be the empty ω -category and i_{-1} to be the unique ω -functor

$$\emptyset \to \mathbb{D}_0.$$

Notice that for every $n \ge 0$, the following commutative square

$$\begin{array}{ccc} \mathbb{S}_{n-1} & \stackrel{i_n}{\longrightarrow} & \mathbb{D}_n \\ & & \downarrow^{i_n} & & \downarrow^{j_n^+} \\ \mathbb{D}_n & \stackrel{j_n^-}{\longrightarrow} & \mathbb{S}_n, \end{array}$$

where we wrote j_n^+ (resp. j_n^-) for the morphism $\langle t(e_{n+1}) \rangle : \mathbb{D}_n \to \mathbb{S}_n$ (resp. $\langle s(e_{n+1}) \rangle : \mathbb{D}_n \to \mathbb{S}_n$), is cocartesian.

Here are some pictures of the n-spheres in low dimension:

$$S_{0} = \bullet,$$

$$S_{1} = \bullet,$$

$$S_{2} = \bullet (1)^{2} \bullet.$$

For an ω -category C and $n \ge 0$, an ω -functor

$$\mathbb{S}_n \to C$$

amounts to the data of two parallel *n*-cells of *C*. In other words, \mathbb{S}_n represents the functor $\omega \mathbf{Cat} \to \mathbf{Set}$ that sends an ω -category to the set of its parallel *n*-cells. For (x, y) a pair of parallel *n*-cells of *C*, we denote by

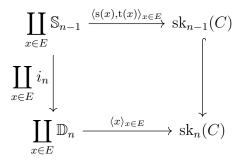
$$\langle x, y \rangle : \mathbb{S}_n \to C$$

the canonically associated ω -functor. For example, the ω -functor i_n is nothing but

$$\langle \mathbf{s}(e_{n+1}), \mathbf{t}(e_{n+1}) \rangle : \mathbb{S}_n \to \mathbb{D}_{n+1}.$$

1.2 FREE ω -CATEGORIES

Definition 1.2.1. Let C be an ω -category and $n \ge 0$. A subset $E \subseteq C_n$ of the n-cells of C is an *n*-basis of C if the commutative square



is cocartesian.

Remark 1.2.2. Note that since for all n < m, we have $sk_n \circ sk_m = sk_n$, an ω -category C has an n-basis if and only if $sk_n(C)$ has an n-basis.

1.2.3. Let us unfold Definition 1.2.1. For n = 0, E is an 0-basis of C if $E = C_0$. For n > 0, E is *n*-basis of C if for every *n*-category D, for every (n - 1)-functor

$$F: \tau^s_{< n-1}(C) \to \tau^s_{< n-1}(D),$$

and for every map

$$f: E \to D_n$$

such that for every $x \in E$,

s(f(x)) = F(s(x)) and t(f(x)) = F(t(x)),

there exists a *unique* n-functor

$$\tilde{F}: \tau^s_{\leq n}(C) \to D$$

such that $\tilde{F}_k = F_k$ for every k < n and $\tilde{F}_n(x) = f(x)$ for every $x \in E$.

Intuitively speaking, this means that $sk_n(C)$ has been obtained from $sk_{n-1}(C)$ by freely adjoining the cells in E.

Example 1.2.4. An *n*-category (seen as an ω -category) always has a *k*-basis for every k > n, namely the empty set.

Less trivial examples will come along soon.

Definition 1.2.5. An ω -category is *free*¹ if it has *n*-basis for every $n \ge 0$.

1.2.6. By considering *n*Cat as a subcategory of ω Cat, the previous definition also works for *n*-categories. It follows from Example 1.2.4 that an *n*-category is free if and only if it has a *k*-basis for every $0 \le k \le n$.

We now wish to recall an important result due to Makkai concerning the uniqueness of the *n*-basis for a free ω -category. First we need the following definition.

Definition 1.2.7. Let C be an ω -category. For n > 0, an n-cell x of C is *indecomposable* if both following conditions are satisfied:

- (a) x is not a unit on a lower dimensional cell,
- (b) if x is of the form

$$x = a \underset{k}{*} b$$

with k < n, then either

$$a = 1_{\mathbf{t}_k(x)}^{(n)},$$

or

$$b = 1_{\mathbf{s}_k(x)}^{(n)}.$$

For n = 0, all 0-cells are, by convention, indecomposable.

We can now state the promised result, whose proof can be found in [Mak05, Section 4, Proposition 8.3].

Proposition 1.2.8 (Makkai). Let C be a free ω -category. For every $n \in \mathbb{N}$, C has a unique n-basis. The cells of this n-basis are exactly the indecomposable n-cells of C.

Remark 1.2.9. Note that there is a subtlety in the previous proposition. It is not true in general that if an ω -category C has an n-basis then it is unique. The point is that we need the existence of the k-bases for k < n in order to prove the uniqueness of the n-basis. (See the paper of Makkai cited previously for details.)

1.2.10. Proposition 1.2.8 allows us to speak of *the* k-basis of a free ω -category C and more generally of the *basis* of C for the sequence

 $(\Sigma_k)_{k\in\mathbb{N}}$

where each Σ_k is the k-basis of C. In the case that C is a free n-category with n finite and in light of Example 1.2.4, we will also speak of *the basis of* C as the finite sequence

 $(\Sigma_k)_{0 \le k \le n}.$

¹Other common terminology for "free ω -category" is " ω -category free on a polygraph" [Bur93] or " ω -category free on a computad" [Str76, Mak05].

We often refer to the elements of the *n*-basis of a free ω -category as the generating *n*-cells. This sometimes leads to use the alternative terminology set of generating *n*-cells instead of *n*-basis.

Definition 1.2.11. Let C and D be two free ω -categories. An ω -functor $f : C \to D$ is *rigid* if for every $n \ge 0$ and every generating n-cell x of C, f(x) is a generating n-cell of D.

So far, we have not yet seen examples of free ω -categories. In order to do so, we will explain in a further section a recursive way of constructing free ω -categories; but let us first take a little detour.

1.3 SUSPENSION OF MONOIDS AND COUNTING GEN-ERATORS

1.3.1. Let M be a monoid. For every n > 0, let $B^n M$ be the n-magma such that:

- it has only one object \star ,
- it has only one k-cell for 0 < k < n, which is $1^{(k)}_{\star}$,
- the set of n-cells is (the underlying set of) M,
- for every k < n, the k-composition of n-cells is given by the composition law of the monoid (which makes sense since all n-cells are k-composable) and the only unital n-cell is given by the neutral element of the monoid.

It is sometimes useful to extend the above construction to the case n = 0 by saying that B^0M is the underlying set of the monoid M.

For n = 1, B^1M is nothing but the monoid M seen as a 1-category with one object.

For n > 1, while it is clear that all first three axioms for *n*-categories (units, functoriality of units and associativity) hold, it is not always true that the exchange rule is satisfied. If * denotes the composition law of the monoid, this axiom states that for all $a, b, c, d \in M$, we must have

$$(a * b) * (c * d) = (a * c) * (b * d).$$

It is straightforward to see that this equation holds if and only if M is commutative. Hence, we have proved the following lemma.

Lemma 1.3.2. Let M be a monoid and $n \in \mathbb{N}$. Then:

- if n = 1, B^1M is a 1-category,

- if n > 1, the *n*-magma $B^n M$ is an *n*-category if and only if M is commutative.

This construction will turn out to be of great use many times in this dissertation and we now explore a few of its properties.

Lemma 1.3.3. Let C be an n-category with $n \ge 1$ and let M = (M, *, 1) be a monoid (commutative if n > 1). The map

$$\operatorname{Hom}_{n\mathbf{Cat}}(C, B^n M) \to \operatorname{Hom}_{\mathbf{Set}}(C_n, M)$$
$$F \mapsto F_n$$

is injective and its image consists exactly of those functions $f: C_n \to M$ such that:

- for every $0 \le k < n$ and every pair (x, y) of k-composable n-cells of C, we have

$$f(x * y) = f(x) * f(y),$$

- for every $x \in C_{n-1}$, we have

 $f(1_x) = 1.$

Proof. The injectivity part follows from the fact that $(B^n M)_k$ is a singleton set for every k < n and hence, an *n*-functor $F : C \to B^n M$ is entirely determined by its restriction to the *n*-cells $F_n : C_n \to M$.

The characterization of the image is immediate once noted that the requirements are only the reformulation of the axioms of n-functors in this particular case.

Lemma 1.3.4. Let C be an n-category with $n \ge 1$ and M a monoid (commutative if n > 1). If C has an n-basis E, then the map

$$\operatorname{Hom}_{n\mathbf{Cat}}(C, B^n M) \to \operatorname{Hom}_{\mathbf{Set}}(E, M)$$
$$F \mapsto F_n|_E$$

is bijective.

Proof. This is an immediate consequence of the universal property of *n*-bases (as explained in Paragraph 1.2.3) \Box

We can now prove the important proposition below.

Proposition 1.3.5. Let C be an ω -category and suppose that C has an n-basis E with $n \ge 0$. For every $\alpha \in E$, there exists a unique function

$$w_{\alpha}: C_n \to \mathbb{N}$$

such that:

- (*a*) $w_{\alpha}(\alpha) = 1$,
- (b) $w_{\alpha}(\beta) = 0$ for every $\beta \in E$ such that $\beta \neq \alpha$,
- (c) for every $0 \le k < n$ and every pair (x, y) of k-composable n-cells of C, we have

$$w_{\alpha}(x * y) = w_{\alpha}(x) + w_{\alpha}(y).$$

Proof. Notice first that C has an n-basis if and only if $sk_n(C)$ has an n-basis (Remark 1.2.2). Hence we can suppose that C is an n-category.

For n = 0, conditions (c) is vacuous and the assertion is trivial.

Now let n > 0 and consider the monoid $\mathbb{N} = (\mathbb{N}, +, 0)$. The existence of a function $C_n \to \mathbb{N}$ satisfying conditions (a) and (b) follows from Lemma 1.3.4 and the fact that it satisfies (c) follows Lemma 1.3.3.

For the uniqueness, notice that for every $x \in C_{n-1}$ we have $1_x = 1_x \underset{n-1}{*} 1_x$ and thus condition (c) implies that

$$w_{\alpha}(1_x) = 0$$

Hence, we can apply Lemma 1.3.3 which shows that w_{α} necessarily comes from an *n*-functor $C \to B^n \mathbb{N}$. Then, the uniqueness follows from conditions (a) and (b) and Lemma 1.3.4.

1.3.6. Let C be an n-category with an n-basis E. For an n-cell x of C, we refer to the integer $w_{\alpha}(x)$ as the *weight of* α *in* x. The reason for such a name will become clearer after Remark 1.8.13 where we give an explicit construction of w_{α} as a function that "counts the number of occurrences of α in an n-cell".

For later reference, let us also highlight the fact that in the proof of the previous proposition, we have shown the important property that if n > 0, then for $y \in C_{n-1}$, we have

$$w_{\alpha}(1_y) = 0$$

This implies that for n > 1, there might be *n*-cells x such that

$$x \neq \alpha$$
 and $w_{\alpha}(x) = 1$.

Indeed, suppose that there exists a k-cell z with 0 < k < n-1 which is not a unit on a lower dimensional cell and such that $t_{k-1}(z) = s_{k-1}(\alpha)$, then we have

$$w_{\alpha}(\alpha *_{k-1} 1_{z}^{(n)}) = w_{\alpha}(\alpha) + w_{\alpha}(1_{z}^{(n)}) = 1.$$

1.4 RECURSIVE CONSTRUCTION OF FREE ω -CATEGORIES

Definition 1.4.1. Let $n \in \mathbb{N}$. An *n*-cellular extension is a quadruplet $\mathcal{E} = (C, \Sigma, \sigma, \tau)$ where:

- C is an n-category,
- Σ is a set, whose elements are referred to as the *indeterminates* of \mathcal{E} ,
- σ and τ are maps $\Sigma \to C_n$ such that for every element $x \in \Sigma$, the *n*-cells $\sigma(x)$ and $\tau(x)$ are parallel.

1.4.2. If we are given an *n*-category *C*, then we also say that an *n*-cellular extension \mathcal{E} is a *cellular extension of C* if it is of the form $\mathcal{E} = (C, \Sigma, \sigma, \tau)$.

Intuitively speaking, the indeterminates are formal extra (n + 1)-cells attached to C via σ and τ . For every $x \in \Sigma$, the *n*-cells $\sigma(x)$ and $\tau(x)$ are understood respectively the source and target of x (which makes sense since these two *n*-cells are parallel). Consequently, we often adopt the notation

$$x: a \to b$$

for an indeterminate such that $\sigma(x) = a$ and $\tau(x) = b$.

Definition 1.4.3. Let $\mathcal{E} = (C, \Sigma, \sigma, \tau)$ and $\mathcal{E}' = (C', \Sigma', \sigma', \tau')$ be two *n*-cellular extensions. A morphism of *n*-cellular extensions $\mathcal{E} \to \mathcal{E}'$ consists of a pair (F, φ) where:

- F is an *n*-functor $C \rightarrow C'$,
- φ is a map $\Sigma \to \Sigma'$,

such that for every $x \in \Sigma$, we have

$$\sigma'(\varphi(x)) = F(\sigma(x))$$
 and $\tau'(\varphi(x)) = F(\tau(x))$.

1.4.4. For $n \ge 0$, we denote by $n\mathbf{Cat}^+$ the category of *n*-cellular extensions and morphisms of *n*-cellular extensions. Every (n+1)-category *C* canonically defines an *n*-cellular extension $(\tau_{\le n}^s(C), C_{n+1}, \mathbf{s}, \mathbf{t})$ where $\mathbf{s}, \mathbf{t} : C_{n+1} \to C_n$ are the source and target maps of *C*. This defines a functor

$$U_n : (n+1)\mathbf{Cat} \to n\mathbf{Cat}^+$$
$$C \mapsto (\tau^s_{\leq n}(C), C_{n+1}, \mathbf{s}, \mathbf{t}).$$

On the other hand, every *n*-cellular extension $\mathcal{E} = (D, \Sigma, \sigma, \tau)$ yields an (n+1)-category \mathcal{E}^* defined as the following amalgamated sum:

This defines a functor

$$n\mathbf{Cat}^+ \to (n+1)\mathbf{Cat}$$
$$\mathcal{E} \mapsto \mathcal{E}^*,$$

which is easily checked to be left adjoint to
$$U_n$$
.

Now let $\phi : \coprod_{x \in \Sigma} \mathbb{D}_n \to \mathcal{E}^*$ be the bottom map of square (1.1). It induces a canonical map

$$j: \Sigma \to (\mathcal{E}^*)_{n+1}$$
$$x \mapsto \phi_x(e_{n+1}),$$

where e_{n+1} is the principal (n + 1)-cell of \mathbb{D}_{n+1} (1.1.7). Notice that this map is natural in that, for every morphism of *n*-cellular extensions

$$G = (F, \varphi) : \mathcal{E} \to \mathcal{E}',$$

the square

$$\begin{array}{cccc}
\Sigma & \xrightarrow{\varphi} & \Sigma' \\
\downarrow j & & \downarrow j' \\
(\mathcal{E}^*)_{n+1} & \xrightarrow{(G^*)_{n+1}} & (\mathcal{E}'^*)_{n+1}
\end{array}$$

is commutative. Notice also that j is compatible with source and target in the sense that for every $x \in \Sigma$, we have

$$s(j(x)) = \sigma(x)$$
 and $t(j(x)) = \tau(x)$.

Lemma 1.4.5. Let $\mathcal{E} = (C, \Sigma, \sigma, \tau)$ be an *n*-cellular extension. The canonical map

$$j: \Sigma \to (\mathcal{E}^*)_{n+1}$$

is injective.

Proof. A thorough reading of the techniques used in the proofs of Lemma 1.3.3, Lemma 1.3.4 and Proposition 1.3.5 shows that the universal property defining \mathcal{E}^* as the amalgamated sum (1.1) is sufficient enough to prove the existence, for each $x \in \Sigma$, of a function

$$w_x: (\mathcal{E}^*)_{n+1} \to \mathbb{N}$$

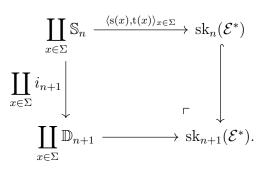
such that $w_x(j(x)) = 1$ and $w_x(j(y)) = 0$ for every $y \in \Sigma$ with $y \neq x$. In particular, this implies that j is injective.

1.4.6. In consequence of the previous lemma, we will often identify Σ with a subset of $(\mathcal{E}^*)_{n+1}$. When we do so, it will *always* be via the map j. This identification is compatible with source and target in the sense that the source (resp. target) of $x \in \Sigma$, seen as an (n + 1)-cell of \mathcal{E}^* , is exactly $\sigma(x)$ (resp. $\tau(x)$).

We can now prove the following proposition, which is the key result of this section. It is slightly less trivial than it appears.

Proposition 1.4.7. For every *n*-cellular extension $\mathcal{E} = (C, \Sigma, \sigma, \tau)$, the subset $\Sigma \subseteq (\mathcal{E}^*)_{n+1}$ is an (n+1)-basis of \mathcal{E}^* .

Proof. Notice first that since the map $i_{n+1} : \mathbb{S}_n \to \mathbb{D}_{n+1}$ is nothing but the canonical inclusion $\mathrm{sk}_n(\mathbb{D}_{n+1}) \to \mathrm{sk}_{n+1}(\mathbb{D}_{n+1}) = \mathbb{D}_{n+1}$, it follows easily from square (1.1) and the fact that the skeleton functors preserve colimits, that C is canonically isomorphic to $\mathrm{sk}_n(\mathcal{E}^*)$ and that the map $C \to \mathcal{E}^*$ can be identified with the canonical inclusion $\mathrm{sk}_n(\mathcal{E}^*) \to \mathrm{sk}_{n+1}(\mathcal{E}^*) = \mathcal{E}^*$. Hence, cocartesian square (1.1) can be identified with



Since we have identified Σ to a subset of the *n*-cells of \mathcal{E}^* via *j*, the above cocartesian square means exactly that Σ is an (n + 1)-basis of \mathcal{E}^* .

1.4.8. Let C be an (n + 1)-category and E be a subset $E \subseteq C_{n+1}$. This defines an *n*-cellular extension

$$\mathcal{E}_E = (\tau^s_{\leq n}(C), E, \mathbf{s}, \mathbf{t}),$$

where s and t are simply the restriction to E of the source and target maps $C_{n+1} \to C_n$. The canonical inclusion $E \hookrightarrow C_{n+1}$ induces a morphism of *n*-cellular extensions

$$\mathcal{E}_E \to U_n(C),$$

and then, by adjunction, an (n + 1)-functor

 $\mathcal{E}_E^* \to C.$

Proposition 1.4.9. Let C be an (n+1)-category. A subset $E \subseteq C_{n+1}$ is an (n+1)-basis of C if and only if the canonical (n+1)-functor

 $\mathcal{E}_E^* \to C$

is an isomorphism.

Proof. It is clear that the canonical (n + 1)-functor $\mathcal{E}_E^* \to C$ sends E, seen as a subset of $(\mathcal{E}_E^*)_{n+1}$, to E, seen as a subset of C_{n+1} . Hence, it follows from Proposition 1.4.7 that if this (n + 1)-functor is an isomorphism, then E is an (n + 1)-base of C.

Conversely, if E is an (n + 1)-base of C, then we can define an (n + 1)-functor $C \to \mathcal{E}_E^*$ that sends E, seen as a subset of C_{n+1} , to E, seen as a subset of $(\mathcal{E}_E^*)_{n+1}$ (and which is obviously the identity on cells of dimension strictly lower than n + 1). The fact that C and \mathcal{E}^* have E as an (n + 1)-base implies that this (n + 1)-functor $C \to \mathcal{E}^*$ is the inverse of the canonical one $\mathcal{E}^* \to C$.

1.4.10. We extend the definitions and the results from 1.4.1 to 1.4.9 to the case n = -1 by saying that a (-1)-cellular extension is simply a set Σ (which is the set of indeterminates) and (-1)Cat⁺ is the category of sets. Since a 0Cat is also the category of sets, it makes sense to define the functors

$$U_{-1}: 0\mathbf{Cat} \to (-1)\mathbf{Cat}^+$$

and

$$(-)^*: (-1)\mathbf{Cat}^+ \to 0\mathbf{Cat}$$

to be both the identity functor on Set.

Proposition 1.4.11. Let $(\mathcal{E}^{(n)})_{n>-1}$ be a sequence where:

-
$$\mathcal{E}^{(-1)}$$
 is a (-1) -cellular extension,

- for every $n \ge 0$, $\mathcal{E}^{(n)}$ is a cellular extension of the *n*-category $(\mathcal{E}^{(n-1)})^*$.

Then, the ω -category defined as the colimit of the canonical diagram

 $(\mathcal{E}^{(-1)})^* \to (\mathcal{E}^{(0)})^* \to \dots \to (\mathcal{E}^{(n)})^* \to \dots$

is free and for every $n \in \mathbb{N}$, its *n*-basis is (canonically isomorphic to) the set of indeterminates of $\mathcal{E}^{(n+1)}$. Moreover, suppose we are given another sequence $(\mathcal{E}'^{(n)})_{n\geq -1}$ as above and a sequence

$$(G^{(n)} = (F^{(n)}, \varphi^{(n)}) : \mathcal{E}^{(n)} \to \mathcal{E}^{\prime(n)})_{n \ge -1}$$

where each $G^{(n)}$ is a morphism of n-cellular extensions such that for every $n \ge 0$

$$F^{(n)} = (G^{(n-1)})^*$$

Then, the ω -functor

$$\operatorname{colim}_{n \ge -1} (\mathcal{E}^{(n)})^* \to \operatorname{colim}_{n \ge -1} (\mathcal{E}'^{(n)})^*$$

induced by colimit is rigid.

Conversely, every free ω -category and every rigid ω -functor arise this way.

Proof. From Proposition 1.4.7, we know that each $(\mathcal{E}^{(n)})^*$ has an (n + 1)-basis, which is canonically isomorphic to the set of indeterminates of $\mathcal{E}^{(n)}$. Besides, since for every $n \ge 0$, $\mathcal{E}^{(n)}$ is a cellular extension of $(\mathcal{E}^{(n-1)})^*$, we have

$$\operatorname{sk}_{n-1}((\mathcal{E}^{(n)})^*) = (\mathcal{E}^{(n-1)})^*$$

by definition. Hence, by a straightforward induction, each $\mathcal{E}^{(n)}$ is a free (n+1)-category and its k-basis for $0 \le k \le n+1$ is (canonically isomorphic to) the set of indeterminates of $\mathcal{E}^{(k-1)}$. Now let $C := \operatorname{colim}_{n \ge -1} (\mathcal{E}^{(n)})^*$. Since for every $k \ge 0$, sk_k preserves colimits and since $\operatorname{sk}_k((\mathcal{E}^{(n)})^*) = (\mathcal{E}^{(k-1)})^*$ for all $0 \le k < n$, we have that

$$\mathrm{sk}_k(C) = (\mathcal{E}^{(k-1)})^*$$

for every $k \ge 0$. Altogether, this proves that C is free and its k-basis is the set of indeterminates of $\mathcal{E}^{(k-1)}$ for every $k \ge 0$.

The fact that a sequence of morphisms of cellular extensions that satisfy the hypothesis given in the statement of the proposition induces a rigid ω -functor is proven in a similar fashion.

For the converse part, notice that a free ω -category C, whose basis is denoted by $(\Sigma_k)_{k \in \mathbb{N}}$, induces a sequence of cellular extensions:

$$\mathcal{E}_C^{(-1)} := \Sigma_0$$

and

$$\mathcal{E}_C^{(n)} := (\mathrm{sk}_n(C), \Sigma_{n+1}, \mathrm{s}, \mathrm{t}) \text{ for } n \ge 0.$$

It follows from Proposition 1.4.9 that $\operatorname{sk}_n(C) \simeq (\mathcal{E}_C^{(n-1)})^*$ and, then, from Lemma 1.1.6 that $C \simeq \operatorname{colim}_{n \ge -1} \mathcal{E}_C^{(n)}$.

Finally, notice that the construction $C \mapsto (\mathcal{E}_C^{(n)})_{n \ge -1}$ described above is obviously functorial with respect to rigid ω -functors and the isomorphism $\mathrm{sk}_n(C) \simeq (\mathcal{E}_C^{(n)})^*$ is natural with respect to rigid ω -functors. Since the statement of Lemma 1.1.6 is also natural in C, this easily implies that every rigid ω -functor arises as the colimit of sequence of morphisms of cellular extensions as described in the statement of the proposition. \Box **Remark 1.4.12.** The previous proposition admits an obvious truncated version for free *n*-categories with *n* finite. In that case, we only need a finite sequence $(\mathcal{E}^{(k)})_{1\leq k\leq n-1}$ of cellular extensions.

Remark 1.4.13. The data of a sequence $(\mathcal{E}^{(n)})_{n\geq -1}$ as in Proposition 1.4.11 is commonly referred to in the literature of the field as a *computad* [Str76] or *polygraph* [Bur93]; consequently a ω -category which is free in the sense of definition 1.2.5 is sometimes referred to as *free on a computad*-or-*polygraph*. Note that the underlying polygraph of a free ω -category is uniquely determined by the free ω -category itself (a straightforward consequence of Proposition 1.2.8), and this is why we chose the shorter terminology *free* ω -category.

1.4.14. Concretely, Proposition 1.4.11 gives us a recipe to construct free ω -categories. It suffices to give a formal list of generating cells of the form:

- generating 0-cells : x^0, y^0, \ldots
- generating 1-cells : $x^1 : \sigma(x^1) \to \tau(x^1), y^1 : \sigma(y^1) \to \tau(y^1), \ldots$
- generating 2-cells: $x^2 : \sigma(x^2) \to \tau(x^2), y^2 : \sigma(y^2) \to \tau(y^2), \dots$

- ...,

where for a generating k-cell x with k > 0, $\sigma(x)$ and $\tau(x)$ are parallel (k - 1)-cells of the free (k - 1)-category recursively generated by the generating cells of dimension strictly lower than k.

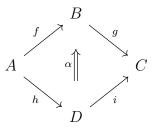
Example 1.4.15. The data of a 1-cellular extension \mathcal{E} is nothing but the data of a graph G (or 1-graph in the terminology of 1.1.4), and in that case it is not hard to see that \mathcal{E}^* is nothing but the free category on G. That is to say, the category whose objects are those of G and whose arrows are strings of composable arrows of G; the composition being given by concatenation of strings. Hence, from Proposition 1.4.11, a (1-)category is free in the sense of Definition 1.2.5 if and only if it is (isomorphic to) a free category on a graph.

Example 1.4.16. The notion of free category on a graph is easily generalized to the notion of free *n*-category on an *n*-graph (with $n \in \mathbb{N} \cup \{\infty\}$). As in the previous example, every free *n*-category on an *n*-graph is free in the sense of Definition 1.2.5. *However*, the converse is not true for n > 1. The point is that in a free *n*-category on an *n*-graph, the source and target of a *k*-generating cell must be (k - 1)-generating cells; whereas for a free *n*-category, they can be any (k-1)-cells (not necessarily generating). For example, the free 2-category described as

- generating 0-cells : A, B, C, D

- generating 1-cells : $f : A \to B, g : B \to C, h : A \to D, i : D \to C$
- generating 2-cells : $\alpha : i * h \Rightarrow g * f$,

which can be pictured as



is not free on a 2-graph. The reason is that the source (resp. the target) of α is g * f (resp. i * h) which are not generating 1-cells.

Example 1.4.17. Let $n \ge 1$ and M be a monoid (commutative if n > 1). The *n*-category $B^n M$ is free if and only if M is a free monoid (free commutative monoid if n > 1). If so, it has exactly one generating cell of dimension 0, no generating cells of dimension 0 < k < n, and the set of generators of the monoid (which is unique) as generating *n*-cells.

1.5 Cells of free ω -categories as words

In this section, we undertake to give a more explicit construction of the (n + 1)-category \mathcal{E}^* generated by an *n*-cellular extension $\mathcal{E} = (C, \Sigma, \sigma, \tau)$. By definition of \mathcal{E}^* , this amounts to give an explicit description of a particular type of colimit in ω Cat. Note also that since $\tau_{\leq n}(\mathcal{E}^*) = C$, all we need to do is to describe the (n + 1)-cells of \mathcal{E}^* . This will take place in two steps: first we construct what ought to be called the *free* (n + 1)-magma generated by \mathcal{E} , for which the (n + 1)-cells are really easy to describe, and then we quotient these cells as to obtain an (n + 1)-category, which will be \mathcal{E}^* .

Recall that an *n*-category is a particular case of *n*-magma.

1.5.1. Let $\mathcal{E} = (C, \Sigma, \sigma, \tau)$ be an *n*-cellular extension. We denote by $\mathcal{W}[\mathcal{E}]$ the set of finite words on the alphabet that has:

- a symbol \mathbf{c}_{α} for each $\alpha \in \Sigma$,
- a symbol \mathbf{i}_x for each $x \in C_n$,
- a symbol $\hat{*}_k$ for each $0 \le k \le n$,
- a symbol of opening parenthesis (,

- a symbol of closing parenthesis).

If w and w' are two elements of $\mathcal{W}[\mathcal{E}]$, we write ww' for their concatenation. We define the subset $\mathcal{T}[\mathcal{E}] \subseteq \mathcal{W}[\mathcal{E}]$ of well formed words (or terms) on the previous alphabet together with two maps s, t : $\mathcal{T}[\mathcal{E}] \to C_n$ in the following recursive way:

- for every $\alpha \in \Sigma$, the word (\mathbf{c}_{α}) is well formed and we have

$$s((\mathbf{c}_{\alpha})) = \sigma(\alpha)$$
 and $t((\mathbf{c}_{\alpha})) = \tau(\alpha)$,

- for every $x \in C_n$, the word (\mathbf{i}_x) is well formed and we have

$$\mathbf{s}((\mathbf{i}_x)) = \mathbf{t}((\mathbf{i}_x)) = x,$$

- if w and w' are well formed words such that s(w) = t(w'), then the word $(w \stackrel{*}{*} w')$ is well formed and we have

$$s((w * w')) = s(w')$$
 and $t((w * w')) = t(w)$,

if w and w' are well formed words such that s_k(s(w)) = t_k(t(w)) with 0 ≤ k < n, then the word (w ^{*}_k w') is well formed and we have

$$\mathbf{s}((w \stackrel{*}{k} w')) = \mathbf{s}(w) \underset{k}{*} s(w') \text{ and } \mathbf{t}((w \stackrel{*}{k} w')) = \mathbf{t}(w) \underset{k}{*} \mathbf{t}(w').$$

As usual, for $0 \le k < n$, we define $s_k, t_k : \mathcal{T}[\mathcal{E}] \to C_k$ to be respectively the iterated source and target (and we set $s_n = s$ and $t_n = t$ for consistency).

Definition 1.5.2. The *size* of a well-formed word w, denoted by |w|, is the number of symbols $\hat{*}_k$ for any $0 \le k \le n$ that appear in w.

1.5.3. Let \mathcal{E} be an *n*-cellular extension and let us write \mathcal{E}^+ for the (n+1)-magma defined in the following fashion:

for every 0 ≤ k ≤ n, we have (E⁺)_k := C_k; the source, target, compositions of k-cells for 0 < k ≤ n and units on k-cells for 0 ≤ k < n are those of C,

-
$$(\mathcal{E}^+)_{n+1} = \mathcal{T}[\mathcal{E}],$$

- the source and target maps $s, t : (\mathcal{E}^+)_{n+1} \to (\mathcal{E}^+)_n$ are the ones defined in 1.5.1,
- for every n-cell x, the unit on x is given by the word

$$(\mathbf{i}_x),$$

- for $0 \le k \le n$, the k-composition of two (n + 1)-cells w and w' such that $s_k(w) = t_k(w')$ is given by the word

$$(w \stackrel{\circ}{*} w').$$

By definition, \mathcal{E}^+ satisfy all the axioms for ω -categories up to dimension n. But on the other hand, the (n+1)-cells of \mathcal{E}^+ make it as far as possible from being an (n+1)-category as *none* of the axioms of ω -categories are satisfied for cells of dimension n+1.

Remark 1.5.4. Notice that, by definition, we have

$$w \underset{k}{*} w' := (w \underset{k}{*} w')$$

for every pair (w, w') of k-composable (n+1)-cells of \mathcal{E}^+ , and we ought to be careful not to confuse the "real" composition symbol "*" with the "formal" composition symbol "*". As a rule of thumb, it is better not to use both symbols in the same expression. Note also that, since we use the usual symbols "(" and ")" as formal symbols of opening and closing parenthesis, things can get really messy if we don't apply the previous rule because it would be hard to distinguish a formal parenthesis from an "non-formal" one.

In the following definition, we consider that a binary relation \mathcal{R} on a set E is nothing but a subset of $E \times E$, and we write $x \mathcal{R} x'$ to say $(x, x') \in \mathcal{R}$.

Definition 1.5.5. Let $n \ge 1$. A *congruence* on an *n*-magma X is a binary relation \mathcal{R} on the set of *n*-cells X_n such that:

- (a) \mathcal{R} is an equivalence relation,
- (b) if $x \mathcal{R} x'$ then x and x' are parallel,
- (c) if $x \mathcal{R} y$ and $x' \mathcal{R} y'$, and if x and x' are k-composable for some $0 \le k < n$ then

$$x \underset{k}{*} x' \mathcal{R} y \underset{k}{*} y'$$

(which makes sense since y and y' are k-composable by the previous axiom).

Remark 1.5.6. Beware that in the previous definition, the relation \mathcal{R} is *only* on the set of cells of dimension n.

Example 1.5.7. Let $F : X \to Y$ be a morphism of *n*-magmas with n > 1. The binary relation \mathcal{R} on X_n defined as

$$x \mathcal{R} y \text{ if } F(x) = F(y)$$

is a congruence.

1.5.8. Let X be an n-magma with $n \ge 1$ and \mathcal{R} a congruence on X. By the first axiom of Definition 1.5.5, \mathcal{R} is an equivalence relation and we can consider the quotient set X_n/\mathcal{R} . We write [x] for the equivalence class of an n-cell x of X. From the second axiom of Definition 1.5.5, we can define unambiguously

$$s([x]) := s(x)$$
 and $t([x]) := t(x)$,

for $x \in X_n$ and from the third axiom, we can define unambiguously

$$[x] \underset{k}{*} [y] := [x \underset{k}{*} y]$$

for every pair (x, y) of k-composable n-cells of X. Altogether, this defines an n-magma, which we denote by X/\mathcal{R} , whose set k-cells is X_k for $0 \le k < n$, and X_n/\mathcal{R} for k = n. The composition, source, target and units of cells of dimension strictly lower than n being those of X and the composition, source and target of n-cells being given by the above formulas.

Definition 1.5.9. Let \mathcal{R} be a congruence on an *n*-magma X with $n \ge 1$. We say that \mathcal{R} is *categorical* if it satisfies all four following axioms:

1. for every k < n and every *n*-cell *x* of *X*, we have

$$1_{\mathbf{t}_{k}(x)}^{(n)} \underset{k}{*} x \mathcal{R} x \text{ and } x \mathcal{R} x \underset{k}{*} 1_{\mathbf{s}_{k}(x)}^{(n)},$$

2. for every k < n and for all k-composable n-cells x and y of X, we have

$$1_{\substack{x \ast y \\ k}} \mathcal{R} \ 1_x \underset{k}{\ast} 1_y,$$

3. for every k < n, for all *n*-cells x, y and z of X such that x and y are k-composable, and y and z are k-composable, we have

$$(x * y) * z \mathcal{R} x * (y * z),$$

- 4. for all $k, l \in \mathbb{N}$ with k < l < n, for all *n*-cells x, x', y and y' of X such that
 - x and y are l-composable, x' and y' are l-composable,
 - x and x' are k-composable, y and y' are k-composable,

we have

$$(x \underset{k}{*} x') \underset{l}{*} (y \underset{k}{*} y') \mathcal{R} (x \underset{l}{*} y) \underset{k}{*} (x' \underset{l}{*} y').$$

Example 1.5.10. Let C be an n-category with n > 1, which we consider as an n-magma. The equality on the set of n-cells of C is, by definition, categorical.

Example 1.5.11. Let $F : X \to Y$ be a morphism of *n*-magmas with n > 1 and suppose that *Y* is an *n*-category. Then the binary relation \mathcal{R} on X_n defined as

$$x \mathcal{R} y \text{ if } F(x) = F(y)$$

is obviously a *categorical* congruence.

Similarly to the "stupid" truncation of ω -categories (1.1.4), given an (n+1)-magma X, we write $\tau_{\leq n}^{s}(X)$ for the *n*-magma obtained by simply forgetting the cells of dimension (n+1).

The following lemma is trivial but nonetheless important. Its immediate proof is omitted.

Lemma 1.5.12. Let X be an n-magma with n > 1 and \mathcal{R} a congruence on X. If $\tau_{\leq n}^{s}(X)$ is an (n-1)-category and \mathcal{R} is categorical, then X/\mathcal{R} is an n-category.

We now wish to see how to prove the existence of a congruence defined with a condition such as "the smallest congruence that contains a given binary relation on the (n + 1)-cells".

Lemma 1.5.13. Let X be an n-magma with $n \ge 1$ and $(\mathcal{R}_i)_{i \in I}$ a non-empty family of congruences on X (i.e. I is not empty). Then, the binary relation

$$\mathcal{R} := igcap_{i \in I} \mathcal{R}_i$$

is a congruence.

Proof. The fact that \mathcal{R} satisfies that first and third axiom of Definition 1.5.5 is immediate and do not even require that I be non empty. The only left thing to prove is that if $x \mathcal{R} x'$, then x and x' are parallel. To see that, notice that since I is not empty, we can choose $i \in I$. Then, by definition, we have $x \mathcal{R}_i x'$ and thus, x and x' are parallel. \Box

Remark 1.5.14. The hypothesis that *I* be non empty in the previous lemma cannot be omitted because, in this case, we would have

$$\mathcal{R} = \bigcap_{\emptyset} \mathcal{R}_i = X_n \times X_n,$$

which means that *all* n-cells would be equivalent under \mathcal{R} . This certainly does not guarantee that the second axiom of Definition 1.5.5, i.e. that equivalent cells are parallel, is satisfied.

Lemma 1.5.15. Let X be an n-magma with $n \ge 1$ and E a set of pairs of parallel *n*-cells of X. There exists a smallest congruence \mathcal{R} on X such that for every $(x, y) \in E$, we have $x \mathcal{R} y$.

Proof. Let I be the set of congruences S on X such that for every $(x, y) \in E$, we have $x \ S \ y$. All we have to prove is that I is not empty, since in that case, we can apply Lemma 1.5.13 to the binary relation

$$\mathcal{R} := \bigcap_{\mathcal{S} \in I} \mathcal{S},$$

which will obviously be the smallest congruence satisfying the desired condition. To see that I is not empty, it suffices to notice that the binary relation "being parallel *n*-cells" is a congruence, which obviously is in I.

Proposition 1.5.16. Let X be an n-magma with $n \ge 1$. There exists a smallest categorical congruence on X.

Proof. Each of the four axioms of Definition 1.5.9 says that some pairs of parallel n-cells must be equivalent under a congruence \mathcal{R} for it to be categorical. The result follows then from Lemma 1.5.15.

1.5.17. In particular, for every *n*-cellular extension \mathcal{E} there exists a smallest categorical congruence on \mathcal{E}^+ , which we denote here by \equiv . Applying Lemma 1.5.12 gives us that \mathcal{E}^+/\equiv is an (n + 1)-category. The construction

$$\mathcal{E}\mapsto \mathcal{E}^+/\equiv$$

clearly defines a functor $\operatorname{Cat}_n^+ \to (n+1)\operatorname{Cat}$.

We can now prove the expected result.

Proposition 1.5.18. Let \mathcal{E} be an *n*-cellular extension and let \equiv be the smallest categorical congruence on \mathcal{E}^+ . The (n + 1)-category \mathcal{E}^+/\equiv is naturally isomorphic to \mathcal{E}^* .

Proof. The strategy of the proof is to show that the functor $\mathcal{E} \mapsto \mathcal{E}^+ /\equiv$ is left adjoint to $U_n : (n+1)\mathbf{Cat} \to \mathbf{Cat}_n^+$. The result will follow then from the uniqueness (up to a natural isomorphism) of left adjoints.

Let $\mathcal{E} = (C, \Sigma, \sigma, \tau)$ be an *n*-cellular extension, let D be an (n+1)-category and let

$$G = (F, \varphi) : \mathcal{E} \to U_n(D)$$

be a morphism of n-cellular extensions. We recursively define a map

$$\overline{\varphi}: \mathcal{T}[\mathcal{E}] \to D_{n+1}$$

as:

-
$$\overline{\varphi}((\mathbf{c}_{\alpha})) = \varphi(\alpha)$$
 for $\alpha \in \Sigma$,

- $\overline{\varphi}((\mathbf{i}_x)) = \mathbb{1}_{F(x)}$ for $x \in C_n$,
- $\overline{\varphi}((w \stackrel{*}{\ast} w')) = \overline{\varphi}(w) \stackrel{*}{\ast} \overline{\varphi}(w').$

A straightforward induction shows that $\overline{\varphi}$ is compatible with source, target and units. Hence, we can define a morphism of (n + 1)-magmas

$$\overline{G}: \mathcal{E}^+ \to D$$

as

$$G_{n+1} = \overline{\varphi}$$

and

$$\overline{G}_k = F_k \text{ for } 0 \le k \le n.$$

Since D is an (n + 1)-category, the binary relation \mathcal{R} on $\mathcal{T}[\mathcal{E}]$ defined as

$$x \mathcal{R} y \text{ if } \overline{\varphi}(x) = \overline{\varphi}(y)$$

is a *categorical* congruence on \mathcal{E}^+ (see Example 1.5.11). Since \equiv is by definition the smallest categorical congruence on \mathcal{E}^+ , \equiv is included in \mathcal{R} . In particular, the map $\overline{\varphi}$ induces a map

$$\widehat{\varphi}: \mathcal{T}[\mathcal{E}] / \equiv \to D_{n+1}.$$

Since \equiv is a congruence, this map is compatible with source and target, and it is straightforward to check that it is also compatible with units. Hence, we have an (n+1)-functor

$$\widehat{G}: \mathcal{E}^+ / \equiv \to D$$

defined as

 $\widehat{G}_{n+1} = \widehat{\varphi}$

and

$$\tau^s_{\le n}(\widehat{G}) = F.$$

Altogether, we have constructed a map

$$\operatorname{Hom}_{\operatorname{Cat}_{n}^{+}}(\mathcal{E}, U_{n}(D)) \to \operatorname{Hom}_{(n+1)\operatorname{Cat}}(\mathcal{E}^{+}/\equiv, D)$$
$$G \mapsto \widehat{G}$$

which is clearly natural in \mathcal{E} and D. What we want to prove is that this map is a bijection.

Let us start with the surjectivity. Let $H: \mathcal{E}^+ / \equiv \to D$ be a (n+1)-functor. We define a map $\varphi: \Sigma \to D_{n+1}$ as

$$\varphi: \Sigma \to D_{n+1}$$
$$\alpha \mapsto H_{n+1}([(\mathbf{c}_{\alpha})]),$$

where [w] is the equivalence class under \equiv of an element $w \in \mathcal{T}[\mathcal{E}]$. All we need to show is that

$$\widehat{\varphi} = H_{n+1}.$$

Let z be an element of $\mathcal{T}[\mathcal{E}]/\equiv$ and let us choose $w \in \mathcal{T}[\mathcal{E}]$ such that z = [w]. We proceed to show that $\widehat{\varphi}(z) = H_{n+1}(z)$ by induction on the size of w (Definition 1.5.2). If |w| = 0, then either $w = (\mathbf{c}_{\alpha})$ for some $\alpha \in \Sigma$ or $w = (\mathbf{i}_x)$ for some $x \in C_n$. In the first case, we have

$$\widehat{\varphi}([(\mathbf{c}_{\alpha})]) = \overline{\varphi}((\mathbf{c}_{\alpha})) = \varphi(\alpha) = H_{n+1}([(\mathbf{c}_{\alpha})]),$$

and in the second case we have,

$$\widehat{\varphi}([(\mathbf{i}_x)]) = \overline{\varphi}((\mathbf{i}_x)) = 1_{H_n(x)} = H_{n+1}([(\mathbf{i}_x)])$$

where for the last equality, we used the fact that $[\mathbf{i}_x]$ is the unit on x in \mathcal{E}^+/\equiv .

Now if |w| = n + 1 with $n \ge 0$, then $w = (w' \stackrel{*}{*} w'')$ for some $w', w'' \in \mathcal{T}[\mathcal{E}]$ that are k-composable with $k \le n$. Hence, using the induction hypothesis, we have

$$\begin{aligned} \widehat{\varphi}([w]) &= \widehat{\varphi}([(w' \stackrel{*}{*} w'')]) = \widehat{\varphi}([w'] \stackrel{*}{*} [w'']) = \widehat{\varphi}([w']) \stackrel{*}{*} \widehat{\varphi}([w'']) \\ &= H_{n+1}([w']) \stackrel{*}{*} H_{n+1}([w'']) \\ &= H_{n+1}([w'] \stackrel{*}{*} [w'']) \\ &= H_{n+1}([(w' \stackrel{*}{*} w'')]) \\ &= H_{n+1}([w]). \end{aligned}$$

We now turn to the injectivity. Let $G = (F, \varphi)$ and $G' = (F', \varphi')$ be two morphisms of *n*-cellular extensions $\mathcal{E} \to U_n(D)$ such that $\widehat{G} = \widehat{G'}$. Since we have

$$F = \tau^s_{\leq n}(\widehat{G}) = \tau^s_{\leq n}(\widehat{G'}) = F',$$

all we have to show is that

$$\varphi = \varphi'.$$

But, by definition, for every $\alpha \in \Sigma$ we have

$$\varphi(\alpha) = \widehat{\varphi}([(\mathbf{c}_{\alpha})]) = \widehat{\varphi'}([(\mathbf{c}_{\alpha})]) = \varphi'(\alpha).$$

1.5.19. In other words, the previous proposition tells us that for an *n*-cellular extension \mathcal{E} , the (n + 1)-cells of \mathcal{E}^* can be represented as equivalence classes of (well formed) words made up of the indeterminates of \mathcal{E} and units on *n*-cells. Note, however, that

the equivalences classes are enormous. For example, for an indeterminate α , all three following words are in the equivalence class of (\mathbf{c}_{α}) :

$$((\mathbf{c}_{\alpha}) \stackrel{*}{\ast} (\mathbf{i}_{\sigma(\alpha)})) \quad (((\mathbf{c}_{\alpha}) \stackrel{*}{\ast} (\mathbf{i}_{\sigma(\alpha)})) \stackrel{*}{\ast} (\mathbf{i}_{\sigma(\alpha)})) \quad ((\mathbf{c}_{\alpha}) \stackrel{*}{\ast} ((\mathbf{i}_{\sigma(\alpha)}) \stackrel{*}{\ast} (\mathbf{i}_{\sigma(\alpha)}))).$$

Note also that when n = 0, we have already said that a 0-cellular extension \mathcal{E} is nothing but a graph and that \mathcal{E}^* is the free category on the graph. In particular, the 1-cells of \mathcal{E}^* can simply be encoded as strings of composable arrows of the graph, which is a much simpler description that the one obtained from the previous proposition since no equivalence relation is involved. A natural question to ask is whether it would be possible to obtain a similar simple description for n > 0. While it is certainly possible to have a simpler description of the (n + 1)-cells of \mathcal{E}^* than the one obtained in the present section (for example by getting rid of some "pathological" well formed words and reducing the size of the equivalence classes), it seems not possible to completely avoid the use of an equivalence relation, at least not in a canonical way. Indeed, we have already seen in Example 1.4.17 that a free *commutative* monoid M gives rise to a free n-category $B^n M$ whose generating n-cells are in bijection with the generators of the monoid. And, as soon as there are at least two generators of the monoid, say a and b, there is no canonical way of representing elements of the monoid by a unique word on the generators, for we have

$$ab = ba.$$

1.6 DISCRETE CONDUCHÉ ω **-FUNCTORS**

1.6.1. Recall that given a category C and a class M of arrows of C, we say that an arrow $f: X \to Y$ of C is *left orthogonal* to M if for every $m: A \to B$ in M and every solid arrow square

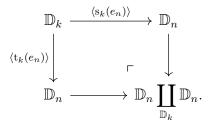
$$\begin{array}{ccc} X & \longrightarrow & A \\ f & & & \downarrow \\ Y & \longrightarrow & B, \end{array}$$

there exists a *unique* arrow $u: Y \to A$ making the whole diagram commute.

In the two following paragraphs, we freely use the notations from 1.1.7 and advise the reader to refer to it if needed.

1.6.2. For all $0 \leq k < n$, we denote by $\mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n$ the *n*-category defined as the

following amalgamated sum:



For an ω -category C, an arrow $\mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n \to C$ amounts to the same data as a pair of k-composable n-cells of C. Hence, $\mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n$ represents the functor

$$\omega \mathbf{Cat} \to \mathbf{Set}$$
$$C \mapsto C_n \times_{C_k} C_n.$$

Given a pair of k-composable n-cells (x, y) of an ω -category C, we write

$$\langle x, y \rangle : \mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n \to C$$

for the associated ω -functor.

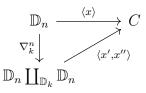
By the Yoneda Lemma, the k-composition operation

$$(-) \underset{k}{*} (-) : C_n \times_{C_k} C_n \to C_n,$$

which is obviously natural in C, yields an n-functor

$$abla_k^n: \mathbb{D}_n \to \mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n.$$

The *n*-cell ∇_k^n is nothing but the *k*-composition of the two generating *n*-cells of $\mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n$. In more practical terms, a commutative triangle



means exactly that x = x' * x''.

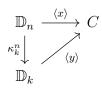
1.6.3. For all $0 \le k < n$, we denote by $\kappa_k^n : \mathbb{D}_n \to \mathbb{D}_k$, the ω -functor

$$\kappa_k^n := \langle 1_{e_k}^{(n)} \rangle : \mathbb{D}_n \to \mathbb{D}_k.$$

More conceptually, for every ω -category C the unit map

$$l_{(-)}^{(n)}: C_k \to C_n,$$

is obviously natural in C and the map κ_k^n is the one induced by the Yoneda Lemma. In practical terms, a commutative triangle



means exactly that $x = 1_y^n$.

Definition 1.6.4. An ω -functor $F : C \to D$ is *discrete Conduché* if it is right orthogonal to the ω -functors

$$\kappa_k^n: \mathbb{D}_n \to \mathbb{D}_k$$

and

$$\nabla_k^n: \mathbb{D}_n \to \mathbb{D}_n \coprod_{\mathbb{D}_k} \mathbb{D}_n$$

for all $k, n \in \mathbb{N}$ such that $0 \le k < n$.

Before unfolding the previous definition, let us put here for later reference the immediate following lemma.

Lemma 1.6.5. The class of discrete Conduché ω -functors is stable by pullback. This means that for every cartesian square in ω Cat

$$\begin{array}{ccc} C' & \longrightarrow & C \\ & & \downarrow & & \downarrow_F \\ F' \downarrow & & & \downarrow_F \\ D' & \longrightarrow & D, \end{array}$$

if F is a discrete Conduché ω -functor, then so is F'.

Proof. This is a standard fact about right orthogonal classes in a category that admits pullbacks. See for example the dual of [FK72, Proposition 2.1.1(e)]. \Box

1.6.6. Let us now give a practical version of Definition 1.6.4. The right orthogonality to κ_k^n means that for every *n*-cell *x* of *C* and every *k*-cell *y* of *D* such that

$$F(x) = 1_y^{(n)},$$

there exists a unique² k-cell z of C such that

$$x = 1_z^{(n)}$$
 and $F(z) = y$.

Similarly, the right orthogonality to ∇_k^n means that for every k-cell x of C, if

$$f(x) = y' * y'$$

with (y', y'') a pair of k-composable n-cells of D, then there exists a *unique* pair (x', x'') of k-composable n-cells of C such that

x = x' * x'',
 F(x') = y' and F(x'') = y''.

As it turns out, the definition we gave of discrete Conduché ω -functor is highly redundant.

Lemma 1.6.7. Let $F : C \to D$ be an ω -functor and let $k < n \in \mathbb{N}$. If F is right orthogonal to ∇_k^n then it is right orthogonal to κ_k^n .

Proof. Let x be an n-cell of C and suppose that $F(x) = 1_y^{(n)}$ with y a k-cell of D. Notice that

$$F(x) = 1_y^{(n)} * 1_y^{(n)}$$

and

$$x = x * 1_{\mathbf{s}_{k}(x)}^{(n)} = 1_{\mathbf{t}_{k}(x)}^{(n)} * x$$

and

$$F(1_{\mathbf{s}_{k}(x)}^{(n)}) = 1_{\mathbf{s}_{k}(F(x))}^{(n)} = 1_{y}^{(n)} = 1_{\mathbf{t}_{k}(F(x))}^{(n)} = F(1_{\mathbf{t}_{k}(x)}^{(n)}).$$

Using the uniqueness part of the right orthogonality to ∇_k^n , we deduce that $x = 1_{s_k(x)}^{(n)} = 1_{t_k(x)}^{(n)}$. Thus, if we set $z = s_k(x) = t_k(x)$, we have $x = 1_z^{(n)}$ and F(z) = y, which is what we needed to prove.

Remark 1.6.8. In light of the previous lemma, the reader might wonder why the right orthogonality to κ_k^n was included in Definition 1.6.4. The motivation for such a choice is that it *should* be possible to apply this definition *mutatis mutandis* to weak ω -categories where Lemma 1.6.7 might not hold anymore.

Somewhat related to Lemma 1.6.7 is the following lemma.

²The map $z \mapsto 1_k^n(z)$ being injective, the uniqueness actually comes for free.

Lemma 1.6.9. Let $F : C \to D$ be an ω -functor and $k < m < n \in \mathbb{N}$. If F is right orthogonal to ∇_k^n and κ_m^n , then it is right orthogonal to ∇_k^m .

Proof. Let x be an m-cell of C and suppose that

$$F(x) = y' * y''_k$$

with (y', y'') a pair of k-composable m-cells of D. Then, we have

$$F(1_x^{(n)}) = 1_{y'}^{(n)} * 1_{y''}^{(n)}.$$

From the right orthogonality to ∇_k^n , we know that there exist z and z'' k-composable n-cells of C such that $F(z') = 1_{y'}^{(n)}$, $F(z'') = 1_{y''}^{(n)}$ and

$$1_x^{(n)} = z' *_k z''.$$

From the right orthogonality to κ_m^n , we know that there exist x' and x'' *m*-cells of C such that $z' = 1_{x'}^{(n)}$, $z'' = 1_{x''}^{(n)}$, F(x') = y' and F(x'') = y''. It follows that $s_k(x') = s_k(x'')$ and

$$1_x^{(n)} = 1_{x'}^{(n)} * 1_{x''}^{(n)} = 1_{x'*x''}^{(n)},$$

hence x = x' * x''. This proves the existence part of the right orthogonality to ∇_k^m .

Now suppose that there are two pairs (x'_1, x''_1) and (x'_2, x''_2) that lift the pair (y', y'') in the required way. It follows that the pairs $(1^{(n)}_{x'_1}, 1^{(n)}_{x''_1})$ and $(1^{(n)}_{x'_2}, 1^{(n)}_{x''_2})$ lift the pair $(1^{(n)}_{y'}, 1^{(n)}_{y''})$ in the required way.

From the uniqueness part of the right orthogonality to ∇_k^n , we deduce that $1_{x'_1}^{(n)} = 1_{x'_2}^{(n)}$ and $1_{x''_1}^{(n)} = 1_{x''_2}^{(n)}$, hence $x'_1 = x'_2$ and $x''_1 = x''_2$.

Definition 1.6.4 admits an obvious truncated version as follows.

Definition 1.6.10. Let $m \in \mathbb{N}$. An *m*-functor $F : C \to D$ is a *discrete Conduché* if it is right orthogonal to κ_k^n and ∇_k^n for all $k, n \in \mathbb{N}$ such that $0 \le k < n \le m$.

1.6.11. Let us dwell on a subtlety here. Since we have considered nCat as a full subcategory of ω Cat for every $n \in \mathbb{N}$ (see 1.1.4), an *n*-functor is a particular case of ω -functor. Hence, it also makes sense to call an *n*-functor *discrete Conduché* when, seen as an ω -functor, it is discrete Conduché in the sense of Definition 1.6.4. This might have been conflicting with Definition 1.6.10 but the following lemma tells us that, in fact, the two notions are equivalent. Consequently, there is no distinction to make.

Lemma 1.6.12. Let $n \in \mathbb{N}$. An *n*-functor $F : C \to D$ is discrete Conduché in the sense of Definition 1.6.10 if and only if, seen as an ω -functor, it is discrete Conduché in the sense of Definition 1.6.4.

Proof. We only have to prove that if F is an discrete Conduché *n*-functor then, when seen as an ω -functor, it is a discrete Conduché ω -functor; the other implication being trivial. From Lemma 1.6.7, this amounts to showing that if F is right orthogonal to ∇_k^m for all $k, m \in \mathbb{N}$ such that $0 \le k < m \le n$, then it is also right orthogonal to ∇_k^m for all $k, m \in \mathbb{N}$ such that $0 \le k < m$ and m > n. This follows easily from the fact that for every *m*-cell *x* in an *n*-category (seen as an ω -category) with m > n, there exists a unique *n*-cell *x'* such that $x = 1_{x'}^{(m)}$. Details are left to the reader.

In practice, we will use the following criterion to detect discrete Conduché *n*-functors.

Proposition 1.6.13. Let $n \in \mathbb{N}$. An *n*-functor $F : C \to D$ is a discrete Conduché *n*-functor if and only if it is right orthogonal to ∇_k^n for every $k \in \mathbb{N}$ such that k < n.

Proof. Immediate consequence of Lemmas 1.6.7 and 1.6.9.

Remark 1.6.14. In the case n = 1, we recover the usual definition of what's commonly referred to as discrete Conduché fibration (see for example [Joh99]).

As we shall now see, discrete Conduché ω -functors have a deep connection with free ω -categories. We begin with an easy property.

Lemma 1.6.15. Let $F : C \to D$ be a discrete Conduché ω -functor and $n \in \mathbb{N}$. An *n*-cell x of C is indecomposable (Definition 1.2.7) if and only if F(x) is indecomposable.

Proof. When n = 0, there is nothing to prove since every 0-cell is indecomposable. We assume now that n > 0.

Suppose that x is indecomposable. Since x is not a unit on a lower dimensional cell, the right orthogonality to κ_k^n for every $0 \le k < n$ implies that F(x) is also not a unit on a lower dimensional cell. Now, if we have

$$F(x) = y' \underset{k}{*} y'',$$

with (y', y'') a pair of k-composable n-cells of D, then the right orthogonality to ∇_k^n implies that

$$x = x' * x''_k$$

with F(x') = y' and F(x'') = y''. Since x is indecomposable, x' or x'' has to be of the form $1_z^{(n)}$ with z and k-cell of C, and thus y' or y'' is of the form $1_{F(z)}^{(n)}$. This proves that F(x) is indecomposable.

Conversely, suppose that F(x) is indecomposable. The *n*-cell x cannot be a unit on a lower dimensional cell as if it were, then F(x) would also be a unit on a lower dimensional cell, which is impossible since it is indecomposable. Now, if

$$x = x' \underset{k}{*} x''$$

with (x', x'') a pair of k-composable n-cells, then

$$F(x) = F(x') \underset{k}{*} F(x'').$$

Hence, either F(x') or F(x'') is a unit on a lower dimensional cell and from the right orthogonality to κ_k^n we deduce that either x' or x'' is a unit on a lower dimensional cell. This proves that x is indecomposable.

From the previous lemma and Proposition 1.2.8, we deduce the following proposition.

Proposition 1.6.16. Let $F : C \to D$ be an ω -functor with C and D free ω -categories. If F is a discrete Conduché ω -functor then it is rigid.

Remark 1.6.17. The converse of the above proposition does not hold. For details, see [Gue20, Appendix A].

We now turn to the main result concerning the relation between discrete Conduché ω -functors and free ω -categories. As a matter of fact, a significant amount of results and definitions that we have seen so far were geared towards the statement and proof of the following theorem.

Theorem 1.6.18. Let $F : C \to D$ be a discrete Conduché ω -functor.

- 1. If D is free then so is C.
- 2. If C is free and if $F_n : C_n \to D_n$ is surjective for every $n \in \mathbb{N}$, then D is also free.

The proof of the previous theorem is long and technical and the next three sections are devoted to it. Before that, let us make Theorem 1.6.18 a little more precise.

1.6.19. Let $F: C \to D$ be a discrete Conduché ω -functor.

If D is free, we know from the previous theorem that C is also free. Now, if we write Σ_n^D for the n-basis of D, then it follows from Proposition 1.2.8 and Lemma 1.6.15 that the n-basis of C is given by

$$\Sigma_n^C := \left\{ x \in C_n | F(x) \in \Sigma_n^D \right\}.$$

2. Similarly, if C is free and if $F_k : C_k \to D_k$ is surjective for every $k \in \mathbb{N}$, we know that D is also free. If we write Σ_n^C for the *n*-basis of C, then it follows once again from Proposition 1.2.8 and Lemma 1.6.15 that the *n*-basis of D is given by

$$\Sigma_n^D := \left\{ F(x) \in D_n | x \in \Sigma_n^C \right\}.$$

1.7 **PROOF OF THEOREM 1.6.18: PART I**

This first part of the proof of Theorem 1.6.18 consists of several technical results on words. They lay a preliminary foundation on which the key arguments of the proof will later rely.

For the whole section, we fix a cellular *n*-extension $\mathcal{E} = (C, \Sigma, \sigma, \tau)$. A "word" always means an element of $\mathcal{W}[\mathcal{E}]$ and a "well formed word" always means an element of $\mathcal{T}[\mathcal{E}]$.

1.7.1. For a word w, the number of symbols that appear in w is referred to as the *length* of w and denoted by $\mathcal{L}(w)$. Recall also from Definition 1.5.2 that when w is well formed, the number of occurrences of symbols $\hat{*}_k$ for any $0 \le k \le n$ that appear in w is referred to as the *size of* w and denoted by |w|.

Definition 1.7.2. A word v is a *subword* of a word w if there exist words a and b such that w can be written as

w = avb.

Remark 1.7.3. Beware that in the previous definition, none of the words were supposed to be well formed. In particular, a subword of a well formed word is not necessarily well formed.

1.7.4. Since a word w is a finite sequence of symbols, it makes sense to write w(i) for the symbol at position i of w, with $0 \le i \le \mathcal{L}(w) - 1$.

For every $0 \le i \le \mathcal{L}(w) - 1$, define $P_w(i)$ to be the number of opening parentheses in w with position $\le i$ minus the number of closing parentheses in w with position $\le i$. This defines a function

$$P_w: \{0,\ldots,\mathcal{L}(w)-1\} \to \mathbb{Z}.$$

Remark 1.7.5. Such a counting function is standard in the literature about formal languages. For example see [HU79, chapter 1, exercice 1.4].

Definition 1.7.6. A word *w* is *well parenthesized* if:

- 1. it is not empty,
- 2. $P_w(i) \ge 0$ for every $0 \le i \le \mathcal{L}(w) 1$,
- 3. $P_w(i) = 0$ if and only if $i = \mathcal{L}(w) 1$.

1.7.7. It follows from the previous definition that the first letter of a well parenthesized word is necessarily an opening parenthesis and that the last letter is necessarily a closing parenthesis. Thus, the length of a well parenthesized word is not less than 2.

Moreover, it is immediate that if w_1 and w_2 are well parenthesized words, then for every $0 \le k \le n$, the word

$$(w_1 \stackrel{\circ}{*} w_2)$$

is well parenthesized.

Lemma 1.7.8. A well formed word is well parenthesized.

Proof. Let w be a well formed word. We proceed by induction on |w|. If |w| = 0, then w is either of the form

 (\mathbf{c}_{α})

or of the form

$$(\mathbf{i}_x).$$

In either case, the assertion is trivial. Now suppose that |w| > 0, we know by definition that

$$w = (w_1 \stackrel{\circ}{*} w_2)$$

with w_1, w_2 well formed words such that $|w_1|, |w_2| < |w|$. The desired properties follow easily from the induction hypothesis. Details are left to the reader.

The converse of the previous lemma is obviously not true. However, Lemma 1.7.10 below is a partial converse.

Lemma 1.7.9. Let w be a well parenthesized word of the form

$$w = (w_1 \stackrel{\circ}{*} w_2)$$

with w_1 and w_2 well parenthesized words and $0 \le k \le n$, and let v be a subword of w. If v is well parenthesized then one the following holds:

- *1.* v = w,
- 2. v is a subword of w_1 ,
- 3. v is a subword of w_2 .

Proof. Let a and b be words such that

$$avb = w = (w_1 \stackrel{\circ}{\ast} w_2).$$

Let $l_1, l_2, l, l_a, l_b, l_v$ respectively be the lengths of w_1, w_2, w, a, b, v . Notice that

$$l_a + l_v + l_b = l = l_1 + l_2 + 3.$$

Notice also that since v is well parenthesized, the following cases are forbidden:

1. $l_1 \leq l_a \leq l_1 + 1$, 2. $l_2 \leq l_b \leq l_2 + 1$, 3. $l_a \geq l - 1$, 4. $l_b \geq l - 1$.

Indeed, the first case would imply that the first letter of v is a closing parenthesis or the symbol $\hat{*}$. Similarly, the second case would imply that the last letter of v is an opening parenthesis or the symbol $\hat{*}$. The third and fourth cases would imply that $l_v < 2$ which is also impossible.

This leaves us with the following cases:

- 1. $l_a = 0$,
- 2. $l_b = 0$,
- 3. $0 < l_a < l_1$ and $0 < l_b < l_2$,
- 4. $0 < l_a < l_1$ and $l_b > l_2 + 1$,
- 5. $l_1 + 1 < l_a$ and $0 < l_b < l_2$.

If we are in the first case, then

$$P_w(j) = P_v(j)$$

for $0 \le j \le l_v - 1$. This implies that $P_w(l_v - 1) = 0$ which means that $l = l_v$, hence w = v.

By a similar argument left to the reader, we can show that the second case implies that w = v.

If we are in the fourth (resp. fifth) case, then it is clear that v is a subword of w_1 (resp. w_2).

Suppose now that we are in the third case. Intuitively, it means that the first letter of v is inside w_1 and the last letter of v is inside w_2 . Notice first that

$$l_a < l_1 < l_a + l_v - 3, \tag{(\star)}$$

where the inequality on the right comes from the fact that $l_v \ge 2$ (because v is well formed).

Besides, by definition of P_w ,

$$P_w(j) = P_v(j - l_a) + P_w(l_a)$$

for $l_a \leq j < l_v + l_a$. In particular, we have

$$1 = P_{w_1}(l_1 - 1) + 1 = P_w(l_1 - 1) = P_v(l_1 - 1) + P_w(l_a).$$

From (\star) and since v is well parenthesized, we deduce that

$$P_v(l_1 - 1) > 0.$$

Hence $P_w(l_a) \leq 0$, which is impossible since w is well formed and since $l_a < l - 1$. \Box

Lemma 1.7.10. Let w be a well formed word. Any well parenthesized subword of w is also well formed.

Proof. We proceed by induction on |w|. If |w| = 0, then w is either of the form

 (\mathbf{c}_{α})

or of the form

 $(\mathbf{i}_x).$

In both cases, the only well parenthesized subword of w is w itself, which is well formed by hypothesis.

Suppose now that |w| > 0 and let u be a well parenthesized subword of w. By definition,

$$w = (w_1 \stackrel{\circ}{*} w_2)$$

with $|w_1|, |w_2| < |w|$. By Lemmas 1.7.8 and 1.7.9, we have that either:

- u = w in which case u is well formed by hypothesis,
- u is a subword of w_1 and from the induction hypothesis we deduce that u is well formed,
- u is a subword of w_2 which is similar to the previous case.

Proposition 1.7.11. Let w be a well formed word of the form

$$w = (w_1 \stackrel{\circ}{\ast} w_2)$$

with w_1 and w_2 well formed words and $0 \le k \le n$, and let v be a subword of w. If v is well formed, then we are in one of the following cases:

l.
$$v = w$$
,

2. v is a subword of w_1 ,

3. v is a subword of w_2 .

Proof. This follows immediately from Lemma 1.7.8 and Lemma 1.7.9.

Proposition 1.7.12. Let u be a well formed word of the form

vew

with v, w and e words and such that e is well formed. If e' is a well formed word that is parallel to e, then the word

ve'w

is also well formed.

Proof. We proceed by induction on |u|.

Base case If |u| = 0, then necessarily v and w are both the empty word and the assertion is trivial.

Inductive step If $|u| \ge 1$, then

$$u = \left(u_1 \stackrel{\circ}{*} u_2\right)$$

with u_1 and u_2 well formed words such that $|u_1|, |u_2| < |u|$. By hypothesis, e is a subword of u and from Proposition 1.7.11, we are in one of the following cases:

- u = e in which case the assertion is trivial.

- e is a subword of u_1 , which means that there exist words \tilde{v}, \tilde{w} such that

$$u_1 = \tilde{v} e \tilde{w}.$$

Moreover, we have

$$v = (\tilde{v}$$

and

$$w = \tilde{w} \stackrel{\circ}{*} u_2).$$

By induction hypothesis, the word

 $\tilde{v}e'\tilde{w}$

is well formed and thus

$$(\tilde{v}e'\tilde{w} \stackrel{*}{\ast} u_2) = vew$$

is well formed.

- e is a subword of u_2 , which is symmetric to the previous case.

Lemma 1.7.13. Let w_1, w_2, w'_1, w'_2 be well parenthesized words, and let $0 \le k \le n$ and $0 \le k' \le n$ be such that

$$(w_1 \stackrel{*}{*} w_2) = (w'_1 \stackrel{*}{*} w'_2).$$

Then, $w_1 = w'_1$, $w_2 = w'_2$ and k = k'.

Proof. Let us define $l := \min(\mathcal{L}(w_1), \mathcal{L}(w'_1))$. Notice that

$$P_w(j) = P_{w_1}(j-1) + 1 = P_{w'_1}(j-1) + 1$$

for $0 < j \leq l$, hence

$$P_{w_1}(l-1) = P_{w'_1}(l-1).$$

Since w_1 and w'_1 are well parenthesized, one of the members of the last equality (and thus both) is equal to 0. This implies that $\mathcal{L}(w_1) = \mathcal{L}(w'_1)$ and the desired properties follow immediately from that.

Proposition 1.7.14. Let w_1, w'_1, w_2, w'_2 be well formed words, and let $0 \le k \le n$ and $0 \le k' \le n$ be such that $(w_1 \stackrel{*}{\ast} w_2)$ and $(w'_1 \stackrel{*}{\ast} w'_2)$ are well formed. If

$$(w_1 \stackrel{\circ}{*} w_2) = (w'_1 \stackrel{\circ}{*} w'_2),$$

then

$$w_1 = w'_1, w_2 = w'_2$$
 and $k = k'$.

Proof. This follows from Lemma 1.7.8 and Lemma 1.7.13.

Corollary 1.7.15. Let w be a well formed word and suppose that it can be written as

$$w = (w_1 \stackrel{\circ}{\ast} w_2)$$

with w_1 and w_2 well formed words and $0 \le k \le n$. Then $s_k(w_1) = t_k(w_2)$.

Proof. By hypothesis, $|w| \ge 1$. From the definition of well formed words, we know that w is of the form

$$(w_1' \stackrel{*}{*} w_2')$$

with w_1' and w_2' well formed words and $0 \leq k' \leq n$ such that

$$\mathbf{s}_{k'}(w_1') = \mathbf{t}_{k'}(w_2').$$

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From Proposition 1.7.14, we have that $w'_1 = w_1$, $w'_2 = w_2$ and k = k'.

1.8 PROOF OF THEOREM 1.6.18: PART II

Recall that we have seen in Proposition 1.5.16 that there exists a smallest categorical congruence on every *n*-magma X with $n \ge 1$. However, the description of this congruence that we used when proving its existence was rather abstract. The main goal of this section is to give a more concrete description of the smallest categorical congruence in the case that $X = \mathcal{E}^+$ for an *n*-cellular extension \mathcal{E} . This description will turn out to be crucial for the last part of the proof of Theorem 1.6.18.

Definition 1.8.1. Let $n \ge 0$, \mathcal{E} be an *n*-cellular extension and $u, u' \in \mathcal{T}[\mathcal{E}]$. An *el*ementary move from u to u' is a quadruple $\mu = (v, w, e, e')$ with $v, w \in \mathcal{W}[\mathcal{E}]$ and $e, e' \in \mathcal{T}[\mathcal{E}]$ such that

$$u = vew,$$
$$u' = ve'w,$$

and such that one of the following holds:

(1) e is of the form

$$((x \stackrel{\circ}{*} y) \stackrel{\circ}{*} z)$$

and e' is of the form

$$(x \stackrel{\circ}{*} (y \stackrel{\circ}{*} z))$$

with $x, y, z \in \mathcal{T}[\mathcal{E}]$ and $0 \le k \le n$,

(2) e is of the form

$$\left(\left(\mathbf{i}_{c}\right) \stackrel{\circ}{\ast}_{k} x\right)$$

r

and e' is of the form

with
$$x \in \mathcal{T}[\mathcal{E}], 0 \leq k \leq n \text{ and } c = 1_{\mathbf{t}_k(x)}^{(n)}$$
,

(3) e is of the form

$$(x \stackrel{\circ}{*}_{k} (\mathbf{i}_{c}))$$

r

and e' is of the form

with
$$x \in \mathcal{T}[\mathcal{E}], 0 \leq k \leq n \text{ and } c = 1_{\mathbf{s}_k(x)}^{(n)}$$
,

(4) e is of the form

$$((\mathbf{i}_c) \stackrel{\circ}{*}_{k} (\mathbf{i}_d))$$

and e' is of the form

$$(\mathbf{i}_{c*d})$$

with (c, d) a pair of k-composable n-cells of C with $0 \le k < n$,

(5) e is of the form

$$((x \stackrel{*}{*} y) \stackrel{*}{*} (z \stackrel{*}{*} t))$$

and e' is of the form

$$(x \stackrel{\circ}{*} z) \stackrel{\circ}{*} (y \stackrel{\circ}{*} t))$$

with $x, y, z, t \in \mathcal{T}[\mathcal{E}]$ and $0 \le l < k \le n$.

1.8.2. We will use the notation

 $\mu: u \to u'$

to say that μ is an elementary move from u to u'.

We write $\mathcal{G}[\mathcal{E}]$ for the graph (or 1-graph in the terminology of 1.1.4) defined as:

- the set of objects of $\mathcal{G}[\mathcal{E}]$ is $\mathcal{T}[\mathcal{E}]$,
- for all u, u' in $\mathcal{T}[\mathcal{E}]$, an arrow of $\mathcal{G}[\mathcal{E}]$ from u to u' is an elementary move from u to u'.

We will use the categorical notation

$$\mathcal{G}[\mathcal{E}](u, u')$$

for the set of arrows from u to u'.

Finally, we will also sometimes write

 $u \leftrightarrow u'$

to say that there exists an elementary move from u to u' or from u' to u.

Definition 1.8.3. Let \mathcal{E} be an *n*-cellular extension and $u, u' \in \mathcal{T}[\mathcal{E}]$. We say that the well formed words u and u' are *equivalent* and write

$$u \sim u'$$

if they are in the same connected component of $\mathcal{G}[\mathcal{E}]$. More precisely, this means that there exists a finite sequence $(u_j)_{0 \le j \le N}$ of well formed words with $u_0 = u$, $u_N = u'$ and $u_j \leftrightarrow u_{j+1}$ for $0 \le j < N$. The equivalence class of a well formed word u will be denoted by [u].

Lemma 1.8.4. Let $u, u' \in \mathcal{T}[\mathcal{E}]$. If $u \sim u'$, then u and u' are parallel.

Proof. Let

$$\mu = (v, w, e, e') : u \to u'$$

be an elementary move from u to u'. Let us show that s(u) = s(u') and t(u) = t(u') with an induction on $\mathcal{L}(v) + \mathcal{L}(w)$. Notice first that, by definition of elementary moves, $|u| \ge 1$ and thus

$$u = \left(u_1 \stackrel{\circ}{*} u_2\right)$$

with $u_1, u_2 \in \mathcal{T}[\mathcal{E}]$.

- **Base case** If $\mathcal{L}(v) + \mathcal{L}(w) = 0$, it means that v and w are both the empty word. It is then straightforward to check the desired property using Definition 1.8.1.
- **Inductive step** Suppose now that $\mathcal{L}(v) + \mathcal{L}(w) \ge 0$. Since *e* is a subword of *u* and is well formed, we know from Proposition 1.7.11 that we are in one of the following cases:
 - e = u, which is exactly the base case.
 - e is a subword of u_1 , which means that there exist $\tilde{v}, \tilde{w} \in \mathcal{T}[\mathcal{E}]$ such that

$$u_1 = \tilde{v} e \tilde{w}.$$

Moreover, we have

and

$$w = \tilde{w} \stackrel{\circ}{*} u_2).$$

 $v = (\tilde{v}$

From Proposition 1.7.12, the word

$$u_1' := \tilde{v}e'\tilde{w}$$

is well formed. Since $\mathcal{L}(\tilde{v}) + \mathcal{L}(\tilde{w}) < \mathcal{L}(v) + \mathcal{L}(w)$, we can use the induction hypothesis on

$$\tilde{\mu} := (\tilde{v}, \tilde{w}, e, e') : u_1 \to u'_1$$

This shows that $s(u_1) = s(u'_1)$ and $t(u_1) = t(u'_1)$ and since

$$u = (u_1 \stackrel{*}{*} u_2)$$
 and $u' = (u'_1 \stackrel{*}{*} u_2)$

it follows easily that s(u) = s(u') and t(u) = t(u').

- e is a subword of u_2 , which is symmetric to the previous case.

By definition of \sim , this suffices to show the desired property.

Lemma 1.8.5. Let $v_1, v_2, v'_1, v'_2 \in \mathcal{T}[\mathcal{E}]$ and $0 \le k \le n$ such that v_1 and v_2 are k-composable, and v'_1 and v'_2 are k-composable. If $v_1 \sim v_2$ and $v'_1 \sim v'_2$, then

$$(v_1 \stackrel{*}{*} v_2) \sim (v'_1 \stackrel{*}{*} v'_2).$$

Proof. Let

$$\mu = (v, w, e, e') : v_1 \to v'_1$$

be an elementary move. Set

$$\tilde{v} := (v)$$

and

$$\tilde{w} := w \stackrel{\circ}{*} v_2).$$

Hence, $(\tilde{v}, \tilde{w}, e, e')$ is an elementary move from $(v_1 \stackrel{*}{k} v_2)$ to $(v'_1 \stackrel{*}{k} v_2)$. Similarly, if we have an elementary move from v_2 to v'_2 , we obtain an elementary move from $(v_1 \stackrel{*}{k} v_2)$ to $(v_1 \stackrel{*}{k} v'_2)$. By definition of \sim , this suffices to show the desired property.

Lemma 1.8.6. Let \mathcal{R} be a categorical congruence on \mathcal{E}^+ and $u, u' \in \mathcal{T}[\mathcal{E}]$. If $u \sim u'$, then $u \mathcal{R} u'$.

Proof. The proof is very similar to the proof of Lemma 1.8.4. Let

$$\mu = (v, w, e, e') : u \to u'$$

be an elementary move from u to u'. Let us prove that $u \mathcal{R} u'$ with an induction on $\mathcal{L}(v) + \mathcal{L}(w)$. Once again, by definition of elementary moves, we have $|u| \ge 1$ and thus

$$u = (u_1 \stackrel{\circ}{*} u_2)$$

with $u_1, u_2 \in \mathcal{T}[\mathcal{E}]$.

- **Base case** If $\mathcal{L}(v) + \mathcal{L}(w) = 0$, then we have u = e and u' = e'. In this case, all the different cases of elementary moves from Definition 1.8.1 correspond to the different axioms of categorical congruences (Definition 1.5.9). Hence, we have $u \mathcal{R}u'$.
- **Inductive step** Suppose now that $\mathcal{L}(v) + \mathcal{L}(w) \ge 0$. Since *e* is a subword of *u*, we know from Proposition 1.7.11 that we are in one of the following cases:

- e = u, which is exactly the base case.

- e is a subword of u_1 , which means that there exist $\tilde{v}, \tilde{w} \in \mathcal{T}[\mathcal{E}]$ such that

$$u_1 = \tilde{v} e \tilde{w}.$$

Moreover, we have

$$v = (\tilde{v}$$

and

$$w = \tilde{w} \stackrel{\circ}{*} u_2).$$

From Proposition 1.7.12, the word

$$u'_1 := \tilde{v}e'\tilde{w}$$

is well formed and since $\mathcal{L}(\tilde{v}) + \mathcal{L}(\tilde{w}) < \mathcal{L}(v) + \mathcal{L}(w)$, we can apply the induction hypothesis on

$$\tilde{\mu} := (\tilde{v}, \tilde{w}, e, e') : u_1 \to u'_1$$

This proves that $u_1 \mathcal{R} u'_1$. Now, since we have

$$u = (u_1 \stackrel{\circ}{*} u_2) = u_1 \stackrel{\circ}{*} u_2$$

and

$$u' = (u'_1 \stackrel{\circ}{*} u_2) = u'_1 \stackrel{\circ}{*} u_2$$

(see Remark 1.5.4 for the distinction between the symbol " $\hat{*}$ " and the symbol " $\hat{*}$ ") and since \mathcal{R} is a congruence, we have

 $u \mathcal{R} u'$.

- e is a subword of u_2 , which is symmetric to the previous case.

Altogether, Lemmas 1.8.4, 1.8.5 and 1.8.6 prove the following proposition, which is what we were aiming for.

Proposition 1.8.7. Let \mathcal{E} be an *n*-cellular extension. The equivalence relation \sim on $\mathcal{T}[\mathcal{E}]$ is the smallest categorical congruence on \mathcal{E}^+ .

We end this section with yet another characterisation of *n*-bases of ω -categories.

1.8.8. Let C be an ω -category, $n \ge 0$ and Σ a subset of C_{n+1} , and consider the n-cellular extension

 $\mathcal{E}_{\Sigma} = (\tau^s_{< n}(C), \Sigma, \mathbf{s}, \mathbf{t})$

(see Paragraph 1.4.8). When there is no ambiguity on the rest of the data, we allow ourselves to write $\mathcal{W}[\Sigma]$, $\mathcal{T}[\Sigma]$ and $\mathcal{G}[\Sigma]$ instead of $\mathcal{W}[\mathcal{E}_{\Sigma}]$, $\mathcal{T}[\mathcal{E}_{\Sigma}]$ and $\mathcal{G}[\mathcal{E}_{\Sigma}]$.

We recursively define a map $\rho_{\Sigma} : \mathcal{T}[\Sigma] \to C_{n+1}$ as

- $\rho_{\Sigma}((\mathbf{c}_{\alpha})) = \alpha$ for every $\alpha \in \Sigma$,
- $\rho_{\Sigma}((\mathbf{i}_x)) = 1_x$ for every $x \in C_n$,
- $\rho_{\Sigma}((w \ \hat{*} \ w')) = \rho(w) \ \hat{*} \ \rho(w')$ for every $0 \le k \le n$ and every pair (w, w') of *k*-composable elements of $\mathcal{T}[\Sigma]$.

Intuitively speaking, ρ_{Σ} is to be understood as an "evaluation map": given w a well formed expression on units and formal cells of Σ , $\rho_{\Sigma}(w)$ is the evaluation of w as an (n+1)-cell of C.

Lemma 1.8.9. The map ρ_{Σ} is compatible with source and target, i.e. for $w \in \mathcal{T}[\Sigma]$, we have

$$s(\rho_{\Sigma}(w)) = s(w)$$
 and $t(\rho_{\Sigma}(w)) = t(w)$.

Proof. This is proved by an immediate induction left to the reader.

Lemma 1.8.10. Let w and w' be elements of $\mathcal{T}[\Sigma]$. If $w \sim w'$, then

$$\rho_{\Sigma}(w) = \rho_{\Sigma}(w').$$

Proof. From Lemma 1.8.9 we know that ρ_{Σ} is compatible with source and target and, by definition, ρ_{Σ} is compatible with composition and units in an obvious sense. Hence, we have a morphism of (n + 1)-magmas

$$F: \mathcal{E}^+_{\Sigma} \to \tau^s_{< n+1}(C)$$

defined as:

- $F_{n+1} = \rho_{\Sigma}$,
- F_k is the identity on C_k for every $0 \le k \le n$.

Since $\tau_{\leq n+1}^s(C)$ is an (n+1)-category, the binary relation \mathcal{R} on $\mathcal{T}[\Sigma]$ defined as $w \mathcal{R} w'$ if $\rho_{\Sigma}(w) = \rho_{\Sigma}(w')$ is a categorical congruence on \mathcal{E}_{Σ} (see Example 1.5.11). The result follows then from 1.8.7

1.8.11. Let C be an ω -category, $\Sigma \subseteq C_{n+1}$ and let a be an (n+1)-cell of C. We define $\mathcal{T}[\Sigma]_a$ to be the set

$$\mathcal{T}[\Sigma]_a := \{ w \in \mathcal{T}[\Sigma] \, | \, \rho_{\Sigma}(w) = a \}.$$

Lemma 1.8.10 implies that if $v \in \mathcal{T}[\Sigma]_a$ and $v \sim w$ then $w \in \mathcal{T}[\Sigma]_a$.

We define $\mathcal{G}[\Sigma]_a$ to be the full subgraph of $\mathcal{G}[\Sigma]$ whose set of objects is $\mathcal{T}[\Sigma]_a$.

Proposition 1.8.12. Let C be an ω -category, $n \ge 0$ and $\Sigma \subseteq C_{n+1}$. The set Σ is an (n+1)-basis of C is and only if for every $a \in C_{n+1}$, the graph $\mathcal{G}[\Sigma]_a$ is 0-connected (i.e. non-empty and connected).

More precisely, this means that for every $a \in C_{n+1}$:

- there exists $w \in \mathcal{T}[\Sigma]$ such that $\rho_{\Sigma}(w) = a$,
- for all $v, w \in \mathcal{T}[\Sigma]$, if $\rho_{\Sigma}(v) = a = \rho_{\Sigma}(w)$, then $v \sim w$.

Proof. From Lemma 1.8.10, the map ρ_{Σ} induces a map

$$\overline{\rho_{\Sigma}}: \mathcal{T}[\Sigma]/\sim \to C_{n+1}.$$

In fact, since \sim is a categorical congruence, this map is the (n + 1)-dimensional part of the (n + 1)-functor

$$\overline{F}: \mathcal{E}^+_{\Sigma} / \sim \to \tau^s_{< n+1}(C),$$

where F is the morphism of (n + 1)-magmas from the proof of Lemma 1.8.10. Since \sim is the smallest categorical congruence on \mathcal{E}_{Σ}^+ , it follows from Proposition 1.5.18 that $\mathcal{E}_{\Sigma}^+/\sim$ is (canonically isomorphic to) \mathcal{E}_{Σ}^* and it is easily seen that the functor \overline{F} is nothing but the (n + 1)-functor $\mathcal{E}_{\Sigma}^* \to \tau_{\leq n+1}^s(C)$ as obtained in 1.4.8. Hence, from Proposition 1.4.9, and the fact that \overline{F} is the identity on cells of dimension not greater than n, we deduce that Σ is an (n + 1)-basis of C if and only if ρ_{Σ} is an isomorphism. The result follows immediately from the fact that equivalence classes of \sim are in bijection with maximal connected components of $\mathcal{G}[\Sigma]$.

Remark 1.8.13. Let n > 0, C be an n-category with an n-basis Σ and $\alpha \in \Sigma$. It is immediate to check from the definition of elementary moves that for two equivalent well formed words $u \sim u'$ of $\mathcal{T}[\Sigma]$, the number of occurrences of \mathbf{c}_{α} in u and u' are the same. In particular, for every $a \in C_n$ we can define the integer $w_{\alpha}(a)$ to be the number of occurrences of \mathbf{c}_{α} in any well formed word u such that $\rho_{\sigma}(u) = a$. An immediate induction using the properties of ρ_{Σ} shows that this function $w_{\alpha} : C_n \to \mathbb{N}$ is the same as the one whose existence was established in Proposition 1.3.5.

1.9 PROOF OF THEOREM 1.6.18: PART III

In this section, we finally go back to Conduché ω -functors. The first two parts might be considered as preliminaries and the key points of the proof of Theorem 1.6.18 lie within this third and last part.

1.9.1. Let $F: C \to D$ be an ω -functor, n > 0, $\Sigma^C \subseteq C_n$ and $\Sigma^D \subseteq D_n$ such that $F_n(\Sigma^C) \subseteq \Sigma^D$. We recursively define a map

$$\widetilde{F}: \mathcal{W}[\Sigma^C] \to \mathcal{W}[\Sigma^D]$$

with

-
$$\widetilde{F}(\mathbf{c}_{\alpha}) = \mathbf{c}_{F(\alpha)}$$
 for $\alpha \in \Sigma^{C}$,

-
$$F(\mathbf{i}_x) = \mathbf{i}_{F(x)}$$
 for $x \in C_n$.

-
$$\tilde{F}(\hat{*}) = \hat{*}$$
 for $0 \le k < n$

-
$$\widetilde{F}(()) = 0$$

-
$$\widetilde{F}()) =).$$

Notice that for every word $w \in \mathcal{W}[\Sigma^C]$, we have

$$|\widetilde{F}(w)| = |w|$$
 and $\mathcal{L}(\widetilde{F}(w)) = \mathcal{L}(w)$.

Lemma 1.9.2. Let $F : C \to D$ be an ω -functor and let $\Sigma^C \subseteq C_n$ and $\Sigma^D \subseteq D_n$ be such that $F_n(\Sigma^C) \subseteq \Sigma^D$. For every $u \in \mathcal{W}[\Sigma^C]$:

- 1. if u is well formed then $\widetilde{F}(u)$ is well formed,
- 2. if $\tilde{F}(u)$ is well formed and if u is a subword (1.7.2) of a well formed word then it is also well formed.

Proof. The first part is proved with a short induction left to the reader. For the second part, first notice that the map

$$\widetilde{F}: \mathcal{W}[\Sigma^C] \to \mathcal{W}[\Sigma^D]$$

satisfies the following property:

For every $w \in \mathcal{W}[\Sigma^C]$, w is well parenthesized if and only if $\widetilde{F}(w)$ is well parenthesized.

It suffices then to apply Lemma 1.7.8 and then Lemma 1.7.10.

1.9.3. The first part of Lemma 1.9.2 shows that \widetilde{F} induces a map

$$\widetilde{F}: \mathcal{T}[\Sigma^C] \to \mathcal{T}[\Sigma^D].$$

Moreover, we have a commutative square

$$\begin{aligned} \mathcal{T}[\Sigma^C] & \stackrel{\rho_C}{\longrightarrow} C_n \\ & \downarrow_{\widetilde{F}} & \downarrow_{F_n} \\ \mathcal{T}[\Sigma^D] & \stackrel{\rho_D}{\longrightarrow} D_n \end{aligned}$$

where ρ_C and ρ_D respectively stand for ρ_{Σ^C} and ρ_{Σ^D} .

Thus, for every $a \in C_n$ we can define a map:

$$\widetilde{F}_a: \mathcal{T}[\Sigma^C]_a \to \mathcal{T}[\Sigma^D]_{F(a)}$$
$$w \mapsto \widetilde{F}(w).$$

Recall from Proposition 1.6.13 that for an ω -functor $F: C \to D$ and $n \ge 0, \tau_{\le n}^s(F)$ is a discrete Conduché *n*-functor if and only if F is right orthogonal to ∇_k^n for every $k \in \mathbb{N}$ such that k < n.

Proposition 1.9.4. Let $F : C \to D$ be an ω -functor and n > 0. The following conditions are equivalent:

- 1. $\tau_{\leq n}^{s}(F): \tau_{\leq n}^{s}(C) \to \tau_{\leq n}^{s}(D)$ is a discrete Conduché n-functor,
- 2. for every $\Sigma^D \subseteq D_n$ and $\Sigma^C := F^{-1}(\Sigma^D)$ and for every $a \in C_n$ the map

$$\widetilde{F}_a: \mathcal{T}[\Sigma^C]_a \to \mathcal{T}[\Sigma^D]_{F(a)}$$

defined above is bijective.

Proof. We begin with $1 \Rightarrow 2$.

Surjectivity Let us prove the following assertion:

$$\forall l \in \mathbb{N}, \forall a \in C_n, \forall w \in \mathcal{T}[\Sigma^D]_{F(a)} \text{ such that } |w| \leq l$$
$$\exists v \in \mathcal{T}[\Sigma^C]_a \text{ such that } \widetilde{F}_a(v) = w.$$

We proceed by induction on l.

Suppose first that l = 0. We are necessarily in one of two cases:

- 1. $w = (\mathbf{c}_{\beta})$ with $\beta \in \Sigma^{D}$. By hypothesis, $\rho_{D}(w) = F(a)$ and by definition of $\rho_{D}, \rho_{D}(w) = \beta$ thus $F(a) = \beta$. By definition of $\Sigma^{C}, a \in \Sigma^{C}$ and we can choose $v = (\mathbf{c}_{a})$.
- w = (i_y) with y ∈ D_{n-1}. By hypothesis, ρ_D(w) = F(a) and by definition of ρ_D, ρ_D(w) = 1_y thus F(a) = 1_y. Since τ^s_{≤n}(F) is a discrete Conduché n-functor, F is right orthogonal to κⁿ_{n-1}. Hence, there exists x ∈ C_{n-1} such that a = 1_x and F(x) = y. We can then choose v = (i_x) ∈ T[Σ^C]_a.

Now suppose that the assertion is true for a fixed $l \in \mathbb{N}$ and let $w \in \mathcal{T}[\Sigma^D]_{F(a)}$ be such that |w| = l + 1.

By definition of well formed words, we have

$$w = (w_1 \stackrel{\circ}{*} w_2)$$

with $0 \le k < n$ and $w_1, w_2 \in \mathcal{T}[\Sigma^D]$ such that $|w_1| \le l$ and $|w_2| \le l$. By hypothesis, $\rho_D(w) = F(a)$ and by definition of ρ_D ,

$$\rho_D(w) = \rho_D(w_1) \underset{k}{*} \rho_D(w_2)$$

and thus,

$$\rho_D(w_1) \underset{k}{*} \rho_D(w_2) = F(a).$$

Since by hypothesis F is right orthogonal to ∇_k^n , we know that there exist $a_1 \in C_n$ and $a_2 \in C_n$ that are k-composable and such that

$$a = a_1 \underset{k}{*} a_2, F(a_1) = \rho_D(w_1) \text{ and } F(a_2) = \rho_D(w_2).$$

Since $|w_1| \leq l$ and $|w_2| \leq l$, we can apply the induction hypothesis. Hence, there exist $v_1 \in \mathcal{T}[\Sigma^C]_{a_1}$ and $v_2 \in \mathcal{T}[\Sigma^C]_{a_2}$ such that

$$\widetilde{F}_{a_1}(v_1) = \widetilde{F}(v_1) = w_1$$
 and $\widetilde{F}_{a_2}(v_2) = \widetilde{F}(v_2) = w_2$.

Since ρ_C commutes with source and target by Lemma 1.8.9, v_1 and v_2 are k-composable and the word $(v_1 \stackrel{*}{*} v_2)$ is well formed. By definition of ρ_C , we have

$$\rho_C((v_1 \stackrel{*}{*} v_2)) = \rho_C(v_1) \stackrel{*}{*} \rho_C(v_2) = a_1 \stackrel{*}{*} a_2 = a_1$$

Thus, $(v_1 \stackrel{\circ}{*} v_2) \in \mathcal{T}[\Sigma^C]_a$ and

$$\widetilde{F}_a((v_1 \stackrel{*}{*} v_2) = \widetilde{F}((v_1 \stackrel{*}{*} v_2)) = (\widetilde{F}(v_1) \stackrel{*}{*} \widetilde{F}(v_2)) = (w_1 \stackrel{*}{*} w_2) = w.$$

Injectivity Let us prove the following assertion:

$$\forall l \in \mathbb{N}, \forall v \in \mathcal{T}[\Sigma^C]_a, \forall w \in \mathcal{T}[\Sigma^C]_a \text{ such that } |v| = |w| \le l$$
$$\widetilde{F}_a(v) = \widetilde{F}_a(w) \Rightarrow v = w$$

We proceed by induction on l.

Suppose first that l = 0. We are necessarily in one of four cases:

- 1. $v = (\mathbf{c}_{\alpha})$ and $w = (\mathbf{c}_{\beta})$ with α and β in Σ^{C} . By definition of ρ_{C} , we have $\alpha = \rho_{C}(v) = a = \rho_{C}(w) = \beta$. Hence, v = w.
- 2. $v = (\mathbf{i}_x)$ and $w = (\mathbf{i}_y)$ with x and y in C_{n-1} . By hypothesis, we have $\rho_C(v) = a = \rho_C(w)$ and by definition of ρ_C , $\mathbf{1}_x = \rho_C(v) = a = \rho_C(w) = \mathbf{1}_y$, thus x = y and v = w.
- 3. $v = (\mathbf{c}_{\alpha})$ and $w = (\mathbf{i}_{x})$ with $\alpha \in \Sigma^{C}$ and $x \in C_{n-1}$. By hypothesis, $(\mathbf{c}_{F(\alpha)}) = \widetilde{F}(v) = \widetilde{F}(w) = (\mathbf{i}_{F(x)})$ which is impossible.
- 4. $v = (\mathbf{i}_x)$ and $w = (\mathbf{c}_\alpha)$ with $\alpha \in \Sigma^C$ and $x \in C_{n-1}$, which is symmetric to the previous case.

Now suppose that the assertion is true for a fixed $l \in \mathbb{N}$ and let $v, w \in \mathcal{T}[\Sigma^C]$ such that |v| = |w| = l + 1 and $\widetilde{F}(v) = \widetilde{F}(w)$. By definition of well formed words, we have

$$v = (v_1 \stackrel{\circ}{*} v_2)$$

and

$$w = (w_1 \stackrel{\circ}{*} w_2)$$

with $|v_1|, |v_2|, |w_1|, |w_2| \le l$.

By hypothesis, we have

$$(\widetilde{F}(v_1) \stackrel{*}{\underset{k}{\ast}} \widetilde{F}(v_2)) = \widetilde{F}(v) = \widetilde{F}(w) = (\widetilde{F}(w_1) \stackrel{*}{\underset{k'}{\ast}} \widetilde{F}(w_2)).$$

From Proposition 1.7.14, we deduce that $\hat{*}_{k} = \hat{*}_{k'}$ and

$$\widetilde{F}(v_j) = \widetilde{F}(w_j)$$

for $j \in \{1, 2\}$.

In order to apply the induction hypothesis, we need to show that $\rho_C(v_j) = \rho_C(w_j)$ for $j \in \{1, 2\}$.

By hypothesis,

$$\rho_C(v_1) * \rho_C(v_2) = \rho_C(v) = a = \rho_C(w) = \rho_C(w_1) * \rho_C(w_2).$$

Hence,

$$F(\rho_C(v_1)) * F(\rho_C(v_2)) = F(a) = F(\rho_C(w_1)) * F(\rho_C(w_2))$$

Besides, $F(\rho_C(v_j)) = \rho_D(\widetilde{F}(v_j)) = \rho_D(\widetilde{F}(w_j)) = F(\rho_C(w_j))$. We deduce from the fact that F is right orthogonal to ∇_k^n that

$$\rho_C(v_j) = \rho_C(w_j)$$

for $j \in \{1, 2\}$.

From the induction hypothesis we have $v_j = w_j$ for $j \in \{1, 2\}$, hence v = w.

Now let us prove $2 \Rightarrow 1$.

Let $a \in C_n$ and suppose that $F(a) = b_1 \underset{k}{*} b_2$. We set $\Sigma^D = \{b_1, b_2\}$. By definition, $((\mathbf{c}_{b_1}) \underset{k}{*} (\mathbf{c}_{b_2})) \in \mathcal{T}[\Sigma^D]_{F(a)}$ and by hypothesis there exists a unique $v \in \mathcal{T}[\Sigma^C]_a$ such that $\widetilde{F}_a(v) = ((\mathbf{c}_{b_1}) \underset{k}{*} (\mathbf{c}_{b_2}))$. Since $|\widetilde{F}_a(v)| = |v| = 1$, we have

 $v = \left(v_1 \stackrel{\circ}{*} v_2\right)$

with $|v_1| = |v_2| = 0$, $s_{k'}(v_1) = t_{k'}(v_2)$ and $0 \le k' < n$. Thus,

$$(\widetilde{F}(v_1) \underset{k'}{*} \widetilde{F}(v_2)) = \widetilde{F}(v) = ((\mathbf{c}_{b_1}) \underset{k}{*} (\mathbf{c}_{b_2})).$$

Using Proposition 1.7.14, we deduce that k = k' and $\widetilde{F}(v_j) = (\mathbf{c}_{b_j})$ for $j \in \{1, 2\}$. We set $a_1 = \rho_C(v_1)$, $a_2 = \rho_C(v_2)$ and we have $s_k(a_1) = t_k(a_2)$,

$$a = \rho_C(v) = \rho_C(v_1) * \rho_C(v_2) = a_1 * a_2$$

and

$$F(a_j) = F(\rho_C(v_j)) = \rho_D(\widetilde{F}(v_j)) = \rho_D(\mathbf{c}_{b_j}) = b_j$$

for $j \in \{1, 2\}$, which proves the existence part of the right orthogonality to ∇_k^n .

Now suppose that we have $a_1, a'_1, a_2, a'_2 \in C_n$ with

$$s_k(a_1) = t_k(a_2) \text{ and } s_k(a'_1) = t_k(a'_2),$$

 $a_1 \underset{k}{*} a_2 = a'_1 \underset{k}{*} a'_2 = a,$
 $F(a_1) = F(a'_1) = b_1 \text{ and } F(a_2) = F(a'_2) = b$

By definition of $\Sigma^C = F^{-1}(\Sigma^D)$, we have $a_1, a'_1, a_2, a'_2 \in \Sigma^C$. Let us define w and w' as

$$w = ((\mathbf{c}_{a_1}) \stackrel{*}{*} (\mathbf{c}_{a_2})) \text{ and } w' = ((\mathbf{c}_{a'_1}) \stackrel{*}{*} (\mathbf{c}_{a'_2})).$$

We have $\rho_C(w) = \rho_C(w') = a$ and $\widetilde{F}(w) = ((\mathbf{c}_{b_1}) \stackrel{*}{\ast} (\mathbf{c}_{b_2})) = \widetilde{F}(w')$. The injectivity of \widetilde{F}_a implies that w = w', hence $a_1 = a'_1$ and $a_2 = a'_2$, which proves the uniqueness part of the right orthogonality to ∇_k^n .

1.9.5. Let $F : C \to D$ be an ω -functor, n > 0, $\Sigma^C \subseteq C_n$ and $\Sigma^D \subseteq D_n$ such that $F(\Sigma^C) \subseteq \Sigma^D$. It follows from the definition of $\widetilde{F} : \mathcal{T}[\Sigma^C] \to \mathcal{T}[\Sigma^D]$ and the definition of elementary move (1.8.1) that for an elementary move

$$\mu = (v, w, e, e') : u \to u'$$

with $u, u' \in \mathcal{T}[\Sigma^C]$, the quadruple

$$(\widetilde{F}(v), \widetilde{F}(w), \widetilde{F}(e), \widetilde{F}(e'))$$

is an elementary move from $\widetilde{F}(u)$ to $\widetilde{F}(u')$. Thus, we have defined a map

$$\mathcal{G}[\Sigma^C](u, u') \to \mathcal{G}[\Sigma^D](\widetilde{F}(u), \widetilde{F}(u')).$$

Together with the map $\widetilde{F} : \mathcal{T}[\Sigma^C] \to \mathcal{T}[\Sigma^D]$, this defines a morphism of graphs

$$\widetilde{F}: \mathcal{G}[\Sigma^C] \to \mathcal{G}[\Sigma^D]$$

and, by restriction, a morphism of graphs

$$\widetilde{F}_a: \mathcal{G}[\Sigma^C]_a \to \mathcal{G}[\Sigma^D]_{F(a)}$$

for every $a \in C_n$.

Lemma 1.9.6. With the notations of the above paragraph, the map

$$\mathcal{G}[\Sigma^C](u, u') \to \mathcal{G}[\Sigma^D](\widetilde{F}(u), \widetilde{F}(u'))$$

is injective.

Proof. Let (v_1, w_1, e_1, e_1') and (v_2, w_2, e_2, e_2') be two elementary moves from u to u' such that

$$(\widetilde{F}(v_1), \widetilde{F}(w_1), \widetilde{F}(e_1), \widetilde{F}(e_1')) = (\widetilde{F}(v_2), \widetilde{F}(w_2), \widetilde{F}(e_2), \widetilde{F}(e_2')).$$

In particular, we have

$$\mathcal{L}(v_1) = \mathcal{L}(v_2)$$
, $\mathcal{L}(w_1) = \mathcal{L}(w_2)$, $\mathcal{L}(e_1) = \mathcal{L}(e_2)$, $\mathcal{L}(e'_1) = \mathcal{L}(e'_2)$.

Since

$$v_1e_1w_1 = u = v_2e_2w_2$$
 and $v_1e'_1w_1 = u' = v_2e'_2w_2$,

we have

$$v_1 = v_2$$
, $w_1 = w_2$, $e_1 = e_2$, $e_1' = e_2'$.

Lemma 1.9.7. With the notations of paragraph 1.9.5, suppose that $\tau_{\leq n-1}^{s}(F)$ is a discrete Conduché (n-1)-functor. Let

 $\mu: v \to v'$

be an elementary move in $\mathcal{T}[\Sigma^D]$. If there exists $u \in \mathcal{T}[\Sigma^C]$ such that

$$\widetilde{F}(u) = v_{1}$$

then there exists $u' \in \mathcal{T}[\Sigma^C]$ and an elementary move

$$\lambda: u \to u'$$

such that

$$\widetilde{F}(u') = v' \text{ and } \widetilde{F}(\lambda) = \mu.$$

Proof. The proof is long and tedious as we have to check all the different cases of elementary moves. For the sake of clarity, we first outline a sketch of the proof that is common to all the cases of elementary moves and then we proceed to fill in the blanks successively for each case.

Let

$$\mu = (v_1, v_2, e, e') : v \to v'$$

be an elementary move. Since, by definition,

$$F(u) = v = v_1 e v_2,$$

u is necessarily of the form

 $u = u_1 \overline{e} u_2$

with $\overline{e}, u_1, u_2 \in \mathcal{W}[\Sigma^C]$ such that

 $\widetilde{F}(\overline{e}) = e$

and

$$F(u_j) = v_j$$

for $j \in \{1, 2\}$. From the second part of Lemma 1.9.2, we deduce that \overline{e} is well formed. In each different case, we will prove the existence of a well formed word $\overline{e'}$ parallel to \overline{e} such that

$$F(\overline{e'}) = e'$$

From Proposition 1.7.12, we deduce that the word

$$\overline{u'} := u_1 \overline{e'} u_2$$

is well formed. By definition, we have

$$\widetilde{F}(u') = v'.$$

Moreover, in each case, it will be immediate that the pair $(\overline{e}, \overline{e'})$ is such that the quadruple

$$\lambda := (u_1, u_2, \overline{e}, \overline{e'})$$

is an elementary move and that

$$F(\lambda) = \mu.$$

All that is left to prove now is the existence of $\overline{e'}$ with the desired properties.

First case: The word e is of the form

$$((x \stackrel{*}{\ast} y) \stackrel{*}{\ast} z)$$

and the word e' is of the form

$$(x \stackrel{\circ}{*}_{k} (y \stackrel{\circ}{*}_{k} z))$$

with $x, y, z \in \mathcal{T}[\Sigma^D]$. The word \overline{e} is then necessarily of the form

$$((\overline{x} \stackrel{\circ}{*} \overline{y}) \stackrel{\circ}{*} \overline{z})$$

Since $\widetilde{F}(\overline{e}) = e$, we deduce from Lemma 1.9.2 that $\overline{x}, \overline{y}, \overline{z}$ and $(\overline{x} \stackrel{*}{\ast} \overline{y})$ are well formed. From Corollary 1.7.15, we deduce that

$$\mathbf{s}_k(\overline{x}) = \mathbf{t}_k(\overline{y})$$

and

$$\mathbf{s}_k(\overline{y}) = \mathbf{t}_k(\overline{z})$$

Thus, the word

$$\overline{e'} := (\overline{x} \stackrel{\circ}{*} (\overline{y} \stackrel{\circ}{*} \overline{z}))$$

is well formed and it satisfies the desired properties.

Second case: The word e is of the form

$$(x \stackrel{\circ}{\ast}_k (\mathbf{i}_{1_z^{(n-1)}}))$$

x

and the word e' is of the form

with $x \in \mathcal{T}[\Sigma^D]$, $0 \le k < n$ and $z = s_k(x)$. Necessarily \overline{e} is of the form

$$(\overline{x} \stackrel{\circ}{*}_{k} (\mathbf{i}_{y}))$$

with $\overline{x} \in \mathcal{T}[\Sigma^C]$ (from Lemma 1.9.2 again) and $y \in C_{n-1}$ such that

 $\widetilde{F}(\overline{x}) = x$

and

$$\widetilde{F}(y) = 1_{z}^{(n-1)}.$$

Then, we set

 $\overline{e'} := \overline{x}.$

The only thing left to show is that

$$y = 1_{\mathbf{s}_k(\overline{x})}^{(n-1)}.$$

If k = n - 1, this follows from Corollary 1.7.15 and the fact that \overline{e} is well formed. If k < n - 1, we need first to use the fact that f is right orthogonal to κ_k^{n-1} to deduce that

$$y = 1_{\overline{z}}^{(n-1)}$$

for some $\overline{z} \in C_k$ such that $F(\overline{z}) = z$ and then use Corollary 1.7.15 and the fact that \overline{e} is well formed.

Third case: Similar to the second one with the unit on the left.

Fourth case: The word e is of the form

$$((\mathbf{i}_x) \stackrel{\circ}{*}_k (\mathbf{i}_y))$$

and the word e' is of the form

$$(\mathbf{i}_{x * y})$$

with $x, y \in D_{n-1}$ such that $s_k(x) = t_k(y)$. Necessarily, \overline{e} is of the form

$$((\mathbf{i}_{\overline{x}}) \stackrel{\circ}{*} (\mathbf{i}_{\overline{y}}))$$

with $\overline{x}, \overline{y} \in C_{n-1}$ such that

$$F(\overline{x}) = x \text{ and } F(\overline{y}) = y.$$

Using Corollary 1.7.15 and the fact that e is well formed, we deduce that $s_k(\overline{x}) = t_k(\overline{y})$. Thus, the word

$$\overline{e'} := (\mathbf{i}_{\overline{x} * \overline{y}})$$

is well formed. It satisfies all the desired properties.

Fifth case: The word e is of the form

$$((x \stackrel{*}{\ast} y) \stackrel{*}{\ast} (z \stackrel{*}{\ast} t))$$

and the word e^\prime is of the form

$$\left(\left(x \stackrel{\circ}{*} z\right) \stackrel{\circ}{*} \left(y \stackrel{\circ}{*} t\right)\right)$$

with $x, y, z, t \in \mathcal{T}[\Sigma^D]$ and $0 \leq l < k < n$ such that all the compatibilities of sources and targets needed are satisfied.

Necessarily, \overline{e} is of the form

$$\left(\left(\overline{x} \stackrel{*}{\ast} \overline{y}\right) \stackrel{*}{\ast} \left(\overline{z} \stackrel{*}{\ast} \overline{t}\right)\right)$$

with $\overline{x}, \overline{y}, \overline{z}, \overline{t} \in \mathcal{W}[\Sigma^C]$ such that

$$\begin{split} \bar{F}(\overline{x}) &= x, \\ \tilde{F}(\overline{y}) &= y, \\ \tilde{F}(\overline{z}) &= z, \\ \tilde{F}(\overline{t}) &= t. \end{split}$$

From Lemma 1.9.2 and the fact that \overline{e} is well formed, we deduce that the words $\overline{x}, \overline{y}, \overline{z}, \overline{t}, (\overline{x} \stackrel{*}{\ast} \overline{y})$ and $(\overline{z} \stackrel{*}{\ast} \overline{t})$ are well formed and from Corollary 1.7.15, we deduce that

$$s_k(\overline{x}) = t_k(\overline{y}),$$

 $s_k(\overline{z}) = t_k(\overline{t})$

and

$$s_l((\overline{x} \stackrel{*}{\ast} \overline{y})) = t_l((\overline{z} \stackrel{*}{\ast} \overline{t})).$$

Since l < k, we deduce from this last equality that

$$s_l(\overline{x}) = s_l(\overline{y}) = t_l(\overline{z}) = t_l(\overline{t}).$$

Thus, the word

$$\overline{e'} := \left(\left(\overline{x} \stackrel{*}{\ast} \overline{z} \right) \stackrel{*}{\ast} \left(\overline{y} \stackrel{*}{\ast} \overline{t} \right) \right)$$

is well formed. It satisfies all the desired properties.

Remark 1.9.8. In the proof of the previous theorem, we have only used the hypothesis that F is right orthogonal to κ_k^n for every k such that $0 \le k < n - 1$.

Lemma 1.9.9. Let $F : C \to D$ be an ω -functor, n > 0, $\Sigma^D \subseteq D_n$ and $\Sigma^C := F^{-1}(\Sigma^D)$. If $\tau^s_{\leq n}(F)$ is a discrete Conduché *n*-functor, then for every $a \in C_n$

$$\widetilde{F}_a: \mathcal{G}[\Sigma^C]_a \to \mathcal{G}[\Sigma^D]_{F(a)}$$

is an isomorphism of graphs.

Proof. Proposition 1.9.4 exactly says that the map

$$\widetilde{F}_a: \mathcal{G}[\Sigma^C]_a \to \mathcal{G}[\Sigma^D]_{F(a)}$$

is an isomorphism on objects and we know from Lemma 1.9.6 that it is a faithful morphism of graphs (same definition as for functors). All that is left to show is that it is also full.

In other words, we have to show that for all $u, u' \in \mathcal{T}[\Sigma^C]_a$ the map

$$\mathcal{G}[\Sigma^C](u, u') \to \mathcal{G}[\Sigma^D](\widetilde{F}(u), \widetilde{F}(u'))$$

is surjective.

Let $\mu : \widetilde{F}(u) \to \widetilde{F}(u')$ be an element of the codomain. From Lemma 1.9.7 we know that there exists

$$\lambda: u \to \iota$$

in $\mathcal{G}[\Sigma^C]$ such that

$$F(\lambda) = \mu$$

In particular, we have

$$\widetilde{F}(v) = \widetilde{F}(u').$$

Since we have an elementary move from u to v and by hypothesis $u \in \mathcal{T}[\Sigma^C]_a$, we also have $v \in \mathcal{T}[\Sigma^C]_a$ (see 1.8.11). Using the injectivity of the map

$$\widetilde{F}_a: \mathcal{T}[\Sigma^C]_a \to \mathcal{T}[\Sigma^D]_{F(a)}$$

we conclude that v = u'.

Proposition 1.9.10. Let $F : C \to D$ be an ω -functor, $n \in \mathbb{N}$, $\Sigma^D \subseteq D_n$ and define $\Sigma^C := F^{-1}(\Sigma^D)$. If $\tau^s_{\leq n}(F)$ is a discrete Conduché *n*-functor, then:

1. if Σ^D is an *n*-basis then so is Σ^C ,

2. if $F_n : C_n \to D_n$ is surjective and Σ^C is an *n*-basis then so is Σ^D .

Proof. The case n = 0 is trivial. We now suppose that n > 0. From Lemma 1.9.9 we have that for every $a \in C_n$, the map

$$\widetilde{F}_a: \mathcal{G}[\Sigma^C]_a \to \mathcal{G}[\Sigma^D]_{F(a)}$$

is an isomorphism of graphs. In particular, $\mathcal{G}[\Sigma^C]_a$ is 0-connected if and only if $\mathcal{G}[\Sigma^D]_{F(a)}$ is 0-connected. We conclude with Proposition 1.8.12.

Putting all the pieces together, we finally have the awaited proof.

1.9.11 (Proof of Theorem 1.6.18). Let $F: C \to D$ be a discrete Conduché ω -functor.

- 1. In the case that D is free, it follows immediately from the first part of the Proposition 1.9.10 that C is free.
- 2. In the case that C is free and $F_n : C_n \to D_n$ is surjective for every $n \in \mathbb{N}$, let us write Σ_n^C for the *n*-basis of C. It follows from Proposition 1.2.8 and Lemma 1.6.15 that

$$F^{-1}(F(\Sigma_n^C)) = \Sigma_n^C$$

Hence, we can apply the second part of Proposition 1.9.10 and C is free.

 \square

CHAPTER 1. YOGA OF ω -CATEGORIES

CHAPTER 2 Homotopical algebra

The present chapter stands out from the others as it contains no original results. Its goal is simply to introduce the language and tools of homotopical algebra that we shall need in the rest of the dissertation. Consequently, most of the results are simply asserted and the reader will find references to the literature for the proofs. The main notion of homotopical algebra we aim for is the one of *homotopy colimits* and our language of choice is that of Grothendieck's theory of *derivators* [Gro90]. We do not assume that the reader is familiar with this theory and will quickly recall the basics. If needed, gentle introductions can be found in [Mal01] and in a letter from Grothendieck to Thomason [Gro91]; more detailed introductions can be found in [Gro13] and in the first section of [Cis03]; finally, a rather complete (yet unfinished and unpublished) textbook on the subject is [Gro16].

2.1 LOCALIZATION, DERIVATION

2.1.1. A *localizer* is a pair $(\mathcal{C}, \mathcal{W})$ where \mathcal{C} is a category and \mathcal{W} is a class of arrows of \mathcal{C} , which we usually refer to as the *weak equivalences*. We denote by $ho^{\mathcal{W}}(\mathcal{C})$, or simply $ho(\mathcal{C})$ when there is no ambiguity, the localization of \mathcal{C} with respect to \mathcal{W} and by

$$\gamma: \mathcal{C} \to \mathrm{ho}(\mathcal{C})$$

the localization functor [GZ67, 1.1]. Recall the universal property of the localization: for every category \mathcal{D} , the functor induced by pre-composition

$$\gamma^* : \underline{\operatorname{Hom}}(\operatorname{ho}(\mathcal{C}), \mathcal{D}) \to \underline{\operatorname{Hom}}(\mathcal{C}, \mathcal{D})$$

is fully faithful and its essential image consists of those functors $F : \mathcal{C} \to \mathcal{D}$ that send the morphisms of \mathcal{W} to isomorphisms of \mathcal{D} .

We shall always consider that C and ho(C) have the same class of objects and implicitly use the equality

$$\gamma(X) = X$$

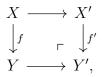
for every object X of C.

The class of arrows W is said to be *saturated* when we have the property:

 $f \in \mathcal{W}$ if and only if $\gamma(f)$ is an isomorphism.

For later reference, we put here the following definition.

Definition 2.1.2. Let $(\mathcal{C}, \mathcal{W})$ be a localizer such that \mathcal{C} has amalgamated sums. A morphism $f : X \to Y$ in \mathcal{W} is a *co-universal weak equivalence* if for every cocartesian square of the form



the morphism f' is also a weak equivalence.

2.1.3. A morphism of localizers $F : (C, W) \to (C', W')$ is a functor $F : C \to C'$ that preserves weak equivalences, i.e. such that $F(W) \subseteq W'$. The universal property of the localization implies that F induces a canonical functor

$$\overline{F}$$
: ho(\mathcal{C}) \rightarrow ho(\mathcal{C}')

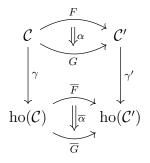
such that the square

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{C}' \\ & & & \downarrow^{\gamma'} \\ & & & \downarrow^{\gamma'} \\ & & & \text{ho}(\mathcal{C}) & \xrightarrow{\overline{F}} & & \text{ho}(\mathcal{C}'). \end{array}$$

is commutative. Let $G : (\mathcal{C}, \mathcal{W}) \to (\mathcal{C}', \mathcal{W}')$ be another morphism of localizers. A 2-morphism of localizers from F to G is simply a natural transformation $\alpha : F \Rightarrow G$. The universal property of the localization implies that there exists a unique natural transformation

$$\operatorname{ho}(\mathcal{C}) \underbrace{ \bigcup_{\overline{\alpha}}^{\overline{F}}}_{\overline{G}} \operatorname{ho}(\mathcal{C}')$$

such that the 2-diagram



is commutative in an obvious sense.

Remark 2.1.4. Since we always consider that for every localizer $(\mathcal{C}, \mathcal{W})$ the categories \mathcal{C} and ho (\mathcal{C}) have the same class of objects and the localization functor is the identity on objects, it follows that for a morphism of localizers $F: (\mathcal{C}, \mathcal{W}) \to (\mathcal{C}', \mathcal{W}')$, we tautologically have

$$\overline{F}(X) = F(X)$$

for every object X of C.

2.1.5. Let $(\mathcal{C}, \mathcal{W})$ and $(\mathcal{C}', \mathcal{W}')$ be two localizers. A functor $F : \mathcal{C} \to \mathcal{C}'$ is *totally left derivable* when there exists a functor

$$\mathbb{L}F : \mathrm{ho}(\mathcal{C}) \to \mathrm{ho}(\mathcal{C}')$$

and a natural transformation

$$\alpha: \mathbb{L}F \circ \gamma \Rightarrow \gamma' \circ F$$

that makes $\mathbb{L}F$ the *right* Kan extension of $\gamma' \circ F$ along γ :

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{C}' \\ & & & & \downarrow \gamma' \\ & & & & \downarrow \gamma' \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\$$

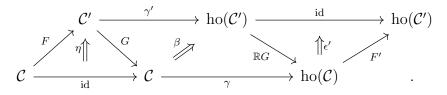
When this right Kan extension is *absolute*, we say that F is *absolutely totally left derivable*. When a functor F is totally left derivable, the pair $(\mathbb{L}F, \alpha)$ is unique up to a unique natural isomorphism and is referred to as *the total left derived functor of* F. Often we will abusively discard α and simply refer to $\mathbb{L}F$ as the total left derived functor of F.

The notion of *total right derivable functor* is defined dually and denoted by $\mathbb{R}F$ when it exists.

Example 2.1.6. Let $F : (\mathcal{C}, \mathcal{W}) \to (\mathcal{C}', \mathcal{W}')$ a morphism of localizers. The universal property of the localization implies that F is absolutely totally left and right derivable and $\mathbb{L}F \simeq \mathbb{R}F \simeq \overline{F}$.

To end this section, we recall a derivability criterion due to Gonzalez, which we shall use in the sequel.

2.1.7. Let $(\mathcal{C}, \mathcal{W})$ and $(\mathcal{C}', \mathcal{W}')$ be two localizers and let $F : \mathcal{C} \longleftrightarrow \mathcal{C}' : G$ be an adjunction whose unit is denoted by η . Suppose that the functor G is totally right derivable with $(\mathbb{R}G, \beta)$ its total right derived functor and suppose that $\mathbb{R}G$ has a left adjoint $F' : \operatorname{ho}(\mathcal{C}) \to \operatorname{ho}(\mathcal{C}')$; the co-unit of this last adjunction being denoted by ϵ' . All this data induces a natural transformation $\alpha : F' \circ \gamma \Rightarrow \gamma' \circ F$ defined as the following composition



Proposition 2.1.8 ([Gon12, Theorem 3.1]). Let (C, W) and (C', W') be two localizers and

$$F: \mathcal{C} \longrightarrow \mathcal{C}': G$$

be an adjunction. If G is absolutely totally right derivable with $(\mathbb{R}G, \beta)$ its left derived functor and if $\mathbb{R}G$ has a left adjoint F'

$$F': \operatorname{ho}(\mathcal{C}) \rightleftharpoons \operatorname{ho}(\mathcal{C}'): \mathbb{R}G_{\mathcal{C}}$$

then F is absolutely totally left derivable and the pair (F', α) , with α defined as in the previous paragraph, is its left derived functor.

2.2 (OP-)DERIVATORS AND HOMOTOPY COLIMITS

Notation 2.2.1. We denote by <u>Cat</u> the 2-category of small categories and <u>CAT</u> the 2-category of big categories. For a 2-category <u>A</u>, the 2-category obtained from <u>A</u> by switching the source and targets of 1-cells is denoted by <u>A</u>^{op}.

The terminal category, i.e. the category with only one object and no non-trivial arrows, is canonically denoted by e. For a (small) category A, the unique functor from A to e is denoted by

$$p_A: A \to e.$$

Definition 2.2.2. An op-prederivator is a (strict) 2-functor

$$\mathbb{D}: \underline{\operatorname{Cat}}^{\operatorname{op}} \to \underline{\operatorname{CAT}}$$

More explicitly, an op-prederivator consists of the data of:

- a big category $\mathbb{D}(A)$ for every small category A,
- a functor $u^* : \mathbb{D}(B) \to \mathbb{D}(A)$ for every functor $u : A \to B$ between small categories,
- a natural transformation

$$\mathbb{D}(B) \underbrace{ \bigvee_{\alpha^*}^{u^*}}_{v^*} \mathbb{D}(A)$$

for every natural transformation

$$A \underbrace{\overset{u}{\underbrace{\forall \alpha}}}_{v} B$$

with A and B small categories,

compatible with compositions and units (in a strict sense).

Remark 2.2.3. Note that some authors call *prederivator* what we have called *op-prederivator*. The terminology we chose in the above definition is compatible with the original one of Grothendieck, who called *prederivator* a 2-functor from <u>Cat</u> to <u>CAT</u> that is contravariant at the level of 1-cells *and* at the level of 2-cells.

Example 2.2.4. Let C be a category. For a small category A, we use the notation C(A) for the category $\underline{\text{Hom}}(A, C)$ of functors $A \to C$ and natural transformations between them. The correspondence $A \mapsto C(A)$ is 2-functorial in an obvious sense and thus defines an op-prederivator

$$\mathcal{C}: \underline{\mathbf{Cat}}^{\mathrm{op}} \to \underline{\mathbf{CAT}}$$
$$A \mapsto \mathcal{C}(A)$$

which we call the op-prederivator represented by C. For $u : A \to B$ in <u>Cat</u>,

$$u^* : \mathcal{C}(B) \to \mathcal{C}(A)$$

is simply the functor induced from u by pre-composition.

We now turn to the most important way of obtaining op-prederivators.

2.2.5. Let $(\mathcal{C}, \mathcal{W})$ be a localizer. For every small category A, we write \mathcal{W}_A the class of *pointwise weak equivalences* of the category $\mathcal{C}(A)$, i.e. the class of arrows $\alpha : d \to d'$ of $\mathcal{C}(A)$ such that $\alpha_a : d(a) \to d'(a)$ belongs to \mathcal{W} for every $a \in Ob(A)$. This defines

a localizer $(\mathcal{C}(A), \mathcal{W}_A)$. The correspondence $A \mapsto (\mathcal{C}(A), \mathcal{W}_A)$ is 2-functorial in that every $u : A \to B$ induces by pre-composition a morphism of localizers

$$u^*: (\mathcal{C}(B), \mathcal{W}_B) \to (\mathcal{C}(A), \mathcal{W}_A)$$

and every $A \underbrace{ \bigvee_{v}}_{v} B$ induces by pre-composition a 2-morphism of localizers

$$(\mathcal{C}(B), \mathcal{W}_B) = \bigvee_{v^*}^{u^*} (\mathcal{C}(A), \mathcal{W}_A).$$

(This last property is trivial since a 2-morphism of localizers is simply a natural transformation between the underlying functors.) Then, by the universal property of the localization, every morphism $u : A \to B$ of Cat induces a functor, again denoted by u^* ,

$$u^* : \operatorname{ho}(\mathcal{C}(B)) \to \operatorname{ho}(\mathcal{C}(A))$$

and every natural transformation $A \underbrace{\stackrel{u}{\stackrel{}}}_{v} A^*$ induces a natural transformation, again denoted by α^* ,

$$\operatorname{ho}(\mathcal{C}(B)) \underbrace{ \bigvee_{\alpha^*}^{u^*} \operatorname{ho}(\mathcal{C}(A))}_{v^*}.$$

Altogether, this defines an op-prederivator

$$\mathcal{H}o^{\mathcal{W}}(\mathcal{C}): \underline{\mathbf{Cat}}^{\mathrm{op}} \to \underline{\mathbf{CAT}}\\ A \mapsto \mathrm{ho}(\mathcal{C}(A)).$$

which we call the *homotopy op-prederivator of* $(\mathcal{C}, \mathcal{W})$. When there is no risk of confusion we will simply write $\mathcal{H}_0(\mathcal{C})$ instead of $\mathcal{H}_0^{\mathcal{W}}(\mathcal{C})$. All the op-prederivators we shall work with arise this way. Notice that for the terminal category *e*, we have a canonical isomorphism

$$\mathcal{H}o(\mathcal{C})(e) \simeq ho(\mathcal{C}),$$

which we shall use without further reference.

Definition 2.2.6. An op-prederivator \mathbb{D} has *left Kan extensions* if for every $u : A \to B$ in Cat, the functor $u^* : \mathbb{D}(B) \to \mathbb{D}(A)$ has a left adjoint

$$u_!: \mathbb{D}(A) \to \mathbb{D}(B).$$

Example 2.2.7. Let C be a category. The op-prederivator represented by C has left Kan extensions if and only if the category C has left Kan extensions along every morphism $u: A \to B$ of Cat in the usual sense. By a standard categorical argument, this means that the op-prederivator represented by C has left Kan extensions if and only if C is cocomplete. Note that for every small category A, the functor

$$p_A^* : \mathcal{C} \simeq \mathcal{C}(e) \to \mathcal{C}(A)$$

is nothing but the diagonal functor that sends an object X of C to the constant diagram with value X. Hence, the functor $p_{A!}$ is nothing but the usual colimit functor of A-shaped diagrams

$$p_{A!} = \operatorname{colim}_A : \mathcal{C}(A) \to \mathcal{C}(e) \simeq \mathcal{C}.$$

2.2.8. We say that a localizer (C, W) has homotopy left Kan extensions when the homotopy op-prederivator of (C, W) has left Kan extensions. In this case, for every small category A, the homotopy colimit functor of A-shaped diagrams is defined as

$$\operatorname{hocolim}_{A} := p_{A!} : \operatorname{ho}(\mathcal{C}(A)) \to \operatorname{ho}(\mathcal{C}).$$

For an object X of $ho(\mathcal{C}(A))$ (which is nothing but a diagram $X : A \to \mathcal{C}$ seen "up to weak equivalence"), the object of $ho(\mathcal{C})$

$$\operatorname{hocolim}_{A}(X)$$

is the *homotopy colimit of X*. For consistency, we also use the notation

$$\operatorname{hocolim}_{a \in A} X(a).$$

When C is also cocomplete (which will always be the case in practice), it follows from Remark 2.1.6 and Proposition 2.1.8 that the functor

$$\operatorname{colim}_A : \mathcal{C}(A) \to \mathcal{C}$$

is left derivable and $hocolim_A$ is the left derived functor of $colim_A$:

$$\mathbb{L}\operatorname{colim}_A \simeq \operatorname{hocolim}_A$$

In particular, for every A-shaped diagram $X : A \to C$, there is a canonical morphism of ho(C)

$$\operatorname{hocolim}_{A}(X) \to \operatorname{colim}_{A}(X).$$

This canonical morphism will be of great importance in the sequel.

2.2.9. Let

$$\begin{array}{ccc} A & \stackrel{f}{\longrightarrow} B \\ \downarrow & \swarrow & \downarrow v \\ C & \stackrel{q}{\longrightarrow} D \end{array}$$

be a 2-square in <u>Cat</u>. Every op-prederivator \mathbb{D} induces a 2-square:

$$\mathbb{D}(A) \xleftarrow{f^*} \mathbb{D}(B)$$

$$u^* \uparrow \swarrow v^*$$

$$\mathbb{D}(C) \xleftarrow{g^*} \mathbb{D}(D).$$

If \mathbb{D} has left Kan extensions, we obtain a canonical natural transformation

$$u_! f^* \Rightarrow g^* v_!$$

referred to as the *homological base change morphism induced by* α and defined as the following composition:

In particular, let $u : A \to B$ be a morphism of <u>Cat</u> and b an object of B seen as a morphism $b : e \to B$. We have a square

$$\begin{array}{ccc} A/b & \stackrel{k}{\longrightarrow} & A \\ p & & & \downarrow \\ p & & & \downarrow \\ e & \stackrel{k}{\longrightarrow} & B \end{array}$$

where :

- A/b is the category whose objects are pairs (a, f : u(a) → b) with a an object of A and f an arrow of B, and morphisms (a, f) → (a', f') are arrows g : a → a' of A such that f' ∘ u(g) = f,
- $k: A/b \to A$ is the functor $(a, p) \mapsto a$,
- ϕ is the natural transformation defined by $\phi_{(a,f)} := f : u(a) \to b$.

Hence, we have a homological base change morphism:

$$p_! k^* \Rightarrow b^* u_!.$$

In the case that \mathbb{D} is the homotopy op-prederivator of a localizer $(\mathcal{C}, \mathcal{W})$, for every object X of $\mathbb{D}(A)$ the above morphism reads

$$\operatorname{hocolim}_{A/b}(X|_{A/b}) \to u_!(X)_b$$

where we use the notation $X|_{A/b}$ for $k^*(X)$ and $u_!(F)_b$ for $b^*(u_!(X))$. Note that this morphism is reminiscent of the formula that computes pointwise left Kan extensions in the "classical" sense (see for example [ML, chapter X, section 3]).

Definition 2.2.10 (Grothendieck). A *right op-derivator* is an op-prederivator \mathbb{D} such that the following axioms are satisfied:

Der 1) For every finite family $(A_i)_{i \in I}$ of small categories, the canonical functor

$$\mathbb{D}(\amalg_{i\in I}A_i)\to\prod_{i\in I}\mathbb{D}(A_i)$$

is an equivalence of categories. In particular, $\mathbb{D}(\emptyset)$ is equivalent to the terminal category.

Der 2) For every small category *A*, the functor

$$\mathbb{D}(A) \to \prod_{a \in \mathrm{Ob}(A)} \mathbb{D}(e)$$

induced by the functors $a^* : \mathbb{D}(A) \to \mathbb{D}(e)$ for all $a \in Ob(A)$ (seen as morphisms $a : e \to A$), is conservative.

- **Der 3d**) \mathbb{D} admits left Kan extensions.
- **Der 4d**) For every $u : A \to B$ in <u>Cat</u> and b object of B, the homological base change morphism

$$p_! k^* \Rightarrow b^* u_!$$

induced by the square

$$\begin{array}{ccc} A/b \xrightarrow{k} A \\ p \downarrow & \swarrow \phi & \downarrow^{i} \\ e \xrightarrow{b} & B \end{array}$$

is an isomorphism.

2.2.11. Let us comment each of the axioms of the previous definition. Axiom **Der 1** ensures that $\mathbb{D}(A)$ "looks like" a category of A-shaped diagrams. Axiom **Der 2** says that isomorphisms in $\mathbb{D}(A)$ can be tested "pointwise". We have already seen that axiom **Der 3d** ensures the existence of left Kan extensions. Finally, axiom **Der 4d** intuitively says that "Kan extensions are computed pointwise".

Example 2.2.12. Let C be a category. The op-prederivator represented by C always satisfy axioms **Der 1** and **Der 2**. We have already seen that axioms **Der 3d** means exactly that C admits left Kan extensions in the classical sense, in which case axiom **Der 4d** is automatically satisfied. Hence, the op-prederivator represented by C is a right op-prederivator if and only if C is cocomplete.

Remark 2.2.13. Beware not to generalize the previous example too hastily. It is not true in general that axiom **Der 3d** implies axiom **Der 4d**; even in the case of the homotopy op-prederivator of a localizer.

This motivates the following definition.

Definition 2.2.14. A localizer (C, W) is *homotopy cocomplete* if the op-prederivator $\mathcal{H}_{O}(C)$ is a right op-derivator.

2.2.15. Axioms **Der 3d** and **Der 4d** can be dualized to obtain axioms **Der 3g** and **Der 4g**, which informally say that the op-prederivator has right Kan extensions and that they are computed pointwise. An op-prederivator satisfying axioms **Der 1**, **Der 2**, **Der 3g** and **Der 4g** is a *left op-derivator*. In fact, an op-prederivator \mathbb{D} is a left op-derivator if and only if the op-prederivator

$$\frac{\underline{\operatorname{Cat}}}{A \mapsto (\mathbb{D}(A^{\operatorname{op}}))^{\operatorname{op}}}$$

is a right op-prederivator. An op-prederivator which is both a left and right op-derivator is an *op-derivator*. For details, the reader can refer to any of the references on derivators previously cited.

2.3 MORPHISMS OF OP-DERIVATORS, PRESERVATION OF HOMOTOPY COLIMITS

We refer to [Lei98] for the precise definitions of pseudo-natural transformation (called strong transformation there) and modification.

2.3.1. Let \mathbb{D} and \mathbb{D}' be two op-prederivators. A *morphism of op-prederivators* $F : \mathbb{D} \to \mathbb{D}'$ is a pseudo-natural transformation from \mathbb{D} to \mathbb{D}' . This means that F consists of:

- a functor $F_A : \mathbb{D}(A) \to \mathbb{D}'(A)$ for every small category A,
- an isomorphism of functors $F_u: F_A u^* \xrightarrow{\sim} u^* F_B$,

$$\begin{split} \mathbb{D}(B) & \xrightarrow{F_B} \mathbb{D}'(B) \\ & \downarrow^{u^*} & \swarrow^{-} & \downarrow^{u^*} \\ \mathbb{D}(A) & \xrightarrow{F_A} \mathbb{D}'(A), \end{split}$$

for every $u: A \to B$ in <u>Cat</u>.

These data being compatible with compositions and units. The morphism is *strict* when F_u is an identity for every $u : A \to B$.

Let $F : \mathbb{D} \to \mathbb{D}'$ and $G : \mathbb{D} \to \mathbb{D}'$ be morphisms of op-prederivators. A 2-morphism $\phi : F \Rightarrow G$ is a modification from F to G. This means that F consists of a natural transformation $\phi_A : F_A \Rightarrow G_A$ for every small category A, and is subject to a coherence axiom similar to the one for natural transformations.

We denote by <u>Pder</u> the 2-category of op-prederivators, morphisms of op-prederivators and 2-morphisms of op-prederivators.

Example 2.3.2. Let $F : \mathcal{C} \to \mathcal{C}'$ be a functor. It induces a strict morphism at the level of op-prederivators, again denoted by F, where for every small category A, the functor $F_A : \mathcal{C}(A) \to \mathcal{C}'(A)$ is induced by post-composition. Similarly, every natural transformation induces a 2-morphism at the level of represented op-prederivators.

Example 2.3.3. Let $F : (\mathcal{C}, \mathcal{W}) \to (\mathcal{C}', \mathcal{W}')$ be a morphism of localizers. For every small category A, the functor $F_A : \mathcal{C}(A) \to \mathcal{C}'(A)$ preserves weak equivalences and the universal property of the localization yields a functor

$$\overline{F}_A : \operatorname{ho}(\mathcal{C}(A)) \to \operatorname{ho}(\mathcal{C}'(A)).$$

This defines a strict morphism of op-prederivators:

$$\overline{F}: \mathcal{H}o(\mathcal{C}) \to \mathcal{H}o(\mathcal{C}').$$

Similarly, every 2-morphism of localizers

$$(\mathcal{C}, \mathcal{W}) \underbrace{\bigvee_{G}^{F}}_{G} (\mathcal{C}', \mathcal{W}')$$

induces a 2-morphism $\overline{\alpha}: \overline{F} \Rightarrow \overline{G}$. Altogether, we have defined a 2-functor

$$\frac{\mathbf{Loc}}{\mathcal{C},\mathcal{W}} \to \frac{\mathbf{Pder}}{\mathcal{H}o(\mathcal{C})},$$

where <u>Loc</u> is the 2-category of localizers.

2.3.4. Let \mathbb{D} and \mathbb{D}' be op-prederivators that admit left Kan extensions and let $F : \mathbb{D} \to \mathbb{D}'$ be a morphism of op-prederivators. For every $u : A \to B$, there is a canonical natural transformation

$$u_! F_A \Rightarrow F_B u_!$$

defined as

$$\mathbb{D}(A) \xrightarrow{u_{!}} \mathbb{D}(B) \xrightarrow{F_{B}} \mathbb{D}'(B)$$

$$\downarrow^{\eta} \qquad \downarrow^{u^{*}} \xrightarrow{\sim}_{F_{u}} \qquad \downarrow^{u^{*}} \xrightarrow{id}$$

$$\mathbb{D}(A) \xrightarrow{}_{F_{A}} \mathbb{D}'(A) \xrightarrow{u_{!}} \mathbb{D}'(B).$$

For example, when \mathbb{D} is the homotopy op-prederivator of a localizer and B is the terminal category e, for every X object of $\mathbb{D}(A)$ the previous canonical morphism reads

$$\operatorname{hocolim}_{A}(F_A(X)) \to F_e(\operatorname{hocolim}_{A}(X)).$$

Definition 2.3.5. Let $F : \mathbb{D} \to \mathbb{D}'$ be a morphism of op-prederivators and suppose that \mathbb{D} and \mathbb{D}' both admit left Kan extensions. We say that F is *cocontinuous*¹ if for every $u : A \to B$, the canonical morphism

$$u_! F_A \Rightarrow F_B u_!$$

is an isomorphism.

Remark 2.3.6. When \mathbb{D} and \mathbb{D}' are homotopy op-prederivators we will often say that a morphism $F : \mathbb{D} \to \mathbb{D}'$ is *homotopy cocontinuous* instead of *cocontinuous* to emphasize the fact that it preserves homotopy Kan extensions.

Example 2.3.7. Let $F : C \to C'$ be a functor and suppose that C and C' are cocomplete. The morphism induced by F at the level of represented op-prederivators is cocontinuous if and only if F is cocontinuous in the usual sense.

2.3.8. As in any 2-category, the notions of equivalence and adjunction make sense in <u>Pder</u>. Precisely, we have that:

- A morphism of op-prederivators F : D → D' is an equivalence when there exists a morphism G : D' → D such that FG is isomorphic to id_{D'} and GF is isomorphic to id_D; the morphism G is a quasi-inverse of F.
- A morphism of op-prederivators F : D → D' is left adjoint to G : D' → D (and G is right adjoint to F) if there exist 2-morphisms η : id_{D'} ⇒ GF and ε : FG ⇒ id_D that satisfy the usual triangle identities.

¹Some authors also say *left exact*.

The following three lemmas are easy 2-categorical routine and are left to the reader.

Lemma 2.3.9. Let $F : \mathbb{D} \to \mathbb{D}'$ be a morphism of op-prederivators. If F is an equivalence then \mathbb{D} is a right op-derivator (resp. left op-derivator, resp. op-derivator) if and only if \mathbb{D}' is one.

Lemma 2.3.10. Let $F : \mathbb{D} \to \mathbb{D}'$ be an equivalence and $G : \mathbb{D}' \to \mathbb{D}$ be a quasi-inverse of G. Then, F is left adjoint to G.

Lemma 2.3.11. Let \mathbb{D} and \mathbb{D}' be op-prederivators that admit left Kan extensions and let $F : \mathbb{D} \to \mathbb{D}'$ be a morphism of op-prederivators. If F is a left adjoint, then it is cocontinuous.

We end this section with a generalization of the notion of localization in the context of op-prederivators.

2.3.12. Let $(\mathcal{C}, \mathcal{W})$ be a localizer. For every small category A, let

$$\gamma_A : \mathcal{C}(A) \to \operatorname{ho}(\mathcal{C}(A))$$

be the localization functor. The correspondence $A \mapsto \gamma_A$ is natural in A and defines a strict morphism of op-prederivators

$$\gamma: \mathcal{C} \to \mathcal{H}o(\mathcal{C}).$$

Definition 2.3.13. Let $(\mathcal{C}, \mathcal{W})$ and $(\mathcal{C}', \mathcal{W}')$ be two localizers and $F : \mathcal{C} \to \mathcal{C}'$ a functor. We say that *F* is *strongly left derivable* if there exists a morphism of op-prederivators

$$\mathbb{L}F:\mathcal{H}o(\mathcal{C})\to\mathcal{H}o(\mathcal{C}')$$

and a 2-morphism of op-prederivators

$$\begin{array}{c} \mathcal{C} & \xrightarrow{F} & \mathcal{C}' \\ \downarrow^{\gamma} & \stackrel{\alpha}{\longrightarrow} & \downarrow^{\gamma'} \\ \mathcal{H}o(\mathcal{C}) & \xrightarrow{\mathbb{L}F} & \mathcal{H}o(\mathcal{C}'). \end{array}$$

such that for every small category A, $((\mathbb{L}F)_A, \alpha_A)$ is the *absolute* total left derived functor of $F_A : \mathcal{C}(A) \to \mathcal{C}'(A)$. The pair $(\mathbb{L}F, \alpha)$ is unique up to a unique isomorphism and is referred to as the *left derived morphism of op-prederivators of* F. Often, we will discard α and simply refer to $\mathbb{L}F$ as the left derived morphism of F. The notion of *strongly right derivable functor* is defined dually and the notation $\mathbb{R}F$ is used. **Example 2.3.14.** Let $(\mathcal{C}, \mathcal{W})$ and $(\mathcal{C}, \mathcal{W}')$ be localizers and $F : \mathcal{C} \to \mathcal{C}'$ a functor. If F preserves weak equivalences (i.e. it is a morphism of localizers), then it is strongly left and right derivable and

$$\overline{F} \simeq \mathbb{L}F \simeq \mathbb{R}G.$$

Gonzalez' criterion (Proposition 2.1.8) admits the following generalization.

Proposition 2.3.15. Let $(\mathcal{C}, \mathcal{W})$ and $(\mathcal{C}', \mathcal{W}')$ be two localizers and

 $F: \mathcal{C} \longleftrightarrow \mathcal{C}': G$

be an adjunction. If G is strongly right derivable and if $\mathbb{R}G$ has a left adjoint F'

$$F': \mathcal{H}o(\mathcal{C}) \longleftrightarrow \mathcal{H}o(\mathcal{C}'): \mathbb{R}G,$$

then F is strongly left derivable and

$$\mathbb{L}F \simeq F'.$$

Proof. Let $\alpha : F' \circ \gamma \Rightarrow F \circ \gamma$ be the 2-morphism of op-prederivators defined *mu*tatis mutandis as in 2.1.7 but at the level of op-prederivators. Proposition 2.1.8 gives us that for every small category A, the functor F_A is absolutely totally left derivable with (F'_A, α_A) its total left derived functor. This means exactly that F' is strongly left derivable and (F', α) is the left derived morphism of op-prederivators of F. \Box

2.4 HOMOTOPY COCARTESIAN SQUARES

2.4.1. Let Δ_1 be the ordered set $\{0 < 1\}$ seen as category. We use the notation \Box for the category $\Delta_1 \times \Delta_1$, which can be pictured as the *commutative* square

$$(0,0) \longrightarrow (0,1)$$
$$\downarrow \qquad \qquad \downarrow$$
$$(1,0) \longrightarrow (1,1)$$

and we use the notation \ulcorner for the full subcategory of \square spanned by (0,0), (0,1) and (1,0), which can be pictured as

$$(0,0) \longrightarrow (0,1)$$

$$\downarrow$$

$$(1,0) \qquad .$$

Finally, we write $i_{\ulcorner}: \ulcorner \rightarrow \Box$ for the canonical inclusion functor.

Definition 2.4.2. Let \mathbb{D} be an op-prederivator. An object X of $\mathbb{D}(\Box)$ is *cocartesian* if for every object Y of $\mathbb{D}(\Box)$, the canonical map

$$\operatorname{Hom}_{\mathbb{D}(\square)}(X,Y) \to \operatorname{Hom}_{\mathbb{D}(\ulcorner)}(i_{\ulcorner}^{*}(X),i_{\ulcorner}^{*}(Y))$$

induced by the functor $i_{r}^{*} : \mathbb{D}(\Box) \to \mathbb{D}(\Gamma)$, is a bijection.

Example 2.4.3. Let C be a category. An object of $C(\Box)$ is nothing but a commutative square in C and it is cocartesian in the sense of the previous definition if and only if it is cocartesian in the usual sense.

For the following definition to make sense, recall that for a localizer $(\mathcal{C}, \mathcal{W})$ and a small category A, the objects of $\mathcal{H}_0(\mathcal{C})(A) = h_0(\mathcal{C}(A))$ are identified with the objects of $\mathcal{C}(A)$ via the localization functor. In particular, an object of $\mathcal{H}_0(\mathcal{C})(\Box)$ is a commutative square of \mathcal{C} (up to weak equivalence).

Definition 2.4.4. Let (C, W) be a localizer. A commutative square of C is said to be *homotopy cocartesian* if it is cocartesian in $H_0(C)$ in the sense of Definition 2.4.2.

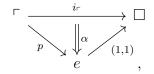
2.4.5. Let \mathbb{D} be an op-prederivator. The object (1, 1) of \Box can be considered as a morphism of Cat

$$(1,1): e \to \square$$

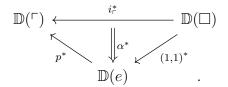
and thus induces a functor $(1,1)^* : \mathbb{D}(\Box) \to \mathbb{D}(e)$. For an object X of $\mathbb{D}(\Box)$, we use the notation

$$X_{(1,1)} := (1,1)^*(X).$$

Now, since (1, 1) is the terminal object of \Box , we have a canonical 2-triangle



where we wrote p instead of p_{\neg} for short and where α is the unique such natural transformation. Hence, we have a 2-triangle



Suppose now that \mathbb{D} has left Kan extensions. For X an object of $\mathbb{D}(\Box)$, we have a canonical morphism $p_!(i^*_{\vdash}(X)) \to X_{(1,1)}$ defined as the composition

$$p_!(i_{\ulcorner}^*(X)) \to p_!p^*(X_{(1,1)}) \to X_{(1,1)},$$

where the arrow on the left is induced by α^* and the arrow on the right is induced by the co-unit of the adjunction $p_! \dashv p^*$.

When \mathbb{D} is the homotopy op-prederivator of a localizer, and X is a commutative square of \mathcal{C}

$$X = \begin{array}{c} A \xrightarrow{u} B \\ \downarrow f & \downarrow g \\ C \xrightarrow{v} D, \end{array}$$

this previous morphism reads

$$\operatorname{hocolim}\left(\begin{array}{c}A \xrightarrow{u} B\\ \star f\\ C\end{array}\right) \to D.$$

Proposition 2.4.6. Let \mathbb{D} be a right op-prederivator. An object X of $\mathbb{D}(\Box)$ is cocartesian if and only if the canonical map $p_!(i^*_{r}(X)) \to X_{(1,1)}$ is an isomorphism.

Proof. Let Y be another object of $\mathbb{D}(\Box)$. Using the adjunction $i_{\lceil !} \dashv i_{\lceil}^*$, the canonical map $\operatorname{Hom}_{\mathbb{D}(\Box)}(X,Y) \to \operatorname{Hom}_{\mathbb{D}(\ulcorner)}(i_{\ulcorner}^*(X), i_{\ulcorner}^*(Y))$ can be identified with

$$\operatorname{Hom}_{\mathbb{D}(\square)}(X,Y) \to \operatorname{Hom}_{\mathbb{D}(\sqcap)}(i_{\sqcap}(X)),Y).$$

Hence, X is cocartesian if and only if the co-unit map $X \to i_{\lceil !}(i_{\lceil}^*(X))$ is an isomorphism. The rest follows then from [Gro16, Lemma 9.2.2(i)].

Hence, for a homotopy cocomplete localizer $(\mathcal{C}, \mathcal{W})$, a commutative square of \mathcal{C} is homotopy cocartesian if and only if the bottom right apex of the square is the homotopy colimit of the upper left corner of the square. This hopefully justifies the terminology of "cocartesian square".

The previous proposition admits the following immediate corollary.

Corollary 2.4.7. Let \mathbb{D} and \mathbb{D}' be right op-derivators and $F : \mathbb{D} \to \mathbb{D}'$ a morphism of op-prederivators. If F is cocontinuous, then it preserves cocartesian squares. This means that if an object X of $\mathbb{D}(\square)$ is cocartesian, then $F_{\square}(X)$ is a cocartesian square of \mathbb{D}' .

We end this section with two useful lemmas which show that homotopy cocartesian squares behave much like "classical" cocartesian squares.

Lemma 2.4.8. Let $(\mathcal{C}, \mathcal{W})$ be a homotopy cocomplete localizer and let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ & & \downarrow^{f} & & \downarrow^{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array}$$

be a commutative square in C. If f and g are weak equivalences then the previous square is homotopy cocartesian.

Proof. Using **Der 2**, one can show that the previous square is isomorphic in $\mathcal{H}_0(\mathcal{C})(\Box)$ to the square

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ \downarrow^{1_A} & \downarrow^{1_B} \\ A & \stackrel{u}{\longrightarrow} & B. \end{array}$$

The result follows then from [Gro13, Proposition 3.12(2)].

Lemma 2.4.9 (Pasting lemma). Let $(\mathcal{C}, \mathcal{W})$ be a homotopy cocomplete localizer and let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B & \stackrel{v}{\longrightarrow} & C \\ \downarrow^{f} & \qquad \downarrow^{g} & \qquad \downarrow^{h} \\ D & \stackrel{w}{\longrightarrow} & E & \stackrel{x}{\longrightarrow} & F \end{array}$$

be a commutative diagram in C. If the square on the left is cocartesian, then the outer square is cocartesian if and only if the right square is cocartesian.

Proof. This is a particular case of [Gro13, Proposition 3.13(1)].

2.5 MODEL CATEGORIES

In this section, we quickly review some aspects of the relation between Quillen's theory of model categories and Grothendieck's theory of derivators. We suppose that the reader is familiar with the former one and refer to the standard textbooks on the subject (such as [Hov07, Hir09, DS95]) for basic definitions and results.

For a model category $\mathcal{M} = (\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$, the homotopy op-prederivator of \mathcal{M} , denoted by $\mathcal{H}_0(\mathcal{M})$, is the homotopy op-prederivator of the localizer $(\mathcal{M}, \mathcal{W})$.

The following theorems are due to Cisinski [Cis03] and can be summed up by the slogan:

Model categories are homotopy cocomplete and left Quillen functors are homotopy cocontinuous.

Theorem 2.5.1 (Cisinski). Let $(\mathcal{M}, \mathcal{W}, Cof, Fib)$ be a model category. The localizer $(\mathcal{M}, \mathcal{W})$ is homotopy cocomplete.

Theorem 2.5.2 (Cisinski). Let \mathcal{M} and \mathcal{M}' be model categories. Let $F \colon \mathcal{M} \to \mathcal{M}'$ be a left Quillen functor (i.e. the left adjoint in a Quillen adjunction). The functor F is strongly left derivable and the morphism of op-prederivators

$$\mathbb{L}F\colon \mathcal{H}o(\mathcal{M})\to \mathcal{H}o(\mathcal{M}')$$

is homotopy cocontinuous.

Remark 2.5.3. The obvious duals of the two above theorems are also true. The reason we put emphasis on cocompleteness rather that completeness is because we will make no use whatsoever of homotopy limits in this dissertation.

Remark 2.5.4. Note that for a model category $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$, its homotopy op-prederivator only depends on its underlying localizer. Hence, the existence of the classes Cof and Fib with the usual properties defining a model structure ought to be thought of as a *property* of the localizer $(\mathcal{M}, \mathcal{W})$, which is sufficient to define a "homotopy theory". For example, Theorem 2.5.1 should have been stated by saying that if a localizer $(\mathcal{M}, \mathcal{W})$ can be extended to a model category $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$, then it is homotopy cocomplete.

Even if Theorem 2.5.1 tells us that (the homotopy op-prederivator of) a model category $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$ has homotopy left Kan extensions, it is not generally true that for a small category A the category of diagrams $\mathcal{M}(A)$ admits a model structure with the pointwise weak equivalences as its weak equivalences. Hence, in general we cannot use the theory of Quillen functors to compute homotopy left Kan extensions (and in particular homotopy colimits). However, in practice all the model categories that we shall encounter are *cofibrantly generated*, in which case the theory is much simpler because $\mathcal{M}(A)$ does admit a model structure with the pointwise weak equivalences as its weak equivalences.

2.5.5. Let C be a category with coproducts and A a small category. For every object X of C and every object a of A, we define $X \otimes a$ as the functor

$$X \otimes a : A \to \mathcal{C}$$
$$b \mapsto \coprod_{\operatorname{Hom}_A(a,b)} X$$

For every object a of A, this gives rise to a functor

$$-\otimes a: \mathcal{C} \to \mathcal{C}(A)$$

 $X \mapsto X \otimes a.$

Proposition 2.5.6. Let $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$ be a cofibrantly generated model category with I (resp. J) as a set of generating cofibrations (resp. trivial cofibrations). For every small category A, there exists a model structure on $\mathcal{M}(A)$ such that:

- the weak equivalences are the pointwise weak equivalences,
- the fibrations are the pointwise fibrations,

- the cofibrations are those morphisms which have the left lifting property to trivial fibrations.

Moreover, this model structure is cofibrantly generated and a set of generating cofibrations (resp. trivial cofibrations) is given by

$$\{f \otimes a : X \otimes a \to Y \otimes a \mid a \in Ob(A), f \in I\}$$

(resp.

$$\{f \otimes a : X \otimes a \to Y \otimes a \mid a \in \mathrm{Ob}(A), f \in J\}\}.$$

Proof. See for example [Cis03, Proposition 3.4].

2.5.7. The model structure of the previous proposition is referred to as the *projective model structure on* $\mathcal{M}(A)$.

Proposition 2.5.8. Let $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$ be a cofibrantly generated model category. For every $u : A \to B$, the adjunction

$$u_!: \mathcal{M}(A) \longleftrightarrow \mathcal{M}(B): u^*$$

is a Quillen adjunction with respect to the projective model structures on $\mathcal{M}(A)$ and on $\mathcal{M}(B)$.

Proof. By definition of the projective model structure, u^* preserve weak equivalences and fibrations.

2.5.9. In particular, in the case that B is the terminal category e, we have that

$$\operatorname{colim}_A: \mathcal{M}(A) \longleftrightarrow \mathcal{M}(e) \simeq \mathcal{M}: p_A^*$$

is a Quillen adjunction. Since $hocolim_A$ is the left derived functor of $colim_A$, we obtain the following immediate corollary of the previous proposition.

Corollary 2.5.10. Let $(\mathcal{M}, \mathcal{W}, \text{Cof}, \text{Fib})$ be a cofibrantly generated model category, A a small category and $X : A \to \mathcal{M}$ a diagram. If X is cofibrant for the projective model structure on $\mathcal{C}(A)$, then the canonical morphism of $\text{ho}(\mathcal{C})$

$$\operatorname{hocolim}_{A}(X) \to \operatorname{colim}_{A}(X)$$

is an isomorphism.

Proof. This is simply a particular case of the general fact that for a left Quillen functor $F : \mathcal{M} \to \mathcal{M}'$ and a cofibrant object X of \mathcal{M} , the canonical map of $ho(\mathcal{M}')$

$$\mathbb{L}F(X) \to F(X)$$

is an isomorphism.

Below is a particular case for which the previous corollary applies.

Proposition 2.5.11. Let $(\mathcal{M}, \mathcal{W}, Cof, Fib)$ be a cofibrantly generated model category and let X be a sequential diagram in \mathcal{M}

$$X_0 \to X_1 \to X_2 \to \cdots$$

(i.e. a diagram $X : (\mathbb{N}, \leq) \to \mathcal{M}$). If X_0 is cofibrant and each $X_i \to X_{i+1}$ is a cofibration, then X is cofibrant for the projective model structure on $\mathcal{M}((\mathbb{N}, \leq))$.

Proof. It is an easy exercise that uses only the fact that cofibrations of the projective model structure are, by definition, the morphisms with left lifting property to pointwise fibrations. For details see [Sch13, Example 2.3.16].

Another setting for which a category of diagrams $\mathcal{M}(A)$ can be equipped with a model structure whose weak equivalences are the pointwise equivalences and for which the A-colimit functor is left Quillen is when the category A is a *Reedy category*. Rather that recalling this theory, we simply put here the only practical result that we shall need in the sequel.

Lemma 2.5.12. Let $(\mathcal{M}, \mathcal{W}, Cof, Fib)$ be a model category and let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ & \downarrow^{f} & & \downarrow^{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array}$$

be a cocartesian square in \mathcal{M} . If either u or f is a cofibration and if A, B and C are cofibrant objects, then this square is homotopy cocartesian.

Proof. See for example [Lur09, Proposition A.2.4.4(i)].

CHAPTER 3

HOMOTOPY THEORY OF ω -CATEGORIES

3.1 NERVE

3.1.1. We denote by Δ the category whose objects are the finite non-empty totally ordered sets $[n] = \{0 < \cdots < n\}$ and whose morphisms are the non-decreasing maps. For n > 0 and $0 \le i \le n$, we denote by

 $\delta^i: [n-1] \to [n]$

the only injective increasing map whose image does not contain i, and for $n \ge 0$ and $0 \le i \le n$, we denote by

$$\sigma^i:[n+1]\to [n]$$

the only surjective non-decreasing map such that the pre-image of $i \in [n]$ contains exactly two elements. The category $\widehat{\Delta}$ of *simplicial sets* is the category of presheaves on Δ . For a simplicial set X, we use the notations

$$X_n := X([n])$$

$$\partial_i := X(\delta^i) : X_n \to X_{n-1}$$

$$s_i := X(\sigma^i) : X_n \to X_{n+1}.$$

Elements of X_n are referred to as *n*-simplices of X, the maps ∂_i are the face maps and the maps s_i are the degeneracy maps.

3.1.2. We denote by $\mathcal{O} : \Delta \to \omega \mathbf{Cat}$ the cosimplicial object introduced by Street in [Str87]. The ω -category \mathcal{O}_n is called the *n*-oriental. There are various ways to give a precise definition of the orientals, but each of them requires some machinery that we don't want to introduce here. Instead, we only recall some important facts on orientals

that we shall need in the sequel and refer to the literature on the subject (such as [Str87, Str91, Str94, Ste04, BG16] or [AM20b, chapitre 7]) for details.

The two main points to retain are:

(OR1) Each \mathcal{O}_n is a free ω -category whose set of generating k-cells is canonically isomorphic to the sets of increasing sequences

$$0 \le i_1 < i_2 < \dots < i_k \le n,$$

or, which is equivalent, to injective increasing maps $[k] \rightarrow [n]$.

We use the notation $\langle i_1 i_2 \cdots i_k \rangle$ for such a cell. In particular, we have that:

- There are no generating k-cells for k > n. Hence, \mathcal{O}_n is an n-category.
- There is exactly one generating *n*-cell of \mathcal{O}_n , which is $\langle 0 \ 1 \cdots n \rangle$. We refer to this cell as the *principal cell of* \mathcal{O}_n .
- There are exactly n + 1 generating (n 1)-cells of \mathcal{O}_n . They correspond to the maps

$$\delta^i: [n-1] \to [n]$$

for $i \in \{0, \dots, n\}$.

(OR2) For n > 0, the source (resp. target) of the principal cell of \mathcal{O}_n can be expressed as a composition of all the generating (n-1)-cells corresponding to δ^i with *i* odd (resp. even); each of these generating (n-1)-cells appearing exactly once in the composite.

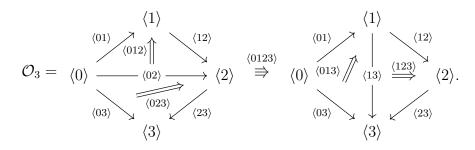
Another way of formulating (**OR2**) is: for n > 0 the weight (1.3.6) of the (n - 1)-cell corresponding to δ_i in the *source* of the principal cell of \mathcal{O}_n is 1 if *i* is odd and 0 if *i* is even and the other way around for the *target* of the principal cell of \mathcal{O}_n . Here are some pictures in low dimension:

$$\mathcal{O}_{0} = \langle 0 \rangle,$$

$$\mathcal{O}_{1} = \langle 0 \rangle \xrightarrow{\langle 01 \rangle} \langle 1 \rangle,$$

$$\mathcal{O}_{2} = \langle 0 \rangle \xrightarrow{\langle 01 \rangle} \langle 012 \rangle \uparrow \uparrow \downarrow \langle 012 \rangle} \langle 12 \rangle,$$

$$\langle 0 \rangle \xrightarrow{\langle 01 \rangle} \langle 012 \rangle \uparrow \downarrow \langle 02 \rangle} \langle 2 \rangle,$$



3.1.3. For every ω -category C, the nerve of C is the simplicial set $N_{\omega}(C)$ defined as

$$N_{\omega}(C) : \Delta^{\mathrm{op}} \to \mathbf{Set}$$

 $[n] \mapsto \mathrm{Hom}_{\omega \mathbf{Cat}}(\mathcal{O}_n, C)$

By post-composition, this yields a functor

$$N_{\omega}: \omega \mathbf{Cat} \to \widehat{\Delta}$$
$$C \mapsto N_{\omega}(C),$$

which we refer to as the *nerve functor for* ω -categories. Furthermore, for every $n \in \mathbb{N}$, we also define a nerve functor for *n*-categories as the restriction of N_{ω} to *n*Cat (seen as a full subcategory of ω Cat)

$$N_n := N_\omega \big|_{n \operatorname{Cat}} : n \operatorname{Cat} \to \widehat{\Delta}.$$

By the usual Kan extension technique, we obtain for every $n \in \mathbb{N} \cup \{\omega\}$ a functor

$$c_n:\widehat{\Delta}\to n\mathbf{Cat}$$

left adjoint to N_n .

3.1.4. For n = 1, the functor N_1 is the usual nerve of categories. Recall that for a (small) category C, an m-simplex X of $N_1(C)$ is a sequence of composable arrows of C

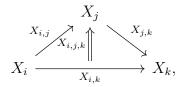
$$X_0 \xrightarrow{X_{0,1}} X_1 \xrightarrow{X_{1,2}} \cdots \longrightarrow X_{m-1} \xrightarrow{X_{m-1,m}} X_m$$

For m > 0 and $0 \le i \le m$, the (m - 1)-simplex $\partial_i(X)$ is obtained by composing arrows at X_i (or simply deleting it for i = 0 or m). For $m \ge 0$ and $0 \le i \le m$, the (m + 1)-simplex $s_i(X)$ is obtained by inserting a unit map at X_i .

For n = 2, the functor N_2 is what is sometimes known as the *Duskin nerve* [Dus02] (restricted from bicategories to 2-categories). For a 2-category C, an *m*-simplex X of $N_2(C)$ consists of:

- for every $0 \le i \le m$, an object X_i of C,

- for all $0 \le i \le j \le m$, an arrow $X_{i,j} : X_i \to X_j$ of C,
- for all $0 \le i \le j \le k \le m$, a 2-triangle



subject to the following axioms:

- for all $0 \le i \le m$, we have

$$X_{i,i} = 1_{X_i},$$

- for all $0 \le i \le j \le m$, we have

$$X_{i,i,j} = X_{i,j,j} = 1_{X_{i,j}},$$

- for all $0 \le i < j < k < l \le m$, we have the equality (known as the *cocycle condition*)

$$(X_{k,l} * X_{i,j,k}) * X_{i,k,l} = (X_{j,k,l} * X_{i,j}) * X_{i,j,l}.$$

For m > 0 and $0 \le l \le m$, the (m - 1)-simplex $\partial_l(X)$ is defined as

$$\partial_l(X)_i = X_{\delta_l(j)}, \quad \partial_l(X)_{i,j} = X_{\delta_l(i),\delta_l(j)} \text{ and } \partial_l(X)_{i,j,k} = X_{\delta_l(i),\delta_l(j),\delta_l(k)}.$$

And similarly, for $m \ge 0$ and $0 \le l \le m$, the (m + 1)-simplex $s_l(X)$ is defined as

$$s_l(X)_i = X_{\sigma_l(j)}, \quad s_l(X)_{i,j} = X_{\sigma_l(i),\sigma_l(j)} \text{ and } s_l(X)_{i,j,k} = X_{\sigma_l(i),\sigma_l(j),\sigma_l(k)}.$$

3.2 THOMASON EQUIVALENCES

3.2.1. From now on, we will consider that the category $\widehat{\Delta}$ is equipped with the model structure defined by Quillen in [Qui67]. A *weak equivalence of simplicial sets* is a weak equivalence for this model structure. The cofibrations for this model structure are the monomorphisms.

Definition 3.2.2. Let $n \in \mathbb{N} \cup \{\omega\}$. A morphism $f : X \to Y$ of *n*Cat is a *Thomason* equivalence when $N_n(f) : N_n(X) \to N_n(Y)$ is a weak equivalence of simplicial sets. We denote by $\mathcal{W}_n^{\text{Th}}$ the class of Thomason equivalences.

3.2.3. We usually make reference to the name "Thomason" in the notations of homotopic constructions induced by Thomason equivalences. For example, we write $\mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}})$ for the homotopy op-prederivator of $(n\mathbf{Cat}, \mathcal{W}_n^{\mathrm{Th}})$ and

$$\gamma^{\mathrm{Th}}: n\mathbf{Cat} \to \mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}})$$

for the localization morphism. The reason is to avoid confusion with other weak equivalences on nCat that we will introduce later.

3.2.4. By definition, the nerve functor induces a morphism of localizers

$$N_n: (n\mathbf{Cat}, \mathcal{W}_n^{\mathrm{Th}}) \to (\widehat{\Delta}, \mathcal{W}_{\Delta})$$

and hence a morphism of op-prederivators

$$\overline{N_n}: \mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta}).$$

Theorem 3.2.5 (Gagna). For every $1 \le n \le \omega$, the morphism

$$\overline{N}_n: \mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta})$$

is an equivalence of op-prederivators.

Proof. Recall from 3.1.3 that $c_n : \widehat{\Delta} \to n\mathbf{Cat}$ denotes the left adjoint of the nerve functor N_n . In [Gag18], Gagna proves that there exists a functor $Q : \widehat{\Delta} \to \widehat{\Delta}$, as well as a zigzag of morphisms of functors

$$N_n c_n Q \stackrel{\alpha}{\longleftarrow} Q \stackrel{\gamma}{\longrightarrow} \mathrm{id}_{\widehat{\Lambda}}$$

and a morphism of functors

$$c_n Q N_n \xrightarrow{\beta} \operatorname{id}_{\widehat{\Delta}},$$

such that $c_n Q$ preserves weak equivalences and α , β and γ are weak equivalences argument by argument. This easily implies that

$$\overline{c_n Q} : \mathcal{H}o(\widehat{\Delta}) \to \mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}})$$

is a quasi-inverse (2.3.8) of

$$\overline{N_n}: \mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta}).$$

From Lemma 2.3.9, we obtain the following corollary.

Corollary 3.2.6. For every $1 \le n \le \omega$, the localizer $(n\mathbf{Cat}^{\mathrm{Th}}, \mathcal{W}_n^{\mathrm{Th}})$ is homotopy cocomplete (Definition 2.2.14).

We will speak of "Thomason homotopy colimits" and "Thomason homotopy cocartesian squares" for homotopy colimits and homotopy cocartesian squares in the localizer $(n\mathbf{Cat}^{\mathrm{Th}}, \mathcal{W}_n^{\mathrm{Th}})$. (See also 3.2.10 below.)

Another consequence of Gagna's theorem is the following corollary.

Corollary 3.2.7. For every $1 \le n \le \omega$, the class $\mathcal{W}_n^{\text{Th}}$ is saturated (2.1.1).

Proof. This follows immediately from the fact that $\overline{N_n}$: ho($n\mathbf{Cat}^{\mathrm{Th}}$) \rightarrow ho($\widehat{\Delta}$) is an equivalence of categories and the fact that weak equivalences of simplicial sets are saturated (because they are the weak equivalences of a model structure).

Remark 3.2.8. Corollaries 3.2.6 and 3.2.7 would also follow from the existence of a model structure on nCat with W_n^{Th} as the weak equivalences. For n = 1, this was established by Thomason [Tho80], and for n = 2, by Ara and Maltsiniotis [AM14]. For n > 3, the existence of such a model structure is conjectured but not yet established.

By definition, for all $1 \le n \le m \le \omega$, the canonical inclusion

$$n\mathbf{Cat} \hookrightarrow m\mathbf{Cat}$$

sends the Thomason equivalences of nCat to Thomason equivalences of mCat. Hence, it induces a morphism of localizers and then a morphism of op-prederivators

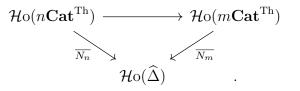
$$\mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(m\mathbf{Cat}^{\mathrm{Th}}).$$

Proposition 3.2.9. For all $1 \le n \le \omega$, the canonical morphism

$$\mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(m\mathbf{Cat}^{\mathrm{Th}})$$

is an equivalence of op-prederivators.

Proof. This follows from Theorem 3.2.5 and the commutativity of the triangle



3.2.10. It follows from the previous proposition that for all $1 \le n \le m \le \omega$, the morphism $\mathcal{H}o(n\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(m\mathbf{Cat}^{\mathrm{Th}})$ of op-prederivators is homotopy cocontinuous and reflects homotopy colimits (in an obvious sense). Hence, given a diagram $d: I \to n\mathbf{Cat}$ with n > 0, we can harmlessly use the notation

$$\operatorname{hocolim}_{i \in I}^{\operatorname{Th}}(d)$$

for both the Thomason homotopy colimits in nCat and in ωCat (or any mCat with $n \leq m$). Similarly, a commutative square of nCat is Thomason homotopy cocartesian in nCat if and only if it is so in ωCat . Hence, there is really no ambiguity when simply calling such a square *Thomason homotopy cocartesian*.

3.3 TENSOR PRODUCT AND OPLAX TRANSFORMATIONS

Recall that ω Cat can be equipped with a monoidal product \otimes , introduced by Al-Agl and Steiner in [AAS93] and by Crans in [Cra95], commonly referred to as the *Gray tensor product*. The implicit reference for this section is [AM20b, Appendices A and B].

3.3.1. The Gray tensor product makes ω Cat into a monoidal category for which the unit is the ω -category \mathbb{D}_0 (which is the terminal ω -category). This monoidal category is *not* symmetric but it is biclosed [AM20b, Theorem A.15], meaning that there exist two functors

$$\underline{\mathrm{hom}}_{\mathrm{oplax}}(-,-), \underline{\mathrm{hom}}_{\mathrm{lax}}(-,-): \omega \mathbf{Cat}^{\mathrm{op}} \times \omega \mathbf{Cat} \to \omega \mathbf{Cat}$$

such that for all ω -categories X, Y and Z, we have isomorphisms

$$\operatorname{Hom}_{\omega \mathbf{Cat}}(X \otimes Y, Z) \simeq \operatorname{Hom}_{\omega \mathbf{Cat}}(X, \underline{\operatorname{hom}}_{\operatorname{oplax}}(Y, Z))$$
$$\simeq \operatorname{Hom}_{\omega \mathbf{Cat}}(Y, \underline{\operatorname{hom}}_{\operatorname{lax}}(X, Z))$$

natural in X, Y and Z. When $X = \mathbb{D}_0$, we have $\mathbb{D}_0 \otimes Y \simeq Y$, and thus

 $\operatorname{Hom}_{\omega \operatorname{Cat}}(Y, Z) \simeq \operatorname{Hom}_{\omega \operatorname{Cat}}(\mathbb{D}_0, \underline{\operatorname{hom}}_{\operatorname{oplax}}(Y, Z)).$

Hence, the 0-cells of the ω -category $\underline{\hom}_{oplax}(Y, Z)$ are the ω -functors $Y \to Z$.

3.3.2. Let $u, v : X \to Y$ be two ω -functors. An *oplax transformation* from u to v is a 1-cell α of $\underline{\hom}_{oplax}(X, Y)$ with source u and target v. We usually use the double arrow notation

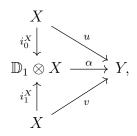
$$\alpha: u \Rightarrow v$$

for oplax transformations. By adjunction, we have

$$\operatorname{Hom}_{\omega \mathbf{Cat}}(\mathbb{D}_1, \underline{\operatorname{hom}}_{\operatorname{oplax}}(X, Y)) \simeq \operatorname{Hom}_{\omega \mathbf{Cat}}(\mathbb{D}_1 \otimes X, Y)$$
$$\simeq \operatorname{Hom}_{\omega \mathbf{Cat}}(X, \underline{\operatorname{hom}}_{\operatorname{lax}}(\mathbb{D}_1, Y)).$$

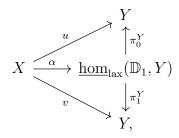
Hence, $\alpha : u \Rightarrow v$ can be encoded in the following two ways:

- As an ω -functor $\alpha : \mathbb{D}_1 \otimes X \to Y$ such that the following diagram



where i_0^X and i_1^X are induced by the two ω -functors $\mathbb{D}_0 \to \mathbb{D}_1$ and where we implicitly used the isomorphism $\mathbb{D}_0 \otimes X \simeq X$, is commutative.

- As an ω -functor $\alpha: X \to \underline{\hom}_{lax}(\mathbb{D}_1, Y)$ such that the following diagram



where π_0^Y and π_1^Y are induced by the two ω -functors $\mathbb{D}_0 \to \mathbb{D}_1$ and where we implicitly used the isomorphism $\underline{\mathrm{hom}}_{\mathrm{lax}}(\mathbb{D}_0, Y) \simeq Y$, is commutative.

The ω -category $\underline{\hom}_{lax}(\mathbb{D}_1, Y)$ is sometimes referred to as the ω -category of cylinders in Y. An explicit description of this ω -category can be found, for example, in [Mét03, Appendix A], [LM09, Section 4] or [AM20b, Appendice B.1].

3.3.3. [Formulas for oplax transformations] We now give a third way of describing oplax transformations based on explicit formulas. The proof that this description is equivalent to those given in the previous paragraph can be found in [AM20b, Appendice B.2].

Let $u, v : X \to Y$ two ω -functors. An oplax transformation $\alpha : u \Rightarrow v$ is given by the data of:

- for every 0-cell x of X, a 1-cell of Y

$$\alpha_x: u(x) \to v(x),$$

- for every *n*-cell of x of X with n > 0, an (n + 1)-cell of Y

$$\alpha_x : \alpha_{t_{n-1}(x)} * \cdots * \alpha_{t_0(x)} * u(x) \to v(x) * \alpha_{s_0(x)} * \cdots * \alpha_{s_{n-1}(x)}$$

subject to the following axioms:

1. for every *n*-cell x of X,

$$\alpha_{1_x} = 1_{\alpha_x}$$

2. for all $0 \le k < n$, for all *n*-cells *x* and *y* of *X* that are *k*-composable,

$$\alpha_{\substack{x \neq y \\ k}} = \left(v(\mathbf{t}_{k+1}(x)) * \alpha_{\mathbf{s}_0(y)} * \cdots * \alpha_{\mathbf{s}_{k-1}(y)} * \alpha_y \right) \\ * \left(\alpha_x * \alpha_{\mathbf{t}_{k-1}(x)} * \cdots * \alpha_{\mathbf{t}_0(x)} * u(\mathbf{s}_{k+1}(y)) \right).$$

Note that to read the formulas correctly, one has to remember the convention that for k < n, the composition $\underset{k}{*}$ has priority over $\underset{n}{*}$ (see 1.1.2).

Example 3.3.4. When C and D are n-categories with n finite and $u, v : C \to D$ are two n-functors, an oplax transformation $\alpha : u \Rightarrow v$ amounts to the data of a (k+1)-cell α_x of D for each k-cell x of C with $0 \le k \le n$, with source and target as in the previous paragraph. These data being subject to the axioms of the previous paragraph. Note that when x is an n-cell of C, α_x is necessarily a unit, which can be expressed as the equality

$$\alpha_{t_{n-1}(x)} \underset{n-1}{*} \cdots \underset{1}{*} \alpha_{t_0(x)} \underset{0}{*} u(x) = v(x) \underset{0}{*} \alpha_{s_0(x)} \underset{1}{*} \cdots \underset{n-1}{*} \alpha_{s_{n-1}(x)}$$

In particular, when n = 1 and C and D are thus (small) categories, an oplax transformation $u \Rightarrow v$ is nothing but a natural transformation from u to v.

3.3.5. Let $u: C \to D$ be an ω -functor. There is an oplax transformation from u to u, denoted by 1_u , which is defined as

$$(1_u)_x := 1_{u(x)}$$

for every cell x of C. More abstractly, this oplax transformation corresponds to the ω -functor

$$\mathbb{D}_1 \otimes C \xrightarrow{p \otimes u} \mathbb{D}_0 \otimes D \simeq D,$$

where p is the only ω -functor $\mathbb{D}_1 \to \mathbb{D}_0$.

3.3.6. Let

$$B \xrightarrow{f} C \xrightarrow{u} D \xrightarrow{g} E$$

be a diagram in ω Cat and $\alpha : u \Rightarrow v$ an oplax transformation. The data of

$$(g \star \alpha)_x := g(\alpha_x)$$

for each cell x of C (resp.

$$(\alpha \star f)_x := \alpha_{f(x)}$$

for each cell x of B) defines an oplax transformation from $g \circ u$ to $g \circ v$ (resp. $u \circ f$ to $v \circ f$) which we denote by $g \star \alpha$ (resp. $\alpha \star f$).

More abstractly, if α is seen as an ω -functor $\mathbb{D}_1 \otimes C \to D$, then $g \star \alpha$ (resp. $\alpha \star f$) corresponds to the ω -functor obtained as the following composition

$$\mathbb{D}_1 \otimes C \xrightarrow{\alpha} D \xrightarrow{f} E$$

(resp.

$$\mathbb{D}_1 \otimes B \xrightarrow{\mathbb{D}_1 \otimes f} \mathbb{D}_1 \otimes C \xrightarrow{\alpha} D).$$

Remark 3.3.7. All the above descriptions of oplax transformations can be easily dualized for *lax transformations* (i.e. the 1-cells of the ω -category $\underline{\hom}_{lax}(X, Y)$ for some ω -categories X and Y). Habit is the only reason why we put emphasis on oplax transformations rather than lax transformations.

3.4 HOMOTOPY EQUIVALENCES AND DEFORMATION RETRACTS

3.4.1. Let C and D be two ω -categories and consider the smallest equivalence relation on the set $\operatorname{Hom}_{\omega \operatorname{Cat}}(C, D)$ such that two ω -functors from C to D are equivalent if there is an oplax direction between them (in any direction). Let us say that two ω -functors $u, v : C \to D$ are *oplax homotopic* if they are equivalent for this equivalence relation.

Definition 3.4.2. An ω -functor $u : C \to D$ is an *oplax homotopy equivalence* if there exists an ω -functor $v : D \to C$ such that $u \circ v$ is oplax homotopic to id_D and $v \circ u$ is oplax homotopic to id_C .

Recall that we write $\gamma^{\text{Th}} : \omega \text{Cat} \to ho(\omega \text{Cat}^{\text{Th}})$ for the localization functor with respect to the Thomason equivalences.

Lemma 3.4.3. Let $u, v : C \to D$ be two ω -functors. If there exists an oplax transformation $\alpha : u \Rightarrow v$, then $\gamma^{\text{Th}}(u) = \gamma^{\text{Th}}(v)$.

Proof. This follows immediately from [AM20c, Théorème B.11].

From this lemma and the fact that the Thomason equivalences are saturated (Corollary 3.2.7), we deduce the following proposition.

Proposition 3.4.4. *Every oplax homotopy equivalence is a Thomason equivalence.*

3.4.5. An ω -functor $i : C \to D$ is an *oplax deformation retract* if there exists an ω -functor $r : D \to C$ such that:

- (a) $r \circ i = \mathrm{id}_C$,
- (b) there exists an oplax transformation $\alpha : id_D \Rightarrow i \circ r$.

Furthermore, *i* is a *strong oplax deformation retract* if α can be chosen such that:

(c)
$$\alpha \star i = 1_i$$
.

An oplax deformation retract is a particular case of homotopy equivalence and thus of Thomason equivalence.

Lemma 3.4.6. The pushout of a strong oplax deformation retract is a strong oplax deformation retract.

Proof. Let $i : A \to B$ be a strong oplax deformation retract and

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & A' \\ \downarrow_{i} & & \downarrow_{i'} \\ B & \stackrel{v}{\longrightarrow} & B' \end{array} \tag{i}$$

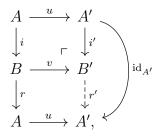
be a cocartesian square. We have to show that i' is also a strong oplax deformation retract. By hypothesis there exist $r: B \to A$ such that $r \circ i = id_A$ and $\alpha : \mathbb{D}_1 \otimes B \to B$ such that the diagrams

and

$$\mathbb{D}_1 \otimes A \xrightarrow{\mathbb{D}_1 \otimes i} \mathbb{D}_1 \otimes B \xrightarrow{\alpha} B, \qquad (iii)$$

where p is the unique morphism $\mathbb{D}_1 \to \mathbb{D}_0$, are commutative.

From the commutativity of the following solid arrow diagram



we deduce the existence of $r': B' \to A'$ that makes the whole diagram commute. In particular, we have $r' \circ i' = id_{A'}$.

From the commutativity of (iii), we easily deduce the commutativity of the following solid arrow diagram

The existence of $\alpha' : \mathbb{D}_1 \otimes B' \to B'$ that makes the whole diagram commutes follows from the fact that the functor $\mathbb{D}_1 \otimes$ – preserves colimits. In particular, we have

$$\alpha' \circ (\mathbb{D}_1 \otimes i') = p \otimes i'.$$

Now, notice that for every ω -category C, the maps

$$i_0^C: C \to \mathbb{D}_1 \otimes C \text{ and } i_1^C: C \to \mathbb{D}_1 \otimes C$$

are natural in C. Using this naturality and simple diagram chasing (left to the reader), we obtain the equalities

$$\alpha' \circ i_0^{B'} \circ v = v,$$

$$\alpha' \circ i_0^{B'} \circ i' = i',$$

and the equalities

$$\alpha' \circ i_1^{B'} \circ v = i' \circ r' \circ v$$
$$\alpha' \circ i_1^{B'} \circ i' = i' \circ r' \circ i'.$$

Using the fact that square (i) is cocartesian, we deduce that $\alpha' \circ i_0^{B'} = \operatorname{id}_{B'}$ and $\alpha' \circ i_1^{B'} = i' \circ r'$. This proves that i' is an oplax deformation retract, which is furthermore strong because of the equality $\alpha' \circ (\mathbb{D}_1 \otimes i') = p \otimes i'$.

In the following proposition, a *co-universal Thomason equivalence* means a couniversal weak equivalence for the localizer ($\omega Cat, \mathcal{W}_{\omega}^{Th}$) (Definition 2.1.2).

Proposition 3.4.7. *Every strong oplax deformation retract is a co-universal Thomason equivalence.*

Proof. Immediate consequence of Lemma 3.4.6 and the fact that oplax transformation retracts are Thomason equivalences. \Box

Remark 3.4.8. All the results we have seen in this section are still true if we replace "oplax" by "lax" everywhere.

3.5 EQUIVALENCES OF ω -CATEGORIES AND THE FOLK MODEL STRUCTURE

3.5.1. Let C be an ω -category. We define the equivalence relation \sim_{ω} on the set C_n by co-induction on $n \in \mathbb{N}$. For $x, y \in C_n$, we have $x \sim_{\omega} y$ when:

- x and y are parallel,

- there exist $r, s \in C_{n+1}$ such that $r: x \to y, s: y \to x$,

$$r *_n s \sim_\omega 1_y$$

and

$$s *_n r \sim_\omega 1_x.$$

For details on this definition and the proof that it is an equivalence relation, see [LMW10, section 4.2].

Example 3.5.2. Let x and y be two 0-cells of an n-category C.

- When n = 1, $x \sim_{\omega} y$ means that x and y are isomorphic.
- When n = 2, $x \sim_{\omega} y$ means that x and y are equivalent, i.e. there exist $f : x \to y$ and $g : y \to x$ such that fg is isomorphic to 1_y and gf is isomorphic to 1_x .

For later reference, we put here the following trivial but important lemma, whose proof is omitted.

Lemma 3.5.3. Let $F : C \to D$ be an ω -functor, $n \ge 0$ and x, y be n-cells of C. If $x \sim_{\omega} y$, then $F(x) \sim_{\omega} F(y)$.

Definition 3.5.4. An ω -functor $F: C \to D$ is an *equivalence of* ω -categories when:

- for every $y \in D_0$, there exists $x \in C_0$ such that

$$F(x) \sim_{\omega} y,$$

- for every $n \ge 0$, for all $x, y \in C_n$ that are parallel and for every $\beta \in D_{n+1}$ such that

$$\beta: F(x) \to F(y),$$

there exists $\alpha \in C_{n+1}$ such that

 $\alpha: x \to y$

and

$$F(\alpha) \sim_{\omega} \beta.$$

Example 3.5.5. If C and D are (small) categories seen as ω -categories, then a functor $F : C \to D$ is an equivalence of ω -categories if and only if it is fully faithful and essentially surjective, i.e. an equivalence of categories.

Theorem 3.5.6. There exists a cofibrantly generated model structure on ω Cat whose weak equivalences are the equivalences of ω -categories, and whose cofibrations are generated by the set $\{i_n : \mathbb{S}_{n-1} \to \mathbb{D}_n | n \in \mathbb{N}\}$ (see 1.1.7).

Proof. This is the main result of [LMW10].

3.5.7. The model structure of the previous theorem is commonly referred to as *folk model structure* on ω Cat. Data of this model structure will often be referred to by using the adjective folk, e.g. *folk cofibration*. Consequently *folk weak equivalence* and *equivalence of* ω -categories mean the same thing.

Furthermore, as in the Thomason case (see 3.2.3), we usually make reference to the word "folk" in the notations of homotopic constructions induced by the folk weak equivalences. For example, we write W^{folk} for the class of folk weak equivalences, $\mathcal{H}_{0}(\omega \mathbf{Cat}^{\text{folk}})$ for the homotopy op-prederivator of $(\omega \mathbf{Cat}, \mathcal{W}_{\omega}^{\text{folk}})$ and

$$\gamma^{\mathrm{folk}}:\omega\mathbf{Cat}
ightarrow\mathcal{H}\mathrm{o}(\omega\mathbf{Cat}^{\mathrm{folk}})$$

for the localization morphism. It follows from the previous theorem and Theorem 2.5.1 that the localizer (ωCat , $\mathcal{W}_{\omega}^{folk}$) is homotopy cocomplete. We will speak of "folk homotopy colimits" and "folk homotopy cocartesian squares" for homotopy colimits and homotopy cocartesian squares in this localizer.

3.5.8. Using the set $\{i_n : \mathbb{S}_{n-1} \to \mathbb{D}_n | n \in \mathbb{N}\}$ of generating folk cofibrations, we obtain that an ω -functor $F : C \to D$ is a *folk trivial fibration* when:

- for every $y \in D_0$, there exists $x \in C_0$ such that

$$F(x) = y$$

- for every $n \ge 0$, for all $x, y \in C_n$ that are parallel and for every $\beta \in D_{n+1}$ such that

$$\beta: F(x) \to F(y)$$

there exists $\alpha \in C_{n+1}$ such that

 $\alpha: x \to y$

and

$$F(\alpha) = \beta.$$

This characterization of folk trivial fibrations is to be compared with Definition 3.5.4 of equivalences of ω -categories.

Proposition 3.5.9. An ω -category is cofibrant for the folk model structure if and only if it is free.

Proof. The fact that every free ω -category is cofibrant follows immediately from the fact that the $i_n : \mathbb{S}_{n-1} \to \mathbb{D}_n$ are cofibrations and that every ω -category C is the colimit of the canonical diagram (Lemma 1.1.6)

$$\operatorname{sk}_0(C) \to \operatorname{sk}_1(C) \to \cdots \to \operatorname{sk}_n(C) \to \operatorname{sk}_{n+1}(C) \to \cdots$$

For the converse, see [Mét08].

3.6 EQUIVALENCES OF ω -CATEGORIES VS THOMASON EQUIVALENCES

Lemma 3.6.1. The nerve functor $N_{\omega} : \omega \mathbf{Cat} \to \widehat{\Delta}$ sends the equivalences of ω -categories to weak equivalences of simplicial sets.

Proof. Since every ω -category is fibrant for the folk model structure [LMW10, Proposition 9], it follows from Ken Brown's Lemma [Hov07, Lemma 1.1.12] that it suffices to show that the nerve sends the folk trivial fibrations to weak equivalences of simplicial sets. In particular, it suffices to show the stronger condition that the nerve sends the folk trivial fibrations to trivial fibrations to trivial fibrations of simplicial sets.

By adjunction, this is equivalent to showing that the functor $c_{\omega} : \widehat{\Delta} \to \omega \mathbf{Cat}$ sends the cofibrations of simplicial sets to folk cofibrations. Since c_{ω} is cocontinuous and the cofibrations of simplicial sets are generated by the inclusions

$$\partial \Delta_n \to \Delta_n$$

for $n \in \mathbb{N}$, it suffices to show that c_{ω} sends these inclusions to folk cofibrations.

Now, it follows from [Str87, Lemma 5.1] that the image of the inclusion $\partial \Delta_n \to \Delta_n$ by c_{ω} can be identified with the canonical inclusion

$$\operatorname{sk}_{n-1}(\mathcal{O}_n) \to \mathcal{O}_n$$

Since \mathcal{O}_n is free, this last morphism is by definition a push-out of a coproduct of folk cofibrations (see Definition 1.2.1), hence a folk cofibration.

As an immediate consequence of the previous lemma, we have the following proposition.

Proposition 3.6.2. Every equivalence of ω -categories is a Thomason equivalence.

Remark 3.6.3. The converse of the above proposition is false. For example, the unique ω -functor

 $\mathbb{D}_1 \to \mathbb{D}_0$

is a Thomason equivalence because its image by the nerve is the unique morphism of simplicial sets $\Delta_1 \rightarrow \Delta_0$ (which obviously is a weak equivalence), but it is *not* an equivalence of ω -categories because \mathbb{D}_1 and \mathbb{D}_0 are not equivalent as categories (see Example 3.5.5).

3.6.4. Proposition 3.6.2 implies that the identity functor on ω Cat induces a morphism of localizers (ω Cat, $\mathcal{W}^{\text{folk}}$) \rightarrow (ω Cat, \mathcal{W}^{Th}), which in turn induces a functor between localized categories

$$\mathcal{J}: ho(\omega \mathbf{Cat}^{\mathrm{folk}}) \to ho(\omega \mathbf{Cat}^{\mathrm{Th}}).$$

This functor cannot be an equivalence since this would imply that every Thomason equivalence is an equivalence of ω -categories.

3.7 SLICE ω -CATEGORIES AND FOLK THEOREM A

3.7.1. Let A be an ω -category and a_0 an object of A. We define the slice ω -category A/a_0 as the following fibred product:

$$\begin{array}{ccc} A/a_0 & \longrightarrow & \underline{\hom}_{\mathrm{lax}}(\mathbb{D}_1, A) \\ \downarrow & & \downarrow & \downarrow \\ \mathbb{D}_0 & & & \downarrow \\ \pi_1^A & & & \downarrow \\ \mathbb{D}_0 & & & A. \end{array}$$

We also define an ω -functor $\pi : A/a_0 \to A$ as the following composition

$$\pi: A/a_0 \to \underline{\mathrm{hom}}_{\mathrm{lax}}(\mathbb{D}_1, A) \xrightarrow{\pi_0^A} A.$$

Let us now give an alternative definition of the ω -category A/a_0 using explicit formulas. The equivalence with the previous definition follows from the dual of [AM20b, Proposition B.5.2]

- An *n*-cell of A/a_0 is a table

$$(x,a) = \begin{pmatrix} (x_0,a_1) & (x_1,a_2) & \cdots & (x_{n-1},a_n) \\ (x'_0,a'_1) & (x'_1,a'_2) & \cdots & (x'_{n-1},a'_n) \end{pmatrix}$$

where x_0 and x'_0 are 0-cells of A, and:

 $\begin{aligned} x_i : x_{i-1} &\longrightarrow x'_{i-1}, & \text{for every } 1 \le i \le n, \\ x'_i : x_{i-1} &\longrightarrow x'_{i-1}, & \text{for every } 1 \le i \le n-1, \\ a_i : a'_{i-1} \underset{i-2}{*} a'_{i-2} \underset{i-3}{*} \cdots \underset{1}{*} a'_1 \underset{0}{*} x_{i-1} &\longrightarrow a_{i-1}, & \text{for every } 1 \le i \le n+1, \\ a'_i : a'_{i-1} \underset{i-2}{*} a'_{i-2} \underset{1}{*} \cdots \underset{1}{*} a'_1 \underset{0}{*} x'_{i-1} &\longrightarrow a_{i-1}, & \text{for every } 1 \le i \le n \end{aligned}$

are *i*-cells of *A*. In low dimension, this gives:

- The source and target of the *n*-cell (a, x) are given by the matrices:

$$s(x,a) = \begin{pmatrix} (x_0,a_1) & (x_1,a_2) & \cdots & (x_{n-2},a_{n-1}) \\ (x'_0,a'_1) & (x'_1,a'_2) & \cdots & (x'_{n-2},a'_{n-1}) \end{pmatrix}$$
$$t(x,a) = \begin{pmatrix} (x_0,a_1) & (x_1,a_2) & \cdots & (x_{n-2},a_{n-1}) \\ (x'_0,a'_1) & (x'_1,a'_2) & \cdots & (x'_{n-2},a'_{n-1}) \end{pmatrix} (x'_{n-1},a'_n) \end{pmatrix}.$$

- The unit of the *n*-cell (a, x) is given by the table:

$$1_{(x,a)} = \begin{pmatrix} (x_0, a_1) & (x_1, a_2) & \cdots & (x_{n-1}, a_n) & (x_n, a_{n+1}) \\ (x'_0, a'_1) & (x'_1, a'_2) & \cdots & (x'_{n-1}, a'_n) & (x_n, a_{n+1}) \end{pmatrix}$$

- The composition of *n*-cells (x, a) and (y, b) such that $s_k(y, b) = t_k(a, x)$, is given by the table:

$$(y,b)_{k}^{*}(x,a) = \begin{pmatrix} (x_{0},a_{1}) & \cdots & (x_{k},a_{k}) & (z_{k+1},c_{k+2}) & \cdots & (z_{n-1},c_{n}) \\ (y'_{0},b'_{1}) & \cdots & (y'_{k},b'_{k}) & (z'_{k+1},c'_{k+2}) & \cdots & (z'_{n-1},c'_{n}) \end{pmatrix},$$

where:

$$\begin{aligned} z_{i} &= y_{i} * x_{i}, & \text{for every } k+1 \leq i \leq n, \\ z'_{i} &= y'_{i} * x'_{i}, & \text{for every } k+1 \leq i \leq n-1, \\ c_{i} &= a_{i} * b_{i} * a'_{k} * a'_{k-1} * \cdots * a'_{1} * x_{k+1}, & \text{for every } k+2 \leq i \leq n+1, \\ c'_{i} &= a'_{i} * b'_{i} * a'_{k} * a'_{k-1} * \cdots * a'_{1} * x'_{k+1}, & \text{for every } k+2 \leq i \leq n+1, \end{aligned}$$

We leave it to the reader to check that the formulas are well defined and that the axioms for ω -categories are satisfied. The canonical forgetful ω -functor $\pi : A/a_0 \to A$ is simply expressed as:

$$A/a_0 \to A$$
$$(x,a) \mapsto x_n.$$

Notice that if A is an n-category, then so is A/a_0 . In this case, for an n-cell (x, a), a_{n+1} is a unit, hence

$$a'_{n} \underset{n-1}{*} a'_{n-1} \underset{n-2}{*} \cdots \underset{1}{*} a'_{1} \underset{0}{*} x_{n} = a_{n}.$$

Example 3.7.2. For a small category A (considered as an ω -category) and an object a_0 of A, the category A/a_0 in the sense of the previous paragraph is nothing but the usual slice category of A over a_0 .

3.7.3. Let $u : A \to B$ be a morphism of ω Cat and b_0 an object of B. We define the ω -category A/b_0 (also denoted by $u \downarrow b_0$) as the following fibred product:

More explicitly, an *n*-cell (x, b) of A/b_0 is a table

$$(x,b) = \begin{pmatrix} (x_0,b_1) & (x_1,b_2) & \cdots & (x_{n-1},b_n) \\ (x'_0,b'_1) & (x'_1,b'_2) & \cdots & (x'_{n-1},b'_n) \end{pmatrix}$$

where the x_i and x'_i are *i*-cells of A such that $x_i: x_{i-1} \longrightarrow x'_{i-1}$, for every $1 \le i \le n$,

 $x'_i: x_{i-1} \longrightarrow x'_{i-1}, \text{ for every } 1 \le i \le n-1,$

and the b_i and b'_i are *i*-cells of B such that

$$\begin{pmatrix} (u(x_0), b_1) & (u(x_1), b_2) & \cdots & (u(x_{n-1}), b_n) \\ (u(x'_0), b'_1) & (u(x'_1), b'_2) & \cdots & (u(x'_{n-1}), b'_n) \end{pmatrix}$$

is an *n*-cell of B/b_0 .

The canonical ω -functor $A/b_0 \rightarrow A$ is simply expressed as

$$A/b_0 \to A$$
$$(x,b) \mapsto x_n,$$

and the ω -functor u/b_0 as

$$u/b_0: A/b_0 \to B/b_0$$
$$(x,b) \mapsto (u(x),b).$$

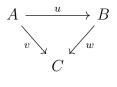
More generally, if we have a commutative triangle in ω Cat

$$\begin{array}{ccc} A & \xrightarrow{u} & B \\ & \swarrow & \swarrow \\ v & \swarrow & v \\ & C & , \end{array}$$

then for every object c_0 of C, we have a functor $u/c_0 : A/c_0 \to B/c_0$ defined as

$$u/c_0 : A/c_0 \to B/c_0$$
$$(x,c) \mapsto (u(x),c)$$

Theorem 3.7.4. (Folk Theorem A) Let



be a commutative triangle in ω Cat. If for every object c_0 of C the induced morphism

$$u/c_0: A/c_0 \to B/c_0$$

is an equivalence of ω -categories, then so is u.

Proof. Before anything else, recall from Lemma 3.5.3 that given an ω -functor $F : X \to Y$ and *n*-cells x and y of X, if $x \sim_{\omega} y$, then $F(x) \sim_{\omega} F(y)$.

- (i) Let b_0 be a 0-cell of B and set $c_0 := w(b_0)$. By definition, the pair $(b_0, 1_{c_0})$ is a 0-cell of B/c_0 . By hypothesis, we know that there exists a 0-cell (a_0, c_1) of A/c_0 such that $(u(a_0), c_1) \sim_{\omega} (b_0, 1_{c_0})$. Hence, by applying the canonical functor $B/c_0 \to B$, we obtain that $u(a_0) \sim_{\omega} b_0$.
- (ii) Let f and f' be parallel n-cells of A and β : u(f) → u(f') an (n + 1)-cell of B. We need to show that there exists an (n + 1)-cell α : f → f' of A such that u(α) ~_ω β.

Let us use the notations:

-
$$a_i := s_i(f) = s_i(f')$$
 for $0 \le i < n$,
- $a'_i := t_i(f) = t_i(f')$ for $0 \le i < n$,
- $a_n := f$ and $a'_n = f'$.

It is straightforward to check that

$$\begin{pmatrix} (a_0, v(a'_1)) & (a_1, v(a'_2)) & \cdots & (a_{n-1}, v(a'_n)) \\ (a'_0, 1_{v(a'_0)}) & (a'_1, 1_{v(a'_1)}) & \cdots & (a'_{n-1}, 1_{(v(a'_{n-1}))}) \end{pmatrix} (a_n, w(\beta)) \end{pmatrix}$$

and

$$\begin{pmatrix} (a_0, v(a'_1)) & (a_1, v(a'_2)) & \cdots & (a_{n-1}, v(a'_n)) \\ (a'_0, 1_{v(a'_0)}) & (a'_1, 1_{v(a'_1)}) & \cdots & (a'_{n-1}, 1_{(v(a'_{(n-1)}))}) \end{pmatrix} (a'_n, 1_{v(a'_n)}) \end{pmatrix}$$

are parallel *n*-cells of A/c_0 where we set $c_0 := v(a'_0)$. Similarly, we have an (n+1)-cell of B/c_0

$$\begin{pmatrix} (u(a_0), v(a'_1)) & \cdots & (u(a_{n-1}), v(a'_n)) & (u(a_n), w(\beta)) \\ (u(a'_0), 1_{v(a'_0)}) & \cdots & (u(a'_{n-1}), 1_{(v(a'_{n-1}))}) & (u(a'_n), 1_{v(a'_n)}) \end{pmatrix}$$

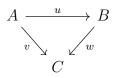
whose source and target respectively are the image by u/c_0 of the above two cells of A/c_0 . By hypothesis, there exists an (n + 1)-cell of A/c_0 of the form

$$\begin{pmatrix} (a_0, v(a'_1)) & \cdots & (a_{n-1}, v(a'_n)) & (a_n, w(\beta)) \\ (a'_0, 1_{v(a'_0)}) & \cdots & (a'_{n-1}, 1_{(v(a'_{n-1}))}) & (a'_n, 1_{v(a'_n)}) \end{pmatrix} \begin{pmatrix} \alpha, \Lambda \end{pmatrix}$$

whose image by u/c_0 is equivalent for the relation \sim_{ω} to the above (n + 1)-cell of B/c_0 . In particular, the source and target of α are respectively f and f'. Finally, we obtain that $u(\alpha) \sim_{\omega} \beta$ by applying the canonical ω -functor $B/c_0 \rightarrow B$.

3.7.5. The name "folk Theorem A" is an explicit reference to Quillen's Theorem A [Qui73, Theorem A] and its ω -categorical generalization by Ara and Maltsiniotis [AM18, AM20c]. For the sake of comparison we recall below the latter one.

Theorem 3.7.6 (Ara and Maltsiniotis' Theorem A). Let



be a commutative triangle in ω **Cat**. If for every object c_0 of C the induced morphism

$$u/c_0: A/c_0 \to B/c_0$$

is a Thomason equivalence, then so is u.

CHAPTER 3. HOMOTOPY THEORY OF ω -CATEGORIES

CHAPTER 4

HOMOLOGY AND ABELIANIZATION OF ω -CATEGORIES

4.1 HOMOLOGY VIA THE NERVE

4.1.1. We denote by $Ch_{\geq 0}$ the category of non-negatively graded chain complexes of abelian groups. Recall that $Ch_{\geq 0}$ can be equipped with a cofibrantly generated model structure, known as the *projective model structure on* $Ch_{>0}$, where:

- the weak equivalences are the quasi-isomorphisms, i.e. morphisms of chain complexes that induce an isomorphism on homology groups,
- the cofibrations are the morphisms of chain complexes f : X → Y such that for every n ≥ 0, f_n : X_n → Y_n is a monomorphism with projective cokernel,
- the fibrations are the morphisms of chain complexes $f : X \to Y$ such that for every n > 0, $f_n : X_n \to Y_n$ is an epimorphism.

(See for example [DS95, Section 7].) From now on, we will implicitly consider that the category $Ch_{\geq 0}$ is equipped with this model structure.

4.1.2. Let X be a simplicial set. We denote by $K_n(X)$ the abelian group of n-chains of X, i.e. the free abelian group on the set X_n . For n > 0, let $\partial : K_n(X) \to K_{n-1}(X)$ be the linear map defined for $x \in X_n$ by

$$\partial(x) := \sum_{i=0}^{n} (-1)^{i} \partial_{i}(x).$$

It follows from the simplicial identities (see [GZ67, section 2.1]) that $\partial \circ \partial = 0$. Hence,

the previous data defines a chain complex K(X) and this defines a functor

$$K:\widehat{\Delta}\to \mathbf{Ch}_{\geq 0}$$
$$X\mapsto K(X)$$

in the expected way.

4.1.3. Recall that an *n*-simplex *x* of a simplicial set *X* is *degenerate* if there exists an epimorphism $\varphi : [n] \to [m]$ with m < n and an *m*-simplex *y* such that $X(\varphi)(y) = x$. We denote by $D_n(X)$ the subgroup of $K_n(X)$ generated by the degenerate *n*-simplices and by $\kappa_n(X)$ the abelian group of *normalized n-chains*:

$$\kappa_n(X) = K_n(X) / D_n(X).$$

Using the simplicial identities, it can be shown that $\partial(D_n(X)) \subseteq D_{n-1}(X)$ for every n > 0. Hence, there is an induced differential which we still denote by ∂ :

$$\partial : \kappa_n(X) \to \kappa_{n-1}(X).$$

This defines a chain complex $\kappa(X)$, which we call the *normalized chain complex of* X. This yields a functor

$$\kappa : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$$
$$X \mapsto \kappa(X).$$

Lemma 4.1.4. The functor $\kappa : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$ is left Quillen and sends the weak equivalences of simplicial sets to quasi-isomorphisms.

Proof. Recall that the Quillen model structure on simplicial sets admits the set of inclusions

$$\{\partial \Delta_n \hookrightarrow \Delta_n | n \in \mathbb{N}\}$$

as generating cofibrations and the set of inclusions

$$\{\Lambda_n^i \hookrightarrow \Delta_n | n \in \mathbb{N}, 0 \le i \le n\}$$

as generating trivial cofibrations (see for example [GJ09, Section I.1] for the notations). A quick computation, which we leave to the reader, shows that the image by κ of $\partial \Delta_n \hookrightarrow \Delta_n$ is a monomorphism with projective cokernel and the image by κ of $\Lambda_n^i \hookrightarrow \Delta_n$ is a quasi-isomorphism. This proves that κ is left Quillen. Since all simplicial sets are cofibrant, it follows from Ken Brown's Lemma [Hov07, Lemma 1.1.12] that κ also preserves weak equivalences. **Remark 4.1.5.** The previous lemma admits also a more conceptual proof as follows. From the Dold–Kan equivalence, we know that $\mathbf{Ch}_{\geq 0}$ is equivalent to the category $\mathbf{Ab}(\Delta)$ of simplicial abelian groups and with this identification the functor $\kappa : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$ is left adjoint of the canonical forgetful functor

$$U: \mathbf{Ch}_{>0} \simeq \mathbf{Ab}(\Delta) \to \widehat{\Delta}$$

induced by the forgetful functor from abelian groups to sets. The fact that U is right Quillen follows then from [GJ09, Lemma 2.9 and Corollary 2.10].

4.1.6. In particular, κ induces a morphism of localizers

$$\kappa : (\widehat{\Delta}, \mathcal{W}_{\Delta}) \to (\mathbf{Ch}_{>0}, \mathcal{W}_{\mathbf{Ch}_{>0}}),$$

where we wrote $\mathcal{W}_{\mathbf{Ch}_{>0}}$ for the class of quasi-isomorphisms.

Definition 4.1.7. The singular homology functor for ω -categories \mathbb{H}^{Sing} is defined as the following composition

$$\mathbb{H}^{\operatorname{Sing}} : \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{Th}}) \xrightarrow{\overline{N_{\omega}}} \operatorname{ho}(\widehat{\Delta}) \xrightarrow{\overline{\kappa}} \operatorname{ho}(\mathbf{Ch}_{\geq 0}).$$

For an ω -category C, $\mathbb{H}^{\text{Sing}}(C)$ is the singular homology of C.

4.1.8. In other words, the singular homology of C is the chain complex $\kappa(N_{\omega}(C))$ seen as an object of $ho(\mathbf{Ch}_{\geq 0})$ (see Remark 2.1.4). For $k \geq 0$, the k-th singular homology group of an ω -category C is defined as

$$H_k^{\operatorname{Sing}}(C) := H_k(\mathbb{H}^{\operatorname{Sing}}(C)) = H_k(\kappa(N_{\omega}(C))),$$

where H_k : ho(Ch_{≥ 0}) \rightarrow Ab is the usual functor that associates to an object of ho(Ch_{≥ 0}) its k-th homology group.

Remark 4.1.9. The adjective "singular" is there to avoid future confusion with another homological invariant for ω -categories that will be introduced later. As a matter of fact, the underlying point of view adopted in this thesis is that *singular homology of* ω -categories ought to be simply called *homology of* ω -categories as it is the only "correct" definition of homology. This assertion will be justified in Remark 4.5.4.

Remark 4.1.10. We could also have defined the singular homology of ω -categories using $K : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$ instead of $\kappa : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$ since these two functors are quasiisomorphic (see [GJ09, Theorem 2.4] for example). An advantage of the latter one is that it is left Quillen. **4.1.11.** We will also denote by \mathbb{H}^{Sing} the morphism of op-prederivators defined as the following composition

$$\mathbb{H}^{\mathrm{Sing}}: \mathcal{H}o(\omega\mathbf{Cat}^{\mathrm{Th}}) \xrightarrow{\overline{N_{\omega}}} \mathcal{H}o(\widehat{\Delta}) \xrightarrow{\overline{\kappa}} \mathcal{H}o(\mathbf{Ch}_{\geq 0}).$$

Proposition 4.1.12. The singular homology

$$\mathbb{H}^{\mathrm{Sing}}: \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\mathbf{Ch}_{>0})$$

is homotopy cocontinuous.

Proof. This follows from the fact that $\overline{N_{\omega}}$ and $\overline{\kappa}$ are both homotopy cocontinuous. For $\overline{N_{\omega}}$, this is because it is an equivalence of op-prederivators and thus we can apply Lemma 2.3.10 and Lemma 2.3.11. For $\overline{\kappa}$, this is because κ is left Quillen and thus we can apply Theorem 2.5.2.

4.2 ABELIANIZATION

We write Ab for the category of abelian groups and for an abelian group G, we write |G| for the underlying set of G.

4.2.1. Let C be an ω -category. For every $n \ge 0$, we define $\lambda_n(C)$ as the abelian group obtained by quotienting $\mathbb{Z}C_n$ (the free abelian group on C_n) by the congruence generated by the relations

$$x * y \sim x + y$$

for all $x, y \in C_n$ that are k-composable for some k < n. For n = 0, this means that $\lambda_0(C) = \mathbb{Z}C_0$. Now let $f : C \to D$ be an ω -functor. For every $n \ge 0$, the definition of ω -functor implies that the linear map

$$\mathbb{Z}C_n \to \mathbb{Z}D_n$$
$$x \in C_n \mapsto f(x)$$

induces a linear map

$$\lambda_n(f) : \lambda_n(C) \to \lambda_n(D).$$

This defines a functor $\lambda_n : \omega \mathbf{Cat} \to \mathbf{Ab}$.

For n > 0, consider the linear map

$$\mathbb{Z}C_n \to \mathbb{Z}C_{n-1}$$
$$x \in C_n \mapsto t(x) - s(x).$$

The axioms of ω -categories imply that it induces a map

$$\partial: \lambda_n(C) \to \lambda_{n-1}(C)$$

which is natural in C. Furthermore, it satisfies the equation $\partial \circ \partial = 0$. Thus, for every ω -category C, we have defined a chain complex $\lambda(C)$:

$$\lambda_0(C) \xleftarrow{\partial} \lambda_1(C) \xleftarrow{\partial} \lambda_2(C) \xleftarrow{\partial} \cdots$$

and for every $f: C \to D$ a morphism of chain complexes

$$\lambda(f):\lambda(C)\to\lambda(D).$$

Altogether, this defines a functor

$$\lambda: \omega \mathbf{Cat} \to \mathbf{Ch}_{>0},$$

which we call the *abelianization functor*.

Lemma 4.2.2. *The functor* λ *is a left adjoint.*

Proof. The category $\mathbf{Ch}_{\geq 0}$ is equivalent to the category $\omega \mathbf{Cat}(\mathbf{Ab})$ of ω -categories internal to abelian groups (see [Bou90, Theorem 3.3]) and with this identification, the functor $\lambda : \omega \mathbf{Cat} \to \omega \mathbf{Cat}(\mathbf{Ab})$ is nothing but the left adjoint of the canonical forgetful functor $\omega \mathbf{Cat}(\mathbf{Ab}) \to \omega \mathbf{Cat}$.

As we shall now see, when the ω -category C is *free* the chain complex $\lambda(C)$ admits a nice expression.

4.2.3. Let $n \ge 0$. Recall that for every monoid M (supposed commutative if $n \ge 1$) we have defined in Section 1.3 an *n*-category $B^n M$ whose set of *n*-cells is isomorphic to the underlying set of M. And the correspondence $M \mapsto B^n M$ defines a functor in the expected way. By considering abelian groups as particular cases of (commutative) monoids, we obtain a functor for each $n \ge 0$

$$B^n : \mathbf{Ab} \to n\mathbf{Cat}$$
$$G \mapsto B^n G,$$

where Ab is the category of abelian groups.

Besides, let us write λ_n again for the functor

$$\lambda_n : n\mathbf{Cat} \to \mathbf{Ab}$$

 $C \mapsto \lambda_n(C)$

(That is the restriction of $\lambda_n : \omega \mathbf{Cat} \to \mathbf{Ab}$ to $n\mathbf{Cat}$.)

Lemma 4.2.4. Let $n \ge 0$. The functor $\lambda_n : n\mathbf{Cat} \to \mathbf{Ab}$ is left adjoint to the functor $B^n : \mathbf{Ab} \to n\mathbf{Cat}$.

Proof. The case n = 0 is immediate since the functor $\lambda_0 : 0\mathbf{Cat} = \mathbf{Set} \to \mathbf{Ab}$ is the "free abelian group" functor and the functor $B^0 : \mathbf{Ab} \to 0\mathbf{Cat} = \mathbf{Set}$ is the "underlying set" functor.

Suppose now that n > 0. From Lemma 1.3.3 we know that for every abelian group G and every *n*-category C, the map

$$\operatorname{Hom}_{n\mathbf{Cat}}(C, B^n G) \to \operatorname{Hom}_{\mathbf{Set}}(C_n, |G|)$$
$$F \mapsto F_n,$$

is injective and its image consists of those functions $f: C_n \to |G|$ such that:

(i) for every $0 \le k < n$ and every pair (x, y) of k-composable n-cells of C, we have

$$f(x * y) = f(x) + f(y),$$

(ii) for every $x \in C_{n-1}$, we have

$$f(1_x) = 0$$

Let us see that because G is an abelian group (recall that Lemma 1.3.3 was stated for the general case of commutative monoids), condition (i) imply condition (ii). Let $f : C_n \to |G|$ be a function that satisfies condition (i) and let $x \in C_{n-1}$. We have $1_x * 1_x = 1_x$, hence

$$f(1_x) = f(1_x *_{n-1} 1_x) = f(1_x) + f(1_x),$$

and then

$$f(1_x) = 0$$

because every element of an (abelian) group has an inverse. Now, because of the adjunction morphism

$$\operatorname{Hom}_{\operatorname{Set}}(C_n, |G|) \simeq \operatorname{Hom}_{\operatorname{Ab}}(\mathbb{Z}C_n, G),$$

we have that $\operatorname{Hom}_{n\mathbf{Cat}}(C, B^n G)$ is naturally isomorphic to the set of morphisms of abelian groups $g : \mathbb{Z}C_n \to G$ such that for every pair (x, y) of k-composable elements of C_n for some k < n, we have

$$g(x * y) = g(x) + g(y).$$

By definition, this set is naturally isomorphic to the set of morphisms of abelian groups from $\lambda_n(C)$ to G. In other words, we have

$$\operatorname{Hom}_{n\mathbf{Cat}}(C, B^{n}G) \simeq \operatorname{Hom}_{\mathbf{Ab}}(\lambda_{n}(C), G).$$

4.2.5. Let C be an ω -category, $n \in \mathbb{N}$ and $E \subseteq C_n$ a subset of the n-cells. We obtain a map $\mathbb{Z}E \to \lambda_n(C)$ defined as the composition

$$\mathbb{Z}E \to \mathbb{Z}C_n \to \lambda_n(C),$$

where the map on the left is induced by the canonical inclusion of E in C_n and the map on the right is the quotient map from the definition of $\lambda_n(C)$.

Lemma 4.2.6. Let C be a free ω -category and let $\Sigma = (\Sigma_n)_{n \in \mathbb{N}}$ be its basis. For every $n \in \mathbb{N}$, the map

$$\mathbb{Z}\Sigma_n \to \lambda_n(C)$$

from the previous paragraph, is an isomorphism.

Proof. Notice first that for every ω -category C, we have $\lambda_n(\tau_{\leq n}^s(C)) = \lambda_n(C)$. Suppose now that C is free with basis $\Sigma = (\Sigma_n)_{n \in \mathbb{N}}$. Using Lemma 4.2.4 and Lemma 1.3.4, we obtain that for every abelian group G, we have

$$\operatorname{Hom}_{\mathbf{Ab}}(\lambda_n(C), G) \simeq \operatorname{Hom}_{\mathbf{Ab}}(\lambda_n(\tau_{\leq n}^s(C)), G)$$
$$\simeq \operatorname{Hom}_{n\mathbf{Cat}}(\tau_{\leq n}^s(C), B^n G)$$
$$\simeq \operatorname{Hom}_{\mathbf{Set}}(\Sigma_n, |G|)$$
$$\simeq \operatorname{Hom}_{\mathbf{Ab}}(\mathbb{Z}\Sigma_n, G),$$

and it is easily checked that this isomorphism is induced by precomposition with the map $\mathbb{Z}\Sigma_n \to \lambda_n(C)$ from the previous paragraph. The result follows then from the Yoneda Lemma.

4.2.7. Let C be a free ω -category and write $\Sigma = (\Sigma_n)_{n \in \mathbb{N}}$ for its basis. For every $n \ge 0$ and every $\alpha \in \Sigma_n$, recall that we have proved in Proposition 1.3.5 the existence of a unique function $w_\alpha : C_n \to \mathbb{N}$ such that:

- (a) $w_{\alpha}(\alpha) = 1$,
- (b) $w_{\alpha}(\beta) = 0$ for every $\beta \in \Sigma_n$ such that $\beta \neq \alpha$,
- (c) for every pair of k-composable n-cells of C for some k < n, we have

$$w_{\alpha}(x * y) = w_{\alpha}(x) + w_{\alpha}(y).$$

We can then define for each $n \ge 0$, a map $w_n : C_n \to \mathbb{Z}\Sigma_n$ with the formula

$$w_n(x) = \sum_{\alpha \in \Sigma_n} w_\alpha(x) \cdot \alpha$$

for every $x \in C_n$.

Condition (c) implies that

$$w_n(x * y) = w_n(x) + w_n(y)$$

for every pair (x, y) of k-composable n-cells of C, and conditions (a) and (b) imply that

 $w_n(\alpha) = \alpha$

for every $\alpha \in \Sigma_n$.

Proposition 4.2.8. Let C be a free ω -category and let $(\Sigma_n)_{n \in \mathbb{N}}$ be its basis. The chain complex $\lambda(C)$ is canonically isomorphic to the chain complex

 $\mathbb{Z}\Sigma_0 \xleftarrow{\partial} \mathbb{Z}\Sigma_1 \xleftarrow{\partial} \mathbb{Z}\Sigma_2 \xleftarrow{\partial} \cdots$

where $\partial : \mathbb{Z}\Sigma_n \to \mathbb{Z}\Sigma_{n-1}$ is the linear map defined by the formula

$$\partial(x) = w_{n-1}(\mathbf{t}(x)) - w_{n-1}(\mathbf{s}(x))$$

for every $x \in \Sigma_n$.

With this identification, if C' is another free ω -category and if $F : C \to C'$ is an ω -functor (not necessarily rigid), then the map $\lambda_n(F) : \lambda_n(C) \to \lambda_n(C')$ reads

$$\lambda_n(F)(x) = w'_n(F(x))$$

for every $x \in \Sigma_n$.

Proof. For $n \ge 0$, write $\phi_n : \mathbb{Z}\Sigma_n \to \lambda_n(C)$ for the map defined in 4.2.5 (which we know is an isomorphism from Lemma 4.2.6).

The map $w_n : C_n \to \mathbb{Z}\Sigma_n$ induces a map $\mathbb{Z}C_n \to \mathbb{Z}\Sigma_n$ by linearity, which in turn induces a map $\lambda_n(C) \to \mathbb{Z}\Sigma_n$ (because $w_n(x * y) = w_n(x) + w_n(y)$ for every pair (x, y)of k-composable n-cells). Write ψ_n for this last map. It is immediate to check that the composition

$$\mathbb{Z}\Sigma_n \xrightarrow{\phi_n} \lambda_n(C) \xrightarrow{\psi_n} \mathbb{Z}\Sigma_n$$

gives the identity on $\mathbb{Z}\Sigma_n$. Hence, ψ_n is the inverse of ϕ_n .

Now, for n > 0, notice that the map $\partial : \mathbb{Z}\Sigma_n \to \mathbb{Z}\Sigma_{n-1}$ given in the statement of the proposition is nothing but the composition

$$\mathbb{Z}\Sigma_n \xrightarrow{\phi_n} \lambda_n(C) \xrightarrow{\partial} \lambda_{n-1}(C) \xrightarrow{\psi_{n-1}} \mathbb{Z}\Sigma_{n-1}.$$

The first part of the proposition follows then from Lemma 4.2.6.

As for the second part, it suffices to notice that if we identify $\lambda_n(C)$ with $\mathbb{Z}\Sigma_n$ via ϕ_n for every free ω -category C, then map $\mathbb{Z}\Sigma_n \to \mathbb{Z}\Sigma'_n$ (where Σ'_n is the *n*-basis of C') induced by F is given by the composition

$$\mathbb{Z}\Sigma_n \xrightarrow{\phi_n} \lambda_n(C) \xrightarrow{\lambda_n(F)} \lambda_n(C') \xrightarrow{\psi_n} \mathbb{Z}\Sigma'_n.$$

4.3 POLYGRAPHIC HOMOLOGY

4.3.1. Let $f, g : K \to K'$ be two morphisms of non-negatively graded chain complexes. Recall that a *chain homotopy* from f to g consists of a sequence of linear maps $(h_n \colon K_n \to K'_{n+1})_{n \in \mathbb{N}}$ such that

$$\partial \circ h_0 = g_0 - f_0$$

and such that for every n > 0, we have

$$\partial \circ h_n + h_{n-1} \circ \partial = g_n - f_n.$$

Recall also that if there is a chain homotopy from f to g, then the localization functor $\gamma^{\mathbf{Ch}_{\geq 0}} : \mathbf{Ch}_{>0} \to \mathrm{ho}(\mathbf{Ch}_{>0})$ identifies f and g, which means that

$$\gamma^{\mathbf{Ch}_{\geq 0}}(f) = \gamma^{\mathbf{Ch}_{\geq 0}}(g).$$

Lemma 4.3.2. Let $u, v : C \to D$ be two ω -functors. If there is an oplax transformation $\alpha : u \Rightarrow v$, then there is a chain homotopy from $\lambda(u)$ to $\lambda(v)$.

Proof. For an *n*-cell x of C (resp. D), let us use the notation [x] for the image of x in $\lambda_n(C)$ (resp. $\lambda_n(D)$).

Let h_n be the map

$$h_n : \lambda_n(C) \to \lambda_{n+1}(D)$$

 $[x] \mapsto [\alpha_x].$

The formulas for oplax transformations from Paragraph 3.3.3 imply that h_n is linear and that for every *n*-cell *x* of *C*, if n = 0, we have

$$\partial(h_0(x)) = [v(x)] - [u(x)],$$

and if n > 0, we have

$$\partial(h_n(x)) + h_{n-1}(\partial(x)) = [v(x)] - [u(x)].$$

Details are left to the reader.

Proposition 4.3.3. The abelianization functor $\lambda : \omega Cat \to Ch_{\geq 0}$ is left Quillen with respect to the folk model structure on ωCat .

Proof. The fact that λ is a left adjoint is Lemma 4.2.2.

A simple computation using Lemma 4.2.8 shows that for every $n \in \mathbb{N}$,

$$\lambda(i_n):\lambda(\mathbb{S}_{n-1})\to\lambda(\mathbb{D}_n)$$

is a monomorphism with projective cokernel. Hence λ sends folk cofibrations to cofibrations of chain complexes.

Then, we know from [LMW10, Sections 4.6 and 4.7] and [AM20b, Remarque B.1.16] (see also [AL20, Paragraph 3.11]) that there exists a set of generating trivial cofibrations J of the folk model structure on ω Cat such that every $j : X \to Y$ in J is a deformation retract (see Paragraph 3.4.5). From Lemma 4.3.2, we conclude that λ sends folk trivial cofibrations to trivial cofibrations of chain complexes.

In particular, λ is totally left derivable (when ω Cat is equipped with folk weak equivalences). This motivates the following definition.

Definition 4.3.4. The polygraphic homology functor

$$\mathbb{H}^{\mathrm{pol}}:\mathrm{ho}(\omega\mathbf{Cat}^{\mathrm{folk}})\to\mathrm{ho}(\mathbf{Ch}_{\geq 0})$$

is the total left derived functor of $\lambda : \omega \mathbf{Cat} \to \mathbf{Ch}_{\geq 0}$ (where $\omega \mathbf{Cat}$ is equipped with folk weak equivalences). For an ω -category C, $\mathbb{H}^{\mathrm{pol}}(C)$ is the *polygraphic homology of* C.

4.3.5. Similarly to singular homology groups, for $k \ge 0$ the k-th polygraphic homology group of an ω -category C is defined as

$$H_k^{\mathrm{pol}}(C) := H_k(\mathbb{H}^{\mathrm{pol}}(C))$$

where H_k : ho($Ch_{\geq 0}$) \rightarrow Ab is the usual functor that associate to an object of ho($Ch_{\geq 0}$) its k-th homology group. In practice, this means that one has to find a cofibrant replacement of C, that is to say a free ω -category P and a folk trivial fibration

$$P \to C$$
,

and then the polygraphic homology groups of C are those of $\lambda(P)$ which are computed using Proposition 4.2.8.

4.3.6. For later reference, let us recall here that since \mathbb{H}^{pol} is the left derived functor of λ , it comes equipped with a universal natural transformation (see 2.1.5)

$$\begin{array}{c} \omega \mathbf{Cat} & \xrightarrow{\lambda} & \mathbf{Ch}_{\geq 0} \\ \gamma^{\mathrm{folk}} & & & & & & \\ \gamma^{\mathrm{folk}} & & & & & & \\ ho(\omega \mathbf{Cat}^{\mathrm{Th}}) & \xrightarrow{\alpha^{\mathrm{pol}}} & ho(\mathbf{Ch}_{\geq 0}). \end{array}$$

Moreover, the functor λ being left Quillen, it is strongly derivable (Definition 2.3.13) and hence induces a morphism of op-prederivators, which we again denote by \mathbb{H}^{pol} :

$$\mathbb{H}^{\mathrm{pol}}: \mathcal{H}o(\omega\mathbf{Cat}^{\mathrm{folk}}) \to \mathcal{H}o(\mathbf{Ch}_{>0}).$$

The following proposition is an immediate consequence of Theorem 2.5.2.

Proposition 4.3.7. *The polygraphic homology*

$$\mathbb{H}^{\mathrm{pol}}:\mathcal{H}\mathrm{o}(\omega\mathbf{Cat}^{\mathrm{folk}})\to\mathcal{H}\mathrm{o}(\mathbf{Ch}_{\geq 0})$$

is homotopy cocontinuous.

As we shall now see, oplax homotopy equivalences (Definition 3.4.2) induce isomorphisms in polygraphic homology. In order to prove that, we first need a couple of technical lemmas.

Lemma 4.3.8. Let

$$\begin{array}{ccc} C' & \xrightarrow{f'_{\epsilon}} & D' \\ \downarrow^{u} & & \downarrow^{v} \\ C & \xrightarrow{f_{\epsilon}} & D \end{array}$$

be commutative squares in ω **Cat** for $\epsilon \in \{0, 1\}$.

If C' is a free ω -category and v a folk trivial fibration, then for every oplax transformation

$$\alpha: f_0 \Rightarrow f_1,$$

there exists an oplax transformation

$$\alpha': f_0' \Rightarrow f_1'$$

such that

$$v \star \alpha' = \alpha \star u.$$

Proof. Notice first that because of the natural isomorphism

$$(\mathbb{D}_0 \amalg \mathbb{D}_0) \otimes C \simeq C \amalg C,$$

we have that $\alpha : f_0 \Rightarrow f_1$ can be encoded in a functor $\alpha : \mathbb{D}_1 \otimes C \to D$ such that the diagram

$$(\mathbb{D}_0 \amalg \mathbb{D}_0) \otimes C \simeq C \amalg C$$
$$\stackrel{i_1 \otimes C}{\longrightarrow} D_1 \otimes C \xrightarrow{\langle u, v \rangle} D$$

(where $i_1 : \mathbb{D}_0 \amalg \mathbb{D}_0 \simeq \mathbb{S}_0 \to \mathbb{D}_1$ is the morphism introduced in 1.1.8) is commutative. Now, the hypotheses of the lemma yield the following commutative square

$$(\mathbb{D}_0 \amalg \mathbb{D}_0) \otimes C' \xrightarrow{\langle f'_0, f'_1 \rangle} D'$$

and since i_1 is a folk cofibration and C' is cofibrant, it follows that the left vertical morphism of the previous square is a folk cofibration (see [Luc17, Proposition 5.1.2.7] or [AL20]). By hypothesis, v is a folk trivial fibration, and so the above square admits a lift

$$\alpha': \mathbb{D}_1 \otimes C' \to D'$$

The commutativity of the two induced triangles shows what we needed to prove. \Box

From now on, for an ω -functor u, we write $\mathbb{H}^{\text{pol}}(u)$ instead of $\mathbb{H}^{\text{pol}}(\gamma^{\text{folk}}(u))$ (where γ^{folk} is the localization functor $\omega \text{Cat} \to \text{ho}(\omega \text{Cat}^{\text{folk}})$) for the morphism induced by u at the level of polygraphic homology.

Lemma 4.3.9. Let $u, v : C \to D$ be two ω -functors. If there exists an oplax transformation $u \Rightarrow v$, then

$$\mathbb{H}^{\mathrm{pol}}(u) = \mathbb{H}^{\mathrm{pol}}(v).$$

Proof. In the case that C and D are both folk cofibrant, this follows immediately from Lemma 4.3.2 and the fact that the localization functor $\mathbf{Ch}_{\geq 0} \to \mathrm{ho}(\mathbf{Ch}_{\geq 0})$ identifies chain homotopic maps (4.3.1).

In the general case, let

$$p: C' \to C$$

and

$$q: D' \to D$$

be folk trivial fibrations with C' and D' cofibrant. Using that q is a trivial fibration and C' is cofibrant, we know that there exist $u' : C' \to D'$ and $v' : C' \to D'$ such that the squares

$$\begin{array}{ccc} C' & \stackrel{u'}{\longrightarrow} & D' & C' & \stackrel{v'}{\longrightarrow} & D' \\ \downarrow^{p} & & \downarrow^{q} & \text{and} & \downarrow^{p} & & \downarrow^{q} \\ C & \stackrel{u}{\longrightarrow} & D & C & \stackrel{v}{\longrightarrow} & D \end{array}$$

are commutative. From Lemma 4.3.8, we deduce the existence of an oplax transformation $u' \Rightarrow v'$. Since C' and D' are cofibrant, we have already proved that

$$\mathbb{H}^{\mathrm{pol}}(u') = \mathbb{H}^{\mathrm{pol}}(v').$$

The commutativity of the two previous squares and the fact that p and q are folk weak equivalences imply the desired result.

The following proposition is an immediate consequence of the previous lemma.

Proposition 4.3.10. Let $u : C \to D$ be an ω -functor. If u is an oplax homotopy equivalence, then the induced morphism

$$\mathbb{H}^{\mathrm{pol}}(u): \mathbb{H}^{\mathrm{pol}}(C) \to \mathbb{H}^{\mathrm{pol}}(D)$$

is an isomorphism.

4.3.11. Oplax homotopy equivalences being particular cases of Thomason equivalences, one may wonder whether it is true that *every* Thomason equivalence induce an isomorphism in polygraphic homology. As we shall see later (Proposition 4.5.3), it is not the case.

Remark 4.3.12. Lemma 4.3.8, Lemma 4.3.9 and Proposition 4.3.10 are also true if we replace "oplax" by "lax" everywhere.

4.4 SINGULAR HOMOLOGY AS DERIVED ABELIANIZA-TION

We have seen in the previous section that the polygraphic homology functor is the total left derived functor of $\lambda : \omega \operatorname{Cat} \to \operatorname{Ch}_{\geq 0}$ when $\omega \operatorname{Cat}$ is equipped with the folk weak equivalences. As it turns out, the abelianization functor is also totally left derivable when $\omega \operatorname{Cat}$ is equipped with the Thomason equivalences and the total left derived functor is the singular homology functor. In order to prove this result, we first need a few technical lemmas.

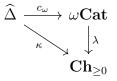
Lemma 4.4.1. Let $\nu : \mathbf{Ch}_{\geq 0} \to \omega \mathbf{Cat}$ be the right adjoint of the abelianization functor (see Lemma 4.2.2). This functor sends the quasi-isomorphisms to Thomason equivalences.

Proof. We have already seen that $\lambda : \omega \operatorname{Cat} \to \operatorname{Ch}_{\geq 0}$ is left Quillen with respect to the folk model structure on $\omega \operatorname{Cat}$. By adjunction, this means that ν is right Quillen for this model structure. In particular, it sends trivial fibrations of chain complexes to folk trivial fibrations. From Ken Brown's Lemma [Hov07, Lemma 1.1.12] and the fact that all chain complexes are fibrant, it follows that ν sends weak equivalences of chain complexes to make equivalences of the folk model structure, which are in particular Thomason equivalences (Lemma 3.6.1).

Remark 4.4.2. The proof of the previous lemma shows the stronger result that ν sends the quasi-ismorphisms to folk weak equivalences. This will be of no use in the sequel.

Recall that we write $c_{\omega} : \widehat{\Delta} \to \omega \mathbf{Cat}$ for the left adjoint of the nerve functor $N_{\omega} : \omega \mathbf{Cat} \to \widehat{\Delta}$ (see Paragraph 3.1.3).

Lemma 4.4.3. *The triangle of functors*



is commutative (up to a canonical isomorphism).

Proof. All the functors involved are cocontinuous, hence it suffices to prove that the triangle is commutative when pre-composed by the Yoneda embedding $\Delta \rightarrow \hat{\Delta}$. This follows immediately from the description of the orientals in [Ste04].

Recall now that the notions of adjunction and equivalence are valid in every 2-category and in particular in the 2-category of pre-derivators (see 2.3.8). We omit the proof of the following lemma, which is the same as when the ambient 2-category is the 2-category of categories.

Lemma 4.4.4. Let $f: y \rightleftharpoons z: g$ be an adjunction and $h: x \to y$ an equivalence with quasi-inverse $k: y \to x$. Then fh is left adjoint to kg.

We can now state and prove the promised result.

Theorem 4.4.5. Consider that $\omega \mathbf{Cat}$ is equipped with the Thomason equivalences. The abelianization functor $\lambda : \omega \mathbf{Cat} \to \mathbf{Ch}_{\geq 0}$ is strongly left derivable and the left derived morphism of op-prederivators

$$\mathbb{L}\lambda^{\mathrm{Th}} : \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\mathbf{Ch}_{>0})$$

is isomorphic to the singular homology

$$\mathbb{H}^{\mathrm{Sing}}: \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\mathbf{Ch}_{\geq 0}).$$

Proof. Let ν be the right adjoint of the abelianization functor (see Lemma 4.2.2) and consider the following adjunctions

$$\widehat{\Delta} \xrightarrow[N_{\omega}]{c_{\omega}} \omega \mathbf{Cat} \xrightarrow[]{\lambda}{} \mathbf{Ch}_{\geq 0}$$

where the functors from left to right are the left adjoints. We know that:

- The functor ν induces a morphism of localizers

$$\nu: (\mathbf{Ch}_{\geq 0}, \mathcal{W}_{\mathbf{Ch}_{\geq 0}}) \to (\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{Th}}),$$

thanks to Lemma 4.4.1.

- The functor N_{ω} induces a morphism of localizers

$$N_{\omega}: (\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{Th}}) \to (\Delta, \mathcal{W}_{\Delta}),$$

by definition of Thomason equivalences.

- There is an isomorphism of functors $\lambda c_{\omega} \simeq \kappa$ (Lemma 4.4.3), hence an induced morphism of localizers

$$(\lambda c_{\omega}) \simeq \kappa : (\overline{\Delta}, \mathcal{W}_{\Delta}) \to (\mathbf{Ch}_{\geq 0}, \mathcal{W}_{\mathbf{Ch}_{\geq 0}}),$$

thanks to Lemma 4.1.4.

It follows that there is an induced adjunction at the level of op-prederivators:

$$\overline{\kappa} \simeq \overline{\lambda c_{\omega}} : \mathcal{H}o(\widehat{\Delta}) \xrightarrow{} \mathcal{H}o(\mathbf{Ch}_{\geq 0}) : \overline{N_{\omega}}\overline{\nu}.$$

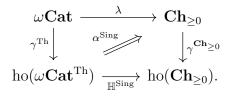
Now, we know from Theorem 3.2.5 that $\overline{N_{\omega}}$ is an equivalence of op-prederivators, and thus admits a quasi-inverse. Let $M : \mathcal{H}o(\widehat{\Delta}) \to \mathcal{H}o(\omega \mathbf{Cat})$ be such a quasi-inverse. From Lemma 4.4.4, we deduce that we have an adjunction:

$$\overline{\kappa}\overline{N_{\omega}}:\mathcal{H}o(\omega\mathbf{Cat}^{\mathrm{Th}}) \longleftrightarrow \mathcal{H}o(\mathbf{Ch}_{\geq 0}):M\overline{N_{\omega}}\overline{\nu}\simeq\overline{\nu}.$$

From Proposition 2.3.15, we conclude that $\lambda : \omega \mathbf{Cat} \to \mathbf{Ch}_{\geq 0}$ is strongly left derivable and that $\mathbb{L}\lambda^{\mathrm{Th}} \simeq \overline{\kappa} \overline{N_{\omega}}$, which is, by definition, the singular homology.

Remark 4.4.6. Beware that neither $c_{\omega} : \widehat{\Delta} \to \omega \mathbf{Cat}$ sends all weak equivalences of simplicial sets to Thomason equivalences nor $\lambda : \omega \mathbf{Cat} \to \mathbf{Ch}_{\geq 0}$ sends all Thomason equivalences to quasi-isomorphisms. However, this does not contradict the fact that $\lambda c_{\omega} : \widehat{\Delta} \to \mathbf{Ch}_{\geq 0}$ does send all weak equivalences of simplicial sets to quasi-isomorphisms.

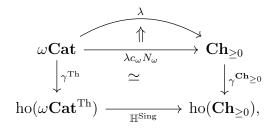
4.4.7. Since \mathbb{H}^{Sing} : ho($\omega \mathbf{Cat}^{\text{Th}}$) \rightarrow ho($\mathbf{Ch}_{\geq 0}$) is the left derived functor of the abelianization functor, it comes with a universal natural transformation



A thorough reading of the proofs of Proposition 2.3.15 and Theorem 4.4.5 enables us to give the following description of α^{Sing} . By post-composing the co-unit of the adjunction $c_{\omega} \dashv N_{\omega}$ with the abelianization functor, we obtain a natural transformation

$$\lambda c_{\omega} N_{\omega} \Rightarrow \lambda$$

Then α^{Sing} is nothing but the following composition of natural transformations



where the square is commutative (up to an isomorphism) because $\mathbb{H}^{\text{Sing}} \simeq \overline{\lambda c_{\omega} N_{\omega}}$.

4.5 COMPARING HOMOLOGIES

4.5.1. Recall from Proposition 3.6.2 that the identity functor on ω Cat induces a morphism of localizers

$$(\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{folk}}) \to (\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{Th}}),$$

which in turn induces a functor

$$\mathcal{J}: \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{folk}}) \to \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}),$$

such that

$$\gamma^{\mathrm{Th}} = \mathcal{J} \circ \gamma^{\mathrm{folk}}.$$

Now, consider the following triangle

$$\begin{array}{ccc} \operatorname{ho}(\omega \mathbf{Cat}^{\mathrm{folk}}) & \xrightarrow{\mathcal{J}} & \operatorname{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) \\ & & & \downarrow_{\mathbb{H}^{\mathrm{Sing}}} \\ & & & \operatorname{ho}(\mathbf{Ch}_{>0}). \end{array}$$

$$(4.1)$$

A natural question to ask is whether this triangle is commutative (up to an isomorphism). Since \mathcal{J} is the identity on objects, this amounts to ask whether for every ω -category C we have an isomorphism (natural in C)

$$\mathbb{H}^{\mathrm{pol}}(C) \simeq \mathbb{H}^{\mathrm{Sing}}(C).$$

As it happens, this is not possible as the following counter-example, due to Ara and Maltsiniotis, shows.

4.5.2 (Ara and Maltsiniotis' counter-example). Write $\mathbb{N} = (\mathbb{N}, +, 0)$ for the commutative monoid of non-negative integers and let C be the 2-category defined as

$$C := B^2 \mathbb{N}$$

(see 1.3.1). As usual, we consider C as an ω -category with only unit cells strictly above dimension 2. This ω -category is free; namely its k-basis is a singleton for k = 0 and k = 2, and the empty set otherwise. In particular C is cofibrant for the folk model structure (Proposition 3.5.9) and it follows from Proposition 4.2.8 that $\mathbb{H}^{\text{pol}}(C)$ is given by the chain complex (seen as an object of $\text{ho}(\mathbf{Ch}_{>0})$)

$$\mathbb{Z} \leftarrow 0 \leftarrow \mathbb{Z} \leftarrow 0 \leftarrow 0 \leftarrow \cdots$$

Hence, the polygraphic homology groups of B are given by

$$H_k^{\text{pol}}(C) = \begin{cases} \mathbb{Z} \text{ if } k = 0, 2\\ 0 \text{ in other cases.} \end{cases}$$

On the other hand, it is proven in [Ara19, Theorem 4.9 and Example 4.10] that (the nerve of) C is a $K(\mathbb{Z}, 2)$. In particular, it has non-trivial singular homology groups in every even dimension. This proves that $\mathbb{H}^{\text{pol}}(C)$ is *not* isomorphic to $\mathbb{H}^{\text{Sing}}(C)$; which means that triangle (4.1) cannot be commutative (up to an isomorphism).

Another consequence of the above counter-example is the following result, which we claimed in 4.3.11. Recall that given a morphism $u : C \to D$ of ω Cat, we write $\mathbb{H}^{\text{pol}}(u)$ instead of $\mathbb{H}^{\text{pol}}(\gamma^{\text{folk}}(u))$.

Proposition 4.5.3. There exists at least one Thomason equivalence

$$u:C\to D$$

such that the induced morphism

$$\mathbb{H}^{\mathrm{pol}}(u):\mathbb{H}^{\mathrm{pol}}(C)\to\mathbb{H}^{\mathrm{pol}}(D)$$

is not an isomorphism of $ho(Ch_{>0})$.

Proof. Suppose the converse, which is that the functor

¢.

$$\mathbb{H}^{\mathrm{pol}} \circ \gamma^{\mathrm{folk}} : \omega \mathbf{Cat} \to \mathrm{ho}(\mathbf{Ch}_{>0})$$

sends the Thomason equivalences to isomorphisms of $ho(\mathbf{Ch}_{\geq 0})$. Because of the inclusion $\mathcal{W}^{\text{folk}} \subseteq \mathcal{W}^{\text{Th}}_{\omega}$, the category $ho(\omega \mathbf{Cat}^{\text{Th}})$ may be identified with the localization of $ho(\omega \mathbf{Cat}^{\text{folk}})$ with respect to $\gamma^{\text{folk}}(\mathcal{W}^{\text{Th}}_{\omega})$ and then the localization functor is nothing but

$$\mathcal{T}: \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{folk}}) \to \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}).$$

From this observation and because of the hypothesis we made on Thomason equivalences inducing isomorphisms in polygraphic homology, we deduce the existence of a functor

$$\mathbb{H}^{\mathrm{pol}}$$
: ho($\omega \mathbf{Cat}^{\mathrm{Th}}$) \rightarrow ho($\mathbf{Ch}_{>0}$)

such that we have

$$\overline{\mathbb{H}^{\mathrm{pol}}} \circ \mathcal{J} = \mathbb{H}^{\mathrm{pol}},$$

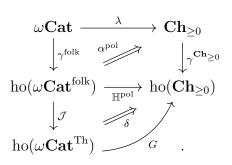
and because of the equality $\gamma^{Th} = \mathcal{J} \circ \gamma^{folk}$, the universal natural transformation α^{pol} now reads

$$\begin{split} \omega \mathbf{Cat} & \xrightarrow{\lambda} \mathbf{Ch}_{\geq 0} \\ & \downarrow^{\gamma^{\mathrm{Th}}} & \stackrel{\alpha^{\mathrm{pol}}}{\longrightarrow} & \downarrow^{\gamma^{\mathbf{Ch}}_{\geq 0}} \\ & \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) & \xrightarrow{\mathbb{H}^{\mathrm{pol}}} & \mathrm{ho}(\mathbf{Ch}_{\geq 0}). \end{split}$$

Let us show that $(\overline{\mathbb{H}^{\text{pol}}}, \alpha^{\text{pol}})$ is the left derived functor of λ when ω Cat is equipped with the Thomason equivalences. Let G and β be as in the following 2-diagram

$$\begin{array}{ccc} \omega \mathbf{Cat} & \xrightarrow{\lambda} & \mathbf{Ch}_{\geq 0} \\ \gamma^{\mathrm{Th}} = \mathcal{J} \circ \gamma^{\mathrm{folk}} & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & &$$

Since \mathbb{H}^{pol} is the left derived functor of λ when $\omega \mathbf{Cat}$ is equipped with the folk weak equivalences, there exists a unique $\delta : G \circ \mathcal{J} \Rightarrow \mathbb{H}^{\text{pol}}$ that factorizes β as



But since \mathcal{J} acts as a localization functor, δ also factorizes uniquely as

$$\mathrm{ho}(\omega\mathbf{Cat}^{\mathrm{folk}}) \xrightarrow{\mathcal{J}} \mathrm{ho}(\omega\mathbf{Cat}^{\mathrm{Th}}) \underbrace{\mathfrak{s}'}_{G} \stackrel{\mathrm{H}^{\mathrm{pol}}}{\overset{\mathrm{ho}}{\overset{\mathrm{ho}}}} (\mathbf{Ch}_{\geq 0}).$$

Altogether we have that β factorizes as

$$\begin{array}{ccc}
\omega \mathbf{Cat} & \xrightarrow{\lambda} & \mathbf{Ch}_{\geq 0} \\
\downarrow^{\gamma^{\mathrm{Th}}} & \stackrel{\alpha^{\mathrm{pol}}}{\longrightarrow} & \downarrow^{\gamma^{\mathbf{Ch}}_{\geq 0}} \\
\mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) & \xrightarrow{\overline{\mathbb{H}^{\mathrm{pol}}}} & \mathrm{ho}(\mathbf{Ch}_{\geq 0}). \\
\overbrace{G}^{\delta' \uparrow \uparrow} & \overbrace{G}^{\bullet} & \overbrace{G}^{\bullet}
\end{array}$$

The uniqueness of such a factorization follows from a similar argument which is left to the reader. This proves that $\overline{\mathbb{H}^{\text{pol}}}$ is the left derived functor of λ when ω Cat is equipped with the Thomason equivalences and in particular we have

$$\mathbb{H}^{\mathrm{Sing}} \simeq \overline{\mathbb{H}^{\mathrm{pol}}}.$$

But since \mathcal{J} is the identity on objects, this implies that for every ω -category C we have

$$\mathbb{H}^{\mathrm{Sing}}(C) \simeq \overline{\mathbb{H}^{\mathrm{pol}}}(C) = \mathbb{H}^{\mathrm{pol}}(C),$$

which we know is impossible.

Remark 4.5.4. It follows from the previous result that if we think of ω -categories as a model for homotopy types (see Theorem 3.2.5), then the polygraphic homology of an ω -category is *not* a well defined invariant. This justifies what we said in remark 4.1.9, which is that *singular homology* is the only "correct" homology of ω -categories.

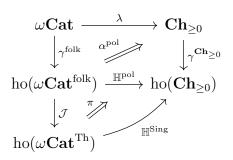
4.5.5. Even though triangle (4.1) is not commutative (even up to an isomorphism), it can be filled up with a 2-morphism. Indeed, consider the following 2-square

$$\begin{split} & \omega \mathbf{Cat} \xrightarrow{\lambda} \mathbf{Ch}_{\geq 0} \\ & & \downarrow^{\gamma^{\mathrm{Th}}} \xrightarrow{\alpha^{\mathrm{Sing}}} & \downarrow^{\gamma^{\mathbf{Ch}}_{\geq 0}} \\ & & \mathrm{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) \xrightarrow{\mathbb{H}^{\mathrm{Sing}}} \mathrm{ho}(\mathbf{Ch}_{\geq 0}), \end{split}$$

Since $\gamma^{\text{Th}} = \mathcal{J} \circ \gamma^{\text{folk}}$ and the polygraphic homology is the total left derived functor of the abelianization functor when ω **Cat** is equipped with the folk weak equivalences, we obtain by universal property (see 2.1.5) a unique natural transformation

$$\begin{array}{c} \operatorname{ho}(\omega \mathbf{Cat}^{\mathrm{folk}}) \\ \downarrow_{\mathcal{J}} & \xrightarrow{\pi} & \\ \operatorname{ho}(\omega \mathbf{Cat}^{\mathrm{Th}}) & \xrightarrow{\pi} & \operatorname{ho}(\mathbf{Ch}_{\geq 0}) \end{array}$$
(4.2)

such that α^{Sing} factorizes as



Since \mathcal{J} is nothing but the identity on objects, for every ω -category C, the natural transformation π yields a map

$$\pi_C : \mathbb{H}^{\mathrm{Sing}}(C) \to \mathbb{H}^{\mathrm{pol}}(C),$$

which we shall refer to as the canonical comparison map.

Remark 4.5.6. When C is free, it follows from the considerations in 4.4.7 that the canonical comparison map π_C can be identified with the image by $\gamma^{\mathbf{Ch}_{\geq 0}}$ of the morphism of $\mathbf{Ch}_{\geq 0}$

$$\lambda c_{\omega} N_{\omega}(C) \to \lambda(C)$$

induced by the co-unit of $c_{\omega} \dashv N_{\omega}$.

Definition 4.5.7. An ω -category C is said to be *homologically coherent* when the canonical comparison map

$$\pi_C : \mathbb{H}^{\mathrm{Sing}}(C) \to \mathbb{H}^{\mathrm{pol}}(C)$$

is an isomorphism of $ho(Ch_{>0})$.

4.5.8. The rest of this dissertation is devoted to the study of homologically coherent ω -categories. Examples of such ω -categories will be presented later. Following the perspective of Remark 4.5.4, polygraphic homology can be thought of as a way to compute singular homology of homologically coherent ω -categories.

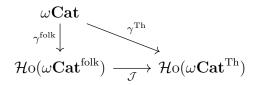
4.5.9. Similarly to 4.5.1, the morphism of localizers

$$(\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{folk}}) \to (\omega \mathbf{Cat}, \mathcal{W}^{\mathrm{Th}})$$

induces a morphism of op-prederivators

$$\mathcal{J}: \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{folk}}) \to \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}})$$

such that the triangle in the category of op-prederivators



is commutative. It follows from what we said in 3.6.4 that the morphism \mathcal{J} cannot be an equivalence of op-prederivators. As we shall see later, \mathcal{J} is not even homotopy cocontinuous. In particular, this implies that given a diagram $d: I \to \omega Cat$, the canonical arrow of $ho(\omega Cat^{Th})$

$$\operatorname{hocolim}_{I}^{\operatorname{Th}}(\mathcal{J}_{I}(d)) \to \mathcal{J}_{e}(\operatorname{hocolim}_{I}(d))$$
(4.3)

induced by \mathcal{J} (see 2.3.4) is generally *not* an isomorphism. Note that since

$$\mathcal{J}_I : \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{folk}})(I) \to \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}})(I)$$

is the identity on objects for every small category I, morphism (4.3) simply reads

$$\operatorname{hocolim}_{I}^{\operatorname{Th}}(d) \to \operatorname{hocolim}_{I}^{\operatorname{folk}}(d).$$

Even if this is not always true, there are some particular diagrams d for which the above morphism is indeed an isomorphism. The criterion to find homologically coherent ω -categories given in the proposition below is based on this observation.

Proposition 4.5.10. Let C be an ω -category. Suppose that there exists a diagram

$$d: I \to \omega \mathbf{Cat}$$

and a co-cone

$$(\varphi_i: d(i) \to C)_{i \in \mathrm{Ob}(I)}$$

such that:

- (i) For every $i \in Ob(I)$, the ω -category d(i) is homologically coherent.
- (ii) The canonical morphism

 $\operatorname{hocolim}^{\operatorname{folk}} d \to C$

is an isomorphism of $ho(\omega Cat^{folk})$.

(iii) The canonical morphism

 $\operatorname{hocolim}^{\operatorname{Th}} d \to C$

is an isomorphism of $ho(\omega Cat^{Th})$.

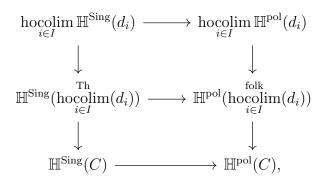
Then the ω -category C is homologically coherent.

Proof. Notice first that all the constructions from 4.5.5 may be reproduced *mutatis mutandis* at the level of op-prederivators. In particular, we obtain a 2-morphism of op-prederivators

$$\mathcal{H}o(\omega\mathbf{Cat}^{\text{folk}})$$

$$\downarrow^{\mathcal{J}} \xrightarrow{\pi} \mathcal{H}o(\omega\mathbf{Cat}^{\text{Th}}) \xrightarrow{\mathbb{H}^{\text{pol}}} \mathcal{H}o(\mathbf{Ch}_{\geq 0}).$$

Then, by naturality, we have a commutative diagram in $ho(Ch_{>0})$:



where:

- the top and bottom horizontal arrows are induced by π ,
- the middle horizontal arrow is induced by π and the canonical morphism

$$\operatorname{hocolim}_{i \in I}^{\operatorname{Th}}(d_i) \to \operatorname{hocolim}_{i \in I}^{\operatorname{folk}}(d_i)$$

from 4.5.9,

- the top vertical arrows are the canonical morphisms induced by every morphism of op-prederivators (see 2.3.4),
- the bottom vertical arrows are induced by the co-cone

$$(\varphi_i: d(i) \to C)_{i \in \mathrm{Ob}(I)}.$$

Since \mathbb{H}^{pol} and \mathbb{H}^{Sing} are both homotopy cocontinuous (Proposition 4.1.12 and Proposition 4.3.7 respectively), both top vertical arrows are isomorphisms. Hypotheses (*ii*) and (*iii*) imply that the bottom vertical arrows are isomorphisms and hypothesis (*i*) imply that the top horizontal arrow is an isomorphism. By a 2-out-of-3 property, the bottom horizontal arrow is an isomorphism, which means exactly that *C* is homologically coherent.

The previous proposition admits the following corollary, which will be of great use in later chapters.

Corollary 4.5.11. Let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ & & \downarrow^{f} & & \downarrow^{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array}$$

be a cocartesian square of ω Cat such that:

- (a) the ω -categories A, B and C are free and homologically coherent,
- (b) at least one of the morphisms $u : A \to B$ or $f : A \to C$ is a folk cofibration,
- (c) the square is Thomason homotopy cocartesian.

Then, the ω -category D is homologically coherent.

Proof. The fact that A,B and C are free and one of the morphism u or f is a folk cofibration ensure that the square is folk homotopy cocartesian (Lemma 2.5.12). The conclusion follows then from Proposition 4.5.10.

4.6 EQUIVALENCE OF HOMOLOGIES IN LOW DIMEN-SION

4.6.1. Recall that for every $n \ge 0$ we have taken the habit of identifying nCat as a full subcategory of ω Cat via the canonical fully faithful functor $\iota_n : n$ Cat $\rightarrow \omega$ Cat (defined in 1.1.4) that sends an *n*-category *C* to the ω -category with the same *k*-cells as *C* for $k \le n$ and only unit cells for k > n. In particular, we abusively wrote

$$C = \iota_n(C).$$

Within this section, and only within this section, we try not to make this abuse of notation and explicitly write ι_n whenever we should.

We have already seen that ι_n has a left adjoint $\tau_{\leq n}^i : \omega \mathbf{Cat} \to n\mathbf{Cat}$, where for an ω -category C, $\tau_{\leq n}^i(C)$ is the *n*-category whose set of *k*-cells is C_k for k < n and whose set of *n*-cells is the quotient of C_n by the equivalence relation \sim generated by

 $x \sim y$ when there exists $z : x \to y$ in C_{n+1} .

As it happens, we can use the adjunction $\tau_{\leq n}^i \dashv \iota_n$ to transport the folk model structure from ω Cat to *n*Cat.

Proposition 4.6.2. There exists a model structure on nCat such that:

- the weak equivalences are exactly those morphisms $f : C \to D$ such that $\iota_n(f)$ is a weak equivalence for the folk model structure on ωCat ,
- the fibrations are exactly those morphisms $f : C \to D$ such that $\iota_n(f)$ is a fibrations for the folk model structure on ω Cat.

Moreover, there exists a set I of generating cofibrations (resp. a set J of generating trivial cofibrations) for the folk model structure on $\omega \mathbf{Cat}$ such that the image by $\tau_{\leq n}^i$ of I (resp. J) is a set of generating cofibrations (resp. generating trivial cofibrations) of the above model structure on $n\mathbf{Cat}$.

Proof. This is [LMW10, Theorem 5]. (Although the part concerning generating cofibrations is not made explicit in the statement of the theorem, it is contained in proof.) \Box

4.6.3. We refer to the model structure of the above proposition as the *folk model struc*ture on *n*Cat. By definition, the functor $\iota_n : n$ Cat $\to \omega$ Cat preserves weak equivalences and fibrations when ω Cat and *n*Cat are equipped with the folk model structure. In particular, the adjunction $\tau_{\leq n}^i \dashv \iota_n$ is a Quillen adjunction. As it happens, the functor $\tau_{\leq n}^i$ also preserves weak equivalences.

Proposition 4.6.4. The functor $\tau_{\leq n}^i : \omega \mathbf{Cat} \to n\mathbf{Cat}$ sends the weak equivalences of the folk model structure on $\omega \mathbf{Cat}$ to weak equivalences of the folk model structure on $n\mathbf{Cat}$.

Proof. Since every ω -category is fibrant for the folk model structure on ω Cat [LMW10, Proposition 9], it suffices to show that $\tau_{\leq n}^i$ sends the folk trivial fibrations of ω Cat to weak equivalences of *n*Cat (in virtue of Ken Brown's Lemma [Hov07, Lemma 1.1.12]).

For convenience, let us write T for the functor $\iota_n \tau_{\leq n}^i : \omega \mathbf{Cat} \to \omega \mathbf{Cat}$. By definition of folk weak equivalences on $n\mathbf{Cat}$, we have to show that for every folk trivial fibration $f: C \to D$ of $\omega \mathbf{Cat}$, the ω -functor

$$T(f): T(C) \to T(D)$$

is a folk weak equivalence on ωCat . Consider the following commutative square

$$C \xrightarrow{f} D$$

$$\downarrow \eta_C \qquad \qquad \downarrow \eta_D$$

$$T(C) \xrightarrow{T(f)} T(D),$$

where η is the unit of the adjunction $\tau_{< n}^i \dashv \iota_n$.

Let us first treat the case of 0-cells. Let y be a 0-cell of T(D). The map η_D being surjective on 0-cells (even if n = 0), there exists y' such that $\eta_D(y') = y$. Since f is a folk trivial fibration, there exists a 0-cell x' of C such that f(x') = y' and then if we set $x := \eta_C(x')$, we have T(f)(x) = y.

Now let x, y be parallel k-cells of T(C) and let $\beta : f(x) \to f(y)$ be a (k + 1)-cell of T(D). We have to distinguish several cases.

Case k < n - 1: Since η_C and η_D are identities on k-cells for every $0 \le k < n$ and since f is a folk trivial fibration, there exists a (k + 1)-cell $\alpha : x \to y$ of T(C) such that

$$T(f)(\alpha) = \beta.$$

- **Case** k = n 1: By definition of T(D), there exists an *n*-cell $\beta' : f(x) \to f(y)$ of D such that $\eta_D(\beta') = \beta$. Since f is a folk trivial fibration, there exists an *n*-cell $\alpha' : x \to y$ of C such that $f(\alpha') = \beta'$. If we set $\alpha := \eta_C(\alpha')$, we have $T(f)(\alpha) = \beta$.
- **Case** k = n: Since all *l*-cells of T(C) and T(D) with l > n are units, we trivially have that f(x) = f(y) and β is the unit on f(x). Now let x' and y' be parallel *n*-cells of *C* such that $\eta_C(x') = x$ and $\eta_C(y') = y$ (this is always possible by definition of T(C)). We have $\eta_D(f(x')) = f(x) = f(y) = \eta_D(f(y'))$. By definition of the functor $\tau_{\leq n}^i$, this means that there exists a zigzag of (n + 1)-cells of *D* from f(x') to f(y'). More precisely, this means that there exists a sequence

$$(z_0,\beta_1,z_1,\cdots,z_{p-1},\beta_p,z_p)$$

where the z_i are all parallel *n*-cells of D with $z_0 = f(x')$ and $z_p = f(y')$, and each β_i is (n+1)-cell of D either from z_{i-1} to z_i or from z_i to z_{i-1} . Using the fact that f is a folk trivial fibration, it is easy to prove the existence of a zigzag from x' to y', which implies in particular that $x = \eta_C(x') = \eta_C(y') = y$.

Case k > n: Since all k-cells of T(C) and T(D) with k > n are units, we trivially have f(x) = f(y) (and β is the unit on f(x)) and x = y.

Altogether, this proves that T(f) is a folk trivial fibration, hence a folk weak equivalence.

For later reference, we put here the following lemma.

Lemma 4.6.5. If an *n*-category *C* has a *k*-basis for every $0 \le k \le n - 1$, then it is cofibrant for the folk model structure on *n*Cat.

Proof. Since $\tau_{\leq n}^i$ is a left Quillen functor, it suffices in virtue of Proposition 3.5.9 to show that there exists a free ω -category C' such that $\tau_{\leq n}^i(C') = C$. First, consider the *n*-category $(U_{n-1}(C))^*$ (for the notations, see 1.4.4 and 1.4.10). This *n*-category has the same *k*-cells as *C* for k < n and has exactly one *generating n*-cell for each *n*-cell of *C*. It is obviously free and we have a canonical *n*-functor

$$\epsilon_C : (U_{n-1}(C))^* \to C,$$

given by the co-unit of the adjunction $(-)^* \dashv U_{n-1}(-)$. Now, let C' be the (n+1)-category (considered as an ω -category) that has the same k-cells as $(U_{n-1}(C))^*$ for $k \leq n$ and whose set of (n+1)-cells is freely generated be the set

 $\{(x, y) \mid x \text{ and } y \text{ are parallel } n \text{-cells of } (U_{n-1}(C))^* \text{ such that } \epsilon_C(x) = \epsilon_C(y) \}.$

The (n + 1)-category C' is obviously free and it is a harmless verification, which we leave to reader, to check that $\tau_{< n}^i(C') = C$.

Example 4.6.6. Every (small) category is cofibrant for the folk model structure on Cat.

We now turn to truncations of chain complexes.

4.6.7. Let $\mathbf{Ch}_{\geq 0}^{\leq n}$ be the category of chain complexes concentrated in degrees between 0 and *n*. This means that an object *K* of $\mathbf{Ch}_{\geq 0}^{\leq n}$ is a diagram of abelian groups of the form

$$K_0 \xleftarrow{\partial} K_1 \xleftarrow{\partial} K_2 \xleftarrow{\partial} \cdots \xleftarrow{\partial} K_n,$$

where $\partial \circ \partial = 0$, and morphisms of $\mathbf{Ch}_{\geq 0}^{\leq n}$ are defined the expected way. We write $\iota_n : \mathbf{Ch}_{\geq 0}^{\leq n} \to \mathbf{Ch}_{\geq 0}$ for the canonical functor that sends an object K of $\mathbf{Ch}_{\geq 0}^{\leq n}$ to the chain complex

$$K_0 \xleftarrow{\partial} K_1 \xleftarrow{\partial} K_2 \xleftarrow{\partial} \cdots \xleftarrow{\partial} K_n \xleftarrow{\partial} 0 \xleftarrow{\partial} \cdots \xleftarrow{\partial} K_n$$

This functor is fully faithful and $\mathbf{Ch}_{\geq 0}^{\leq n}$ may be identified with the full subcategory of $\mathbf{Ch}_{\geq 0}$ spanned by chain complexes K such that $K_k = 0$ for every k > n.

Similarly to the case of *n*-categories, the functor $\iota_n : \mathbf{Ch}_{\geq 0}^{\leq n} \to \mathbf{Ch}_{\geq 0}$ has a left adjoint $\tau_{\leq n}^i : \mathbf{Ch}_{\geq 0} \to \mathbf{Ch}_{\geq 0}^{\leq}$, where for a chain complex K, $\tau_{\leq n}^i(K)$ is the object of $\mathbf{Ch}_{\geq 0}^{\leq n}$ defined as

$$K_0 \xleftarrow{\partial} K_1 \xleftarrow{\partial} K_{n-1} \xleftarrow{\partial} K_{n-1}$$

Again, as with n-categories, we can use the adjunction

to create a model structure on $\mathbf{Ch}_{\geq 0}^{\leq n}$.

Proposition 4.6.8. There exists a model structure on $\mathbf{Ch}_{\geq 0}^{\leq n}$ such that:

- the weak equivalences are exactly those morphisms $f : K \to K'$ such that $\iota_n(f)$ is a weak equivalence for the projective model structure on $\mathbf{Ch}_{\geq 0}$,
- the fibrations are exactly those morphisms $f : K \to K'$ such that $\iota_n(f)$ is a fibration for the projective model structure on $\mathbf{Ch}_{>0}$.

Proof. This is a typical example of a transfer of a cofibrantly generated model structure along a right adjoint as in [Bek01, Proposition 2.3]. Since the weak equivalences of the projective model structure on $Ch_{\geq 0}$ are closed under filtered colimits [Wei95, Theorem 2.6.15], the only *a priori* non-trivial hypothesis to check is that there exists a set J of generating trivial cofibrations of the projective model structure on $Ch_{\geq 0}$ such that for every $j: A \to B$ in J and every cocartesian square

$$\begin{array}{ccc} \tau^{i}_{\leq n}(A) & \longrightarrow & X \\ \tau^{i}_{\leq n}(j) & & & \downarrow^{g} \\ \tau^{i}_{\leq n}(B) & \longrightarrow & Y, \end{array}$$

the morphism $\iota_n(g)$ is a weak equivalence of $\mathbf{Ch}_{\geq 0}$. As explained in [DS95, Proposition 7.19], there exists a set of generating trivial cofibrations of the projective model structure on $\mathbf{Ch}_{\geq 0}$ consisting of the maps

$$0 \to D_k$$

for each k > 0, where D_k is the following chain complex concentrated in degree k and k - 1

$$0 \leftarrow \dots \leftarrow 0 \leftarrow \mathbb{Z} \stackrel{\mathrm{id}}{\leftarrow} \mathbb{Z} \leftarrow 0 \leftarrow \dots$$

What is left to show then is that for every k > 0 and every object X of $\mathbf{Ch}_{\geq 0}^{\leq n}$, the canonical inclusion map

$$X \to X \oplus \tau^i_{< n}(D_k)$$

is sent by ι_n to a weak equivalence of $\mathbf{Ch}_{\geq 0}$. This follows immediately from the fact that homology groups commute with direct sums.

4.6.9. We refer to the model structure of the previous proposition as the *projective model* structure on $\mathbf{Ch}_{\geq 0}^{\leq n}$.

Lemma 4.6.10. For every chain complex K, the unit map

$$K \to \iota_n \tau^i_{< n}(K)$$

induces isomorphisms

$$H_k(K) \simeq H_k(\iota_n \tau^i_{< n}(K))$$

for every $0 \le k \le n$.

Proof. For $0 \le k < n-1$, this is trivial. For k = n-1, this follows easily from the fact that the image of $\partial : K_k / \partial(K_{k+1}) \to K_{k-1}$ is equal to the image of $\partial : K_k \to K_{k-1}$. Finally for k = n, it is straightforward to check that

$$H_n(K) = \frac{\operatorname{Ker}(\partial : K_n \to K_{n-1})}{\operatorname{Im}(\partial : K_{n+1} \to K_n)}$$

is isomorphic to

$$H_n(\iota_n \tau^i_{< n}(K)) = \operatorname{Ker}(\partial : K_n / \partial(K_{n+1}) \to K_{n-1}).$$

The isomorphism being obviously induced by the unit map $K \to \iota_n \tau^i_{\leq n}(K)$.

As a consequence of this lemma, we have the analogous of Proposition 4.6.4.

Proposition 4.6.11. The functor $\tau_{\leq n}^i : \mathbf{Ch}_{\geq 0} \to \mathbf{Ch}_{\geq 0}^{\leq n}$ sends the weak equivalences of the projective model structure on $\mathbf{Ch}_{\geq 0}$ to weak equivalences of the projective model structure on $\mathbf{Ch}_{\geq 0}^{\leq n}$.

Proof. Let $f : K \to K'$ be a weak equivalence for the projective model structure on $Ch_{\geq 0}$ and consider the naturality square

$$K \xrightarrow{f} K' \downarrow_{\eta_{K}} \downarrow_{\eta'_{K}} \downarrow_{\eta'_{K}} \downarrow_{\eta'_{K}} \downarrow_{\eta'_{K}} \iota_{n}\tau^{i}_{\leq n}(K) \xrightarrow{\iota_{n}\tau^{i}_{\leq n}(f)} \iota_{n}\tau^{i}_{\leq n}(K'),$$

where η is the unit map of the adjunction $\tau_{<n}^i \dashv \iota_n$. It follows from Lemma 4.6.10 that

$$H_k(\iota_n \tau^i_{\leq n}(f)) : H_k(\iota_n \tau^i_{\leq n}(K)) \to H_k(\iota_n \tau^i_{\leq n}(K'))$$

is an isomorphism for every $k \leq n$. Since obviously $H_k(\iota_n \tau^i_{\leq n}(f))$ is also an isomorphism for k > n, this proves the result.

We now investigate the relation between truncation and abelianization.

4.6.12. Let C be n-category. A straightforward computation shows that the chain complex $\lambda(\iota_n(C))$ is such that

$$\lambda_k(\iota_n(C)) = 0$$

for every k > n and thus $\lambda(\iota_n(C))$ can be seen as an object of $\mathbf{Ch}_{\geq 0}^{\leq n}$. Hence, we can define a functor $\lambda_{\leq n} : n\mathbf{Cat} \to \mathbf{Ch}_{\geq 0}^{\leq n}$ as

$$\lambda_{\leq n} : n\mathbf{Cat} \to \mathbf{Ch}_{\geq 0}^{\leq n}$$
$$C \mapsto \lambda(\iota_n(C)),$$

and we tautologically have that the square

$$n\mathbf{Cat} \xrightarrow{\lambda_{\leq n}} \mathbf{Ch}_{\geq 0}^{\leq n}$$
$$\downarrow^{\iota_n} \qquad \qquad \downarrow^{\iota_n}$$
$$\omega\mathbf{Cat} \xrightarrow{\lambda} \mathbf{Ch}_{>0}$$

is commutative.

Lemma 4.6.13. The square

$$egin{aligned} & \omega \mathbf{Cat} & \stackrel{\lambda}{\longrightarrow} & \mathbf{Ch}_{\geq 0} \ & & & \downarrow^{ au^i_{\leq n}} & & \downarrow^{ au^i_{\leq n}} \ & & n\mathbf{Cat} & \stackrel{\lambda_{\leq n}}{\longrightarrow} & \mathbf{Ch}^{\leq n}_{\geq 0} \end{aligned}$$

is commutative (up to a canonical isomorphism).

Proof. Notice first that we have a natural transformation

$$\beta:\tau^i_{\leq n}\circ\lambda\Rightarrow\lambda_{\leq n}\circ\tau^i_{\leq n}$$

defined as

$$\omega \mathbf{Cat} \xrightarrow{\tau_{\leq n}^{*}} n\mathbf{Cat} \xrightarrow{\lambda_{\leq n}} \mathbf{Ch}_{\geq 0}^{\leq n}$$

$$\underset{id}{\overset{\eta}{\longrightarrow}} \downarrow_{\iota_{n}} \qquad \underset{\iota_{n}}{\overset{\iota_{n}}{\longrightarrow}} \underbrace{\tau_{\epsilon_{n}}^{*}}_{\overset{id}{\longrightarrow}} \mathbf{Ch}_{\geq 0}^{\leq n}.$$

Since for every ω -category C and every k < n, we have

$$C_k = \iota_n(\tau^i_{\le n}(C))_k$$

and for every chain complex K and every k < n, we have

$$\iota_n(\tau^i_{< n}(K))_k = K_k,$$

it follows that for every ω -category C and every k < n, the morphism β_k is nothing but the equality

$$\tau^i_{\leq n}(\lambda(C))_k = \lambda_{\leq n}(\tau^i_{\leq n}(C))_k.$$

Hence, all we have to prove is that

$$\beta_n : \tau^i_{\leq n}(\lambda(C))_n \to \lambda_{\leq n}(\tau^i_{\leq n}(C))_n$$

is an isomorphism for every ω -category C.

Recall from Lemma 4.2.4 that $\lambda_n \circ \iota_n : n\mathbf{Cat} \to \mathbf{Ab}$ (which we abusively wrote as λ_n) is left adjoint to the functor $B^n : \mathbf{Ab} \to n\mathbf{Cat}$. In particular, for every ω -category C and every abelian group G, we have

$$\operatorname{Hom}_{\mathbf{Ab}}(\lambda_n\iota_n\tau^i_{\leq n}(C),G)\simeq\operatorname{Hom}_{n\mathbf{Cat}}(\tau^i_{\leq n}(C),B^nG)\simeq\operatorname{Hom}_{\omega\mathbf{Cat}}(C,\iota_n(B^nG)).$$

Then, it follows from an argument similar to the proof of Lemma 1.3.3 (see also the proof of Lemma 4.2.4) that this last set is naturally isomorphic to the set of functions $f_n: C_n \to G$ such that:

- for every $0 \le k < n$ and every pair (x, y) of k-composable n-cells of C, we have

$$f(x * y) = f(x) + f(y),$$

- for every (n+1)-cell $\alpha : x \to y$ of C, we have

$$f(x) = f(y).$$

By definition of $\lambda_n(C)$ and of $\partial : \lambda_{n+1}(C) \to \lambda_n(C)$, we thus have

$$\operatorname{Hom}_{\mathbf{Ab}}(\lambda_n \iota_n \tau^i_{\leq n}(C), G) \simeq \operatorname{Hom}_{\mathbf{Ab}}(\lambda_n(C)/\partial(\lambda_{n+1}(C)), G).$$

Hence, we have $\lambda_{\leq n}(\tau_{\leq n}^{i}(C))_{n} \simeq \tau_{\leq n}^{i}(\lambda(C))_{n}$ and a thorough analysis of naturality shows that this isomorphism is nothing but β_{n} .

In the following lemma, nCat is equipped with the folk model structure and $Ch_{\geq 0}^{\geq n}$ with the projective model structure.

Lemma 4.6.14. The functor $\lambda_{\leq n} : n\mathbf{Cat} \to \mathbf{Ch}_{\geq 0}^{\leq n}$ is left Quillen.

Proof. Let I and J respectively be sets of generating cofibrations and generating trivial cofibrations of the folk model structure on ω Cat such that $\tau_{\leq n}^i(I)$ and $\tau_{\leq n}^i(J)$ respectively are sets of generating cofibrations and generating trivial cofibrations of the projective model structure on nCat (which we know exist by the second part of Proposition 4.6.2). What we have to show is that for every f in I (resp. J), $\lambda_{\leq n}(\tau_{\leq n}^i(f))$ is a cofibration (resp. generating cofibration) for the folk model structure on $\mathbf{Ch}_{\geq 0}^{\leq n}$. From Lemma 4.6.13, we have

$$\lambda_{\leq n}(\tau^i_{\leq n}(f)) \simeq \tau^i_{\leq n}(\lambda(f)).$$

Since λ and $\tau_{\leq n}^i$ are both left Quillen functors, this proves the result.

As an immediate consequence of the previous lemma, the functor $\lambda_{\leq n}$ is left derivable and we have the following key result.

Proposition 4.6.15. The square

$$\begin{array}{c} \operatorname{ho}(\omega \mathbf{Cat}^{\operatorname{folk}}) \xrightarrow{\mathbb{L}\lambda} \operatorname{ho}(\mathbf{Ch}_{\geq 0}) \\ & \downarrow^{\overline{\tau_{\leq n}^{i}}} & \downarrow^{\overline{\tau_{\leq n}^{i}}} \\ \operatorname{ho}(n\mathbf{Cat}^{\operatorname{folk}}) \xrightarrow{\mathbb{L}\lambda_{\leq n}} \operatorname{ho}(\mathbf{Ch}_{\geq 0}^{\leq n}) \end{array}$$

is commutative (up to a canonical isomorphism).

Proof. Straightforward consequence of Lemma 4.6.13 and the fact that the left derived functor of a composition of left Quillen functors is the composition of the left derived functors (see for example [Hov07, Theorem 1.3.7]). \Box

Remark 4.6.16. Beware that the square

$$\begin{array}{c} \operatorname{ho}(n\mathbf{Cat}^{\operatorname{folk}}) \xrightarrow{\mathbb{L}\lambda_{\leq n}} \operatorname{ho}(\mathbf{Ch}_{\geq 0}^{\leq n}) \\ & \downarrow^{\overline{\iota_n}} & \downarrow^{\overline{\iota_n}} \\ \operatorname{ho}(\omega\mathbf{Cat}^{\operatorname{folk}}) \xrightarrow{\mathbb{L}\lambda} \operatorname{ho}(\mathbf{Ch}_{>0}) \end{array}$$

is *not* commutative. If it were, then for every *n*-category *C* and every k > n, we would have $H_k^{\text{pol}}(\iota_n(C)) = 0$ for every k > n, which is not even true for the case n = 1 as we shall see in the following chapter.

A useful consequence of Proposition 4.6.15 is the following corollary.

Corollary 4.6.17. Let $n \ge 0$ and C be an ω -category. If C has a k-basis for every $0 \le k \le n - 1$, then the canonical map of $ho(\mathbf{Ch}_{>0})$

$$\alpha_C^{\mathrm{pol}} : \mathbb{H}^{\mathrm{pol}}(C) \to \lambda(C)$$

induces isomorphisms

$$H_k^{\text{pol}}(C) \simeq H_k(\lambda(C))$$

for every $0 \le k \le n$.

Proof. From Lemma 4.6.13 and Proposition 4.6.15, we deduce that the morphism $\overline{\tau_{\leq n}^{i}}(\alpha_{C}^{\text{pol}})$ of ho($\mathbf{Ch}_{\geq 0}^{\leq n}$) can be identified with the canonical morphism

$$\mathbb{L}\lambda_{\leq n}(\tau^{i}_{\leq n}(C)) \to \lambda_{\leq n}(\tau^{i}_{\leq n}(C)).$$

From Lemma 4.6.5, we have that $\tau_{\leq n}^i(C)$ is cofibrant for the folk model structure on nCat, and the result follows immediately from the fact that $\lambda_{\leq n}$ is left Quillen.

4.6.18. Since every ω -category trivially admits its set of 0-cells as a 0-base, it follows from the previous proposition that for every ω -category C we have

$$\mathbb{H}_0^{\mathrm{pol}}(C) \simeq H_0(\lambda(C))$$

and

$$\mathbb{H}_1^{\mathrm{pol}}(C) \simeq H_1(\lambda(C)).$$

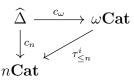
Intuitively speaking, this means that no cofibrant resolution of C is needed to compute its first two polygraphic homology groups.

We now turn to the relation between truncation and singular homology of ω -categories. Recall that for every $n \ge 0$, the nerve functor $N_n : n\mathbf{Cat} \to \widehat{\Delta}$ is defined as the following composition

$$N_n: n\mathbf{Cat} \xrightarrow{\iota_n} \omega\mathbf{Cat} \xrightarrow{N_\omega} \widehat{\Delta}_n$$

and for $n \in \mathbb{N} \cup \{\omega\}$ we write $c_n : \widehat{\Delta} \to n\mathbf{Cat}$ for the left adjoint of N_n .

Lemma 4.6.19. For every $n \in \mathbb{N}$, the following triangle of functors



is commutative (up to an isomorphism).

Proof. Straightforward consequence of the fact that $N_n = N_\omega \circ \iota_n$ and the fact that the composition of left adjoints is the left adjoint of the composition.

4.6.20. In particular, it follows from the previous lemma that the co-unit of the adjunction $c_{\omega} \dashv N_{\omega}$ induces for every ω -category C and every $n \ge 0$, a canonical morphism of nCat

$$c_n N_{\omega}(C) \simeq \tau^i_{\leq n} c_{\omega} N_{\omega}(C) \to \tau^i_{\leq n}(C),$$

which is natural in C.

Lemma 4.6.21. For every ω -category C, the canonical morphism of Cat

$$c_1 N_\omega(C) \to \tau^i_{<1}(C)$$

is an isomorphism.

Proof. Let C be an ω -category and D be a (small) category. By adjunction, we have

$$\operatorname{Hom}_{\operatorname{Cat}}(c_1 N_{\omega}(C), D) \simeq \operatorname{Hom}_{\widehat{\Delta}}(N_{\omega}(C), N_1(D)).$$
(4.4)

Now let $\Delta_{\leq 2}$ be the full subcategory of Δ spanned by [0], [1] and [2] and let $i : \Delta_{\leq 2} \to \Delta$ be the canonical inclusion. This inclusion induces by pre-composition a functor i^* : $\widehat{\Delta} \to \widehat{\Delta_{\leq 2}}$ which has a right-adjoint $i_* : \widehat{\Delta_{\leq 2}} \to \widehat{\Delta}$. Recall that the nerve of a (small) category is 2-coskeletal (see for example [Str87, Theorem 5.2]), which means that for every category D, the unit morphism $N_1(D) \to i_*i^*(N_1(D))$ is an isomorphism of simplicial sets. In particular, we have

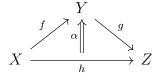
$$\operatorname{Hom}_{\widehat{\Delta}}(N_{\omega}(C), N_{1}(D)) \simeq \operatorname{Hom}_{\widehat{\Delta}}(N_{\omega}(C), i_{*}i^{*}(N_{1}(D)))$$
$$\simeq \operatorname{Hom}_{\widehat{\Delta_{\leq 2}}}(i^{*}(N_{\omega}(C)), i^{*}(N_{1}(D))).$$

Using the description of \mathcal{O}_0 , \mathcal{O}_1 and \mathcal{O}_2 from 3.1.2, we deduce that a morphism F: $i^*(N_{\omega}(C)) \to i^*(N_1(D))$ of $\widehat{\Delta_{\leq 2}}$ consists of a function $F_0: C_0 \to D_0$ and a function $F_1: C_1 \to D_1$ such that

- (a) for every $x \in C_0$, we have $F_1(1_x) = 1_{F_0(x)}$,
- (b) for every $x \in C_1$, we have

$$s(F_1(x)) = F_0(s(x))$$
 and $t(F_1(x)) = F_0(t(x)))$,

(c) for every 2-triangle



in C, we have $F_1(g) * F_1(f) = F_1(h)$.

In particular, it follows that F_1 is compatible with composition of 1-cells in an obvious sense and that for every 2-cell $\alpha : f \Rightarrow g$ of C, we have $F_1(f) = F_1(g)$. And conversely, this last condition implies condition (c) above. This means exactly that we have a natural isomorphism

$$\operatorname{Hom}_{\widehat{\Delta_{\leq 2}}}(i^*(N_{\omega}(C)), i^*(N_1(D))) \simeq \operatorname{Hom}_{\operatorname{Cat}}(\tau^i_{\leq 1}(C), D).$$

Altogether, we have

$$\operatorname{Hom}_{\operatorname{Cat}}(c_1 N_{\omega}(C), D) \simeq \operatorname{Hom}_{\operatorname{Cat}}(\tau^i_{\leq 1}(C), D),$$

which proves that

$$c_1 N_{\omega}(C) \simeq \tau^i_{<1}(C)$$

and a thorough analysis of naturality shows that this isomorphism is nothing but the canonical morphism $c_1 N_{\omega}(C) \to \tau^i_{<1}(C)$.

We can now prove the important following proposition.

Proposition 4.6.22. For every ω -category *C*, the canonical map of $ho(\mathbf{Ch}_{\geq 0})$

$$\alpha_C^{\operatorname{Sing}} : \mathbb{H}^{\operatorname{Sing}}(C) \to \lambda(C)$$

induces isomorphisms

$$H_k^{\text{Sing}}(C) \simeq H_k(\lambda(C))$$

for $k \in \{0, 1\}$.

Proof. Let C be an ω -category. Recall from 4.4.7 that the canonical morphism α_C^{Sing} : $\mathbb{H}^{\text{Sing}}(C) \to \lambda(C)$ is nothing but the image by the localization functor $\mathbf{Ch}_{\geq 0} \to \mathrm{ho}(\mathbf{Ch}_{\geq 0})$ of the morphism

$$\lambda c_{\omega} N_{\omega}(C) \to \lambda(C)$$

induced by the co-unit of the adjunction $c_{\omega} \dashv N_{\omega}$. From 4.6.15 we have that

$$\tau_{\leq 1}^{i}\lambda c_{\omega}N_{\omega}(C)\simeq\lambda_{\leq 1}\tau_{\leq 1}^{i}c_{\omega}N_{\omega}(C)=\lambda_{\leq 1}c_{1}N_{\omega}(C),$$

and from Lemma 4.6.21 we obtain

$$\tau_{\leq 1}^i \lambda c_\omega N_\omega(C) \simeq \lambda_{\leq 1} \tau_{\leq 1}^i(C) \simeq \tau_{\leq 1}^i \lambda(C).$$

This means exactly that the image by $\overline{\tau_{\leq 1}^i}$ of α_C^{Sing} is an isomorphism, which is what we wanted to prove.

Finally, we obtain the result we were aiming for.

Proposition 4.6.23. For every ω -category C, the canonical comparison map

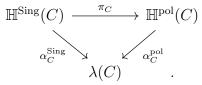
$$\pi_C : \mathbb{H}^{\mathrm{Sing}}(C) \to \mathbb{H}^{\mathrm{pol}}(C)$$

induces isomorphisms

$$H_k^{\operatorname{Sing}}(C) \simeq H_k^{\operatorname{pol}}(C)$$

for $k \in \{0, 1\}$.

Proof. Let C be an ω -category and consider the following commutative triangle of $ho(\mathbf{Ch}_{\geq 0})$



From Proposition 4.6.22, we know that α_C^{Sing} induces isomorphisms

$$H_k^{\operatorname{Sing}}(C) \simeq H_k(\lambda(C))$$

for $k \in \{0,1\}$ and from Corollary 4.6.17 and Paragraph 4.6.18 we know that α_C^{pol} induces isomorphisms $H_k^{\text{pol}}(C) \simeq H_k(\lambda(C))$ for $k \in \{0,1\}$. The result follows then from an immediate 2-out-of-3 property.

4.6.24. A natural question following the above proposition is:

For which
$$k \ge 0$$
 do we have $H_k^{\text{Sing}}(C) \simeq H_k^{\text{pol}}(C)$ for every ω -category C ?

We have already seen in 4.5.2 that when $C = B^2 \mathbb{N}$ we have

$$H_{2p}^{\text{Sing}}(B^2\mathbb{N}) \not\simeq H_{2p}^{\text{pol}}(B^2\mathbb{N})$$

for every $p \ge 2$. Furthermore, with a similar argument to the one given in 4.5.2, we have that for every $k \ge 3$, the (nerve of the) ω -category $B^k \mathbb{N}$ is a $K(\mathbb{Z}, k)$. In particular, we have

$$H_{2p+3}^{\mathrm{Sing}}(B^{2p+1}\mathbb{N}) \simeq \mathbb{Z}/2\mathbb{Z}$$

for every $p \ge 1$ (see [EML54, Theorem 23.1]). On the other hand, since $B^k \mathbb{N}$ is a free k-category, we have $H_n^{\text{pol}}(B^k \mathbb{N}) = 0$ for all $n \ge k$. All in all, we have proved that for every $k \ge 4$, there exists at least one ω -category C such that

$$H_k^{\operatorname{Sing}}(C) \not\simeq H_k^{\operatorname{pol}}(C)$$

However, it is still an open question to know whether for $k \in \{2, 3\}$ we have

$$H_k^{\operatorname{Sing}}(C) \simeq H_k^{\operatorname{pol}}(C)$$

for every ω -category C. The only missing part to adapt the proof of Proposition 4.6.23 for these values of k is the analogue of Lemma 4.6.21. But contrary to the case k = 1, it is not generally true that the canonical morphism $c_k N_{\omega}(C) \rightarrow \tau^i_{\leq k}(C)$ is an isomorphism when $k \geq 2$. However, what we really need is that the image by λ of this morphism be a quasi-isomorphism. In the case k = 2, it seems that this canonical morphism admits an oplax 2-functor as an inverse up to oplax transformation which could be an hint towards the conjecture that $H_2^{\text{Sing}}(C) \simeq H_2^{\text{pol}}(C)$ for every ω -category C.

CHAPTER 4. HOMOLOGY OF ω -CATEGORIES

CHAPTER 5

HOMOLOGY OF CONTRACTIBLE ω -CATEGORIES AND ITS CONSEQUENCES

5.1 CONTRACTIBLE ω **-CATEGORIES**

Recall that for every ω -category C, we write $p_C : C \to \mathbb{D}_0$ for the canonical morphism to the terminal object \mathbb{D}_0 of ω Cat.

Definition 5.1.1. An ω -category C is *oplax contractible* when the canonical morphism $p_C: C \to \mathbb{D}_0$ is an oplax homotopy equivalence (Definition 3.4.2).

Proposition 5.1.2. Every oplax contractible ω -category C is homologically coherent and we have

$$\mathbb{H}^{\mathrm{pol}}(C) \simeq \mathbb{H}^{\mathrm{Sing}}(C) \simeq \mathbb{Z}$$

where \mathbb{Z} is seen as an object of $ho(\mathbf{Ch}_{>0})$ concentrated in degree 0.

Proof. Consider the commutative square

$$\mathbb{H}^{\mathrm{pol}}(C) \xrightarrow{\pi_C} \mathbb{H}^{\mathrm{Sing}}(C)
\downarrow_{\mathbb{H}^{\mathrm{pol}}(p_C)} \qquad \qquad \downarrow_{\mathbb{H}^{\mathrm{Sing}}(p_C)}
\mathbb{H}^{\mathrm{pol}}(\mathbb{D}_0) \xrightarrow{\pi_{\mathbb{D}_0}} \mathbb{H}^{\mathrm{Sing}}(\mathbb{D}_0).$$

It follows respectively from Proposition 3.4.4 and Proposition 4.3.10 that the right and left vertical morphisms of the above square are isomorphisms. Then, an immediate computation left to the reader shows that \mathbb{D}_0 is homologically coherent and that $\mathbb{H}^{\text{pol}}(\mathbb{D}_0) \simeq$ $\mathbb{H}^{\text{Sing}}(\mathbb{D}_0) \simeq \mathbb{Z}$. By a 2-out-of-3 property, we deduce that $\pi_C : \mathbb{H}^{\text{Sing}}(C) \to \mathbb{H}^{\text{pol}}(C)$ is an isomorphism and $\mathbb{H}^{\text{pol}}(C) \simeq \mathbb{H}^{\text{Sing}}(C) \simeq \mathbb{Z}$. \Box **Remark 5.1.3.** Definition 5.1.1 admits an obvious "lax" variation and Proposition 5.1.2 is also true for lax contractible ω -categories.

We end this section with an important result on slice ω -categories (Paragraph 3.7.1).

Proposition 5.1.4. Let A be an ω -category and a_0 an object of A. The ω -category A/a_0 is oplax contractible.

Proof. This follows from the dual of [AM20c, Proposition 5.22].

5.2 HOMOLOGY OF GLOBES AND SPHERES

Lemma 5.2.1. For every $n \in \mathbb{N}$, the ω -category \mathbb{D}_n is oplax contractible.

Proof. Recall that we write e_n for the unique non-trivial *n*-cell of \mathbb{D}_n and that by definition \mathbb{D}_n has exactly two non-trivial *k*-cells for every *k* such that $0 \le k < n$. These two *k*-cells are parallel and are given by $s_k(e_n)$ and $t_k(e_n)$.

Let $r : \mathbb{D}_0 \to \mathbb{D}_n$ be the ω -functor that points to $t_0(e_n)$ (which means that $r = \langle t_0(e_n) \rangle$ with the notations of 1.1.7). For every k-cell x of \mathbb{D}_n , we have

$$r(p(x)) = 1_{t_0(e_n)}^{(k)}$$

where we write p for the unique ω -functor $\mathbb{D}_n \to \mathbb{D}_0$.

Now for $0 \le k < n$, we define $\alpha_{s_k(e_n)}$ and $\alpha_{t_k(e_n)}$ as

$$\alpha_{\mathbf{s}_k(e_n)} = \begin{cases} \mathbf{t}_{k+1}(e_n), \text{ if } k < n-1\\ e_n, \text{ if } k = n-1 \end{cases} \quad \text{ and } \alpha_{\mathbf{t}_k(e_n)} = \begin{cases} \mathbf{1}_{\mathbf{t}_k(e_n)}, \text{ if } k < n-1\\ e_n, \text{ if } k = n-1. \end{cases}$$

It is straightforward to check that this data defines an oplax transformation

$$\alpha: \mathrm{id}_{\mathbb{D}_n} \Rightarrow r \circ p$$

(see 3.3.3 and Example 3.3.4), which proves the result.

In particular, for every $n \in \mathbb{N}$, \mathbb{D}_n is homologically coherent. Recall from 1.1.8 that for every $n \ge 0$, we have a cocartesian square

$$\begin{array}{ccc} \mathbb{S}_{n-1} & \stackrel{i_n}{\longrightarrow} & \mathbb{D}_n \\ & \downarrow^{i_n} & & \downarrow^{j_n^+} \\ \mathbb{D}_n & \stackrel{f_n}{\longrightarrow} & \mathbb{S}_n. \end{array}$$

Lemma 5.2.2. For every $n \ge 0$, the commutative square of simplicial sets

$$\begin{array}{cccc}
N_{\omega}(\mathbb{S}_{n-1}) & \xrightarrow{N_{\omega}(i_n)} & N_{\omega}(\mathbb{D}_n) \\
& & & \downarrow_{N_{\omega}(i_n)} & & \downarrow_{N_{\omega}(j_n^+)} \\
& & & N_{\omega}(\mathbb{D}_n) & \xrightarrow{N_{\omega}(j_n^-)} & N_{\omega}(\mathbb{S}_n)
\end{array}$$

is cocartesian.

Proof. Since colimits in presheaf categories are computed pointwise, what we need to show is that for every $k \ge 0$, the following commutative square is cocartesian

Notice first that the square

$$\begin{array}{ccc} \mathbb{S}_{n-1} & \stackrel{i_n}{\longrightarrow} & \mathbb{D}_n \\ & & \downarrow^{i_n} & & \downarrow^{j_n^+} \\ \mathbb{D}_n & \stackrel{j_n^-}{\longrightarrow} & \mathbb{S}_n. \end{array}$$

is cartesian and all four morphisms are monomorphisms. Since the functor

$$\operatorname{Hom}_{\omega \operatorname{Cat}}(\mathcal{O}_k, -) : \omega \operatorname{Cat} \to \operatorname{Set}$$

preserves limits, the square (5.1) is a cartesian square of Set all of whose four morphisms are monomorphisms. Hence, in order to prove that square (5.1) is cocartesian, we only need to show that for every $k \ge 0$ and every ω -functor $\varphi : \mathcal{O}_k \to \mathbb{S}_n$, there exists an ω -functor $\varphi' : \mathcal{O}_k \to \mathbb{D}_n$ such that either $j_n^+ \circ \varphi' = \varphi$ or $j_n^- \circ \varphi' = \varphi$.

For convenience, let us write h_n^+ (resp. h_n^-) for the only generating *n*-cell of \mathbb{S}_n contained in the image of j_n^+ (resp. j_n^-). The cells h_n^+ and h_n^- are the only non-trivial *n*-cells of \mathbb{S}_n . We also write α_k for the principal cell of \mathcal{O}_k (see 3.1.2). This is the only non-trivial *k*-cell of \mathcal{O}_k .

Now, let $\varphi : \mathcal{O}_k \to \mathbb{S}_n$ be an ω -functor. There are several cases to distinguish.

Case k < n: Since every generating cell of γ of \mathcal{O}_k is of dimension not greater than k, the cell $\varphi(\gamma)$ is of dimension strictly lower than n. Since all cells of dimension strictly lower than n are both in the image of j_n^+ and in the image of j_n^- , φ obviously factors through j_n^+ (and j_n^-).

- **Case** k = n: The image of α_n is either a non-trivial *n*-cell of \mathbb{S}_n or a unit on a strictly lower dimensional cell. In the second situation, everything works like the case k < n. Now suppose for example that $\varphi(\alpha_n)$ is h_n^+ , which is in the image of j_n^+ . Since all of the other generating cells of \mathcal{O}_n are of dimension strictly lower than n, their images by φ are also of dimension strictly lower than n and hence, are all contained in the image of j_n^+ . Altogether this proves that φ factors through j_n^+ . The case where $\varphi(\alpha_n) = h_n^-$ is symmetric.
- Case k > n: Since \mathbb{S}_n is an *n*-category, the image of α_k is necessarily of the form $\varphi(\alpha_k) = 1_x^{(k)}$ with x a cell of \mathbb{S}_n of dimension non-greater than n. If x is a unit on a cell whose dimension is strictly lower than n, then everything works like in the case k < n. If not, this means that x is a non-trivial n-cell of \mathbb{S}_n . Suppose for example that $x = h_n^+$. Now let γ be a generator of \mathcal{O}_k of dimension k-1. We have $\varphi(\gamma) = 1_y^{(k-1)}$ with y which is either a unit on a cell of dimension strictly lower than n, or a non-degenerate n-cell of \mathbb{S}_n (if k-1 = n, we use the convention that $1_y^{(k-1)} = y$). In the first situation, y is in the image of j_n^+ as in the case k < n, and thus, so is $1_y^{(k-1)}$. In the second situation, this means a priori that either $y = h_n^+$ or $y = h_n^-$. But we know that γ is part of a composition that is equal to either the source or the target of α_k (see 3.1.2) and thus, y is part of a composition that is equal to either the source or the target of $x = h_n^+$. Since no composition involving h_n^- can be equal to h_n^+ (one could invoke the function introduced in 1.3.5), this implies that $y = h_n^+$ and hence, $f(\gamma)$ is in the image of j_n^+ . This goes for all generating cells of dimension k-1 of \mathcal{O}_k and we can recursively apply the same reasoning for generating cells of dimension k - 2, then k - 3 and so forth. Altogether, this proves that φ factorizes through j_n^+ . The case where $x = h_n^-$ and φ factorizes through j_n^- is symmetric.

From these two lemmas, follows the important proposition below.

Proposition 5.2.3. For every $n \ge -1$, the ω -category \mathbb{S}_n is homologically coherent.

Proof. Recall that the cofibrations of simplicial sets are exactly the monomorphisms, and in particular all simplicial sets are cofibrant. Since $i_n : \mathbb{S}_{n-1} \to \mathbb{D}_n$ is a monomorphism for every $n \ge 0$ and since N_{ω} preserves monomorphisms (as a right adjoint), it follows from Lemma 5.2.2 and Lemma 2.5.12 that the square

$$N_{\omega}(\mathbb{S}_{n-1}) \xrightarrow{N_{\omega}(i_n)} N_{\omega}(\mathbb{D}_n)$$
$$\downarrow^{N_{\omega}(i_n)} \qquad \qquad \downarrow^{N_{\omega}(j_n^-)}$$
$$N_{\omega}(\mathbb{D}_n) \xrightarrow{N_{\omega}(j_n^-)} N_{\omega}(\mathbb{S}_n)$$

is a homotopy cocartesian square of simplicial sets. Since N_{ω} induces an equivalence of op-prederivators $\mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta})$ (Theorem 3.2.5), it follows that the square of $\omega \mathbf{Cat}$

$$\begin{array}{ccc} \mathbb{S}_{n-1} & \stackrel{i_n}{\longrightarrow} & \mathbb{D}_n \\ & \downarrow^{i_n} & & \downarrow^{j_n^+} \\ \mathbb{D}_n & \stackrel{j_n^-}{\longrightarrow} & \mathbb{S}_n \end{array}$$

is Thomason homotopy cocartesian for every $n \ge 0$. Finally, since $i_n : \mathbb{S}_{n-1} \to \mathbb{D}_n$ is a folk cofibration and \mathbb{S}_{n-1} and \mathbb{D}_n are folk cofibrant for every $n \ge 0$, we deduce the desired result from Corollary 4.5.11 and an immediate induction. The base case being simply that $\mathbb{S}_{-1} = \emptyset$ is obviously homologically coherent.

5.2.4. The previous proposition implies what we claimed in Paragraph 4.5.9, which is that the morphism of op-prederivators

$$\mathcal{J}: \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{folk}}) \to \mathcal{H}o(\omega \mathbf{Cat}^{\mathrm{Th}})$$

induced by the identity functor of ω Cat is *not* homotopy cocontinuous. Indeed, recall from Paragraph 4.5.2 that the category $B^2\mathbb{N}$ is *not* homologically coherent; but on the other hand we have a cocartesian square

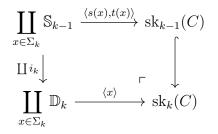
$$\begin{array}{ccc} \mathbb{S}_1 & \longrightarrow & \mathbb{D}_0 \\ & \downarrow^{i_2} & & \downarrow \\ \mathbb{D}_2 & \longrightarrow & B^2 \mathbb{N}, \end{array}$$

where the map $\mathbb{D}_2 \to B^2 \mathbb{N}$ points the unique generating 2-cell of $B^2 \mathbb{N}$ and $\mathbb{D}_0 \to B^2 \mathbb{N}$ points to the only object of $B^2 \mathbb{N}$. Since \mathbb{S}_1 , \mathbb{D}_0 and \mathbb{D}_2 are free and i_2 is a folk cofibration, the square is also folk homotopy cocartesian. If \mathcal{J} was homotopy cocontinuous, then this square would also be Thomason homotopy cocartesian. Since we know that \mathbb{S}_1 , \mathbb{D}_0 and \mathbb{D}_2 are homologically coherent, this would imply that $B^2 \mathbb{N}$ is homologically coherent.

From Proposition 5.2.3, we also deduce the proposition below which gives a criterion to detect homologically coherent ω -categories when we already know that they are free.

Proposition 5.2.5. Let C be a free ω -category and for every $k \in \mathbb{N}$ let Σ_k be its k-basis.

If for every $k \in \mathbb{N}$, the following cocartesian square (see 1.2.1)



is Thomason homotopy cocartesian, then C is homologically coherent.

Proof. Since the morphisms i_k are folk cofibrations and the ω -categories \mathbb{S}_{k-1} and \mathbb{D}_k are folk cofibrant and homologically coherent, it follows from Corollary 4.5.11 and an immediate induction that all $\mathrm{sk}_k(C)$ are homologically coherent. The result follows then from Lemma 1.1.6, Corollary 2.5.10 and Proposition 2.5.11.

5.3 The miraculous case of 1-categories

Recall that the terms *1-category* and *(small) category* are synonymous. While we have used the latter one more often so far, in this section we will mostly use the former one. As usual, the canonical functor $\iota_1 : \mathbf{Cat} \to \omega \mathbf{Cat}$ is treated as an inclusion functor and hence we always consider 1-categories as particular cases of ω -categories.

The goal of what follows is to show that every 1-category is homologically coherent. In order to do that, we will prove that every 1-category is a canonical colimit of contractible 1-categories and that this colimit is homotopic both with respect to the folk weak equivalences and with respect to the Thomason equivalences. We call the reader's attention to an important subtlety here: even though the desired result only refers to 1-categories, we have to work in the setting of ω -categories. This can be explained from the fact that if we take a cofibrant resolution of a 1-category C in the folk model structure on ω Cat

$$P \to C$$
,

then P is not necessarily a 1-category. In particular, polygraphic homology groups of a 1-category need *not* be trivial in dimension higher than 1.

5.3.1. Let A be a 1-category and a an object of A. Recall that we write A/a for the slice 1-category of A over a, that is the 1-category whose description is as follows:

- an object of A/a is a pair $(a', p : a' \to a)$ where a' is an object of A and p is an arrow of A,

an arrow of A/a is a pair (q, p : a' → a) where p is an arrow of A and q is an arrow of A of the form q : a" → a'. The target of (q, p) is given by (a', p) and the source by (a", p ∘ q).

We write π_a for the canonical forgetful functor

$$\pi_a: A/a \to A$$
$$(a', p) \mapsto a'.$$

This is a special case of the more general notion of slice ω -category introduced in 3.7.1. In particular, given an ω -category X and an ω -functor $f : X \to A$, we have defined the ω -category X/a and the ω -functor

$$f/a: X/a \to A/a$$

as the following pullback

More explicitly, the *n*-cells of X/a can be described as pairs (x, p) where x is an *n*-cell of X and p is an arrow of A of the form

$$p: f(x) \to a \text{ if } n = 0$$

and

$$p: f(\mathbf{t}_0(x)) \to a \text{ if } n > 0.$$

From now on, let us use the convention that $t_0(x) = x$ when x is a 0-cell of X.

When n > 0, the source and target of an *n*-cell (x, p) of X/a are given by

$$s((x, p)) = (s(x), p)$$
 and $t((x, p)) = (t(x), p)$.

Moreover, the ω -functor f/a is described as

$$(x,p) \mapsto (f(x),p),$$

and the canonical ω -functor $X/a \to X$ as

 $(x,p) \mapsto x.$

5.3.2. Let $f: X \to A$ be an ω -functor with A a 1-category. Every arrow $\beta: a \to a'$ of A induces an ω -functor

$$X/\beta: X/a \to X/a'$$
$$(x, p) \mapsto (x, \beta \circ p),$$

which takes part in a commutative triangle

$$\begin{array}{c} X/a \xrightarrow{X/\beta} X/a' \\ \swarrow \\ X \end{array}$$

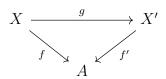
This defines a functor

$$X/-: A \to \omega \mathbf{Cat}$$
$$a \mapsto X/a$$

and a canonical ω -functor

$$\operatorname{colim}_{a\in A}(X/a)\to X$$

Let $f': X' \to A$ be another ω -functor and let



be a commutative triangle in ω Cat. Recall from 3.7.1 that for every object a of A, there is an ω -functor g/a defined as

$$g/a: X/a \to X'/a$$
$$(x,p) \mapsto (g(x),p).$$

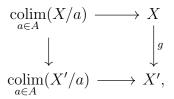
This defines a natural transformation

$$g/-: X/- \Rightarrow X'/-,$$

and thus induces an $\omega\text{-functor}$

$$\operatorname{colim}_{a \in A}(X/a) \to \operatorname{colim}_{a \in A}(X'/a).$$

Furthermore, it is immediate to check that the square



is commutative.

Lemma 5.3.3. Let $f : X \to A$ be an ω -functor such that A is a 1-category. The canonical ω -functor

$$\operatorname{colim}_{a \in A} (X/a) \to X$$

is an isomorphism.

Proof. We have to show that the cocone

$$(X/a \to X)_{a \in \mathrm{Ob}(A)}$$

is colimiting. Let

$$(\phi_a: X/a \to C)_{a \in \mathrm{Ob}(A)}$$

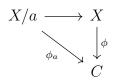
be another cocone and let x be a n-arrow of X. Notice that the pair

 $(x, 1_{f(t_0(x))})$

is a *n*-arrow of $X/f(t_0(x))$. We leave it to the reader to check that the formula

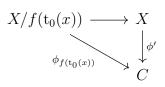
$$\phi: X \to C$$
$$x \mapsto \phi_{f(\mathbf{t}_0(x))}(x, \mathbf{1}_{f(\mathbf{t}_0(x))}).$$

defines an ω -functor and it is straightforward to check that for every object a of A the triangle



is commutative. This proves the existence part of the universal property.

Now let $\phi' : X \to C$ be another ω -functor that makes the previous triangle commute for every object a of A and let x be an n-cell of X. Since the triangle



is commutative, we necessarily have

$$\phi'(x) = \phi_{f(t_0(x))}(x, 1_{f(t_0(x))})$$

which proves that $\phi' = \phi$.

5.3.4. In particular, when we apply the previous lemma to $id_A : A \to A$, we obtain that every 1-category A is (canonically isomorphic to) the colimit

$$\operatorname{colim}_{a \in A} (A/a).$$

We now proceed to prove that this colimit is homotopic with respect to the folk weak equivalences.

Up to Lemma 5.3.7, we fix once and for all an ω -functor $f : X \to A$ with A a 1-category.

Lemma 5.3.5. If X is free, then for every object a of A, the ω -category X/a is free. More precisely, if Σ_n^X is the n-basis of X, then the n-basis of X/a is the set

$$\Sigma_n^{X/a} := \{ (x, p) \in (X/a)_n | x \in \Sigma_n^X \}.$$

Proof. It is immediate to check that for every object a of A, the canonical forgetful functor $\pi_a : A/a \to A$ is a discrete Conduché functor (see Section 1.6). Hence, from Lemma 1.6.5 we know that $X/a \to X$ is a discrete Conduché ω -functor. The result follows then from Theorem 1.6.18.

5.3.6. When X is free, every arrow $\beta : a \to a'$ of A induces a map

$$\Sigma_n^{X/a} \to \Sigma_n^{X/a'}$$
$$(x, p) \mapsto (x, \beta \circ p).$$

This defines a functor

$$\Sigma_n^{X/-} : A \to \mathbf{Set}$$
$$a \mapsto \Sigma_n^{X/a}$$

Lemma 5.3.7. If X is free, then there is an isomorphism of functors

$$\Sigma_n^{X/-} \simeq \prod_{x \in \Sigma_n^X} \operatorname{Hom}_A \left(f(\mathbf{t}_0(x)), - \right).$$

Proof. For every object a of A and every $x \in \Sigma_n^X$, we have a canonical map

$$\operatorname{Hom}_{A} \left(f(\mathbf{t}_{0}(x)), a \right) \to \Sigma_{n}^{X/a}$$
$$p \mapsto (x, p).$$

By universal property, this induces a map

$$\coprod_{x \in \Sigma_n^X} \operatorname{Hom}_A \left(f(\mathsf{t}_0(x)), a \right) \to \Sigma_n^{X/a},$$

which is natural in a. A simple verification shows that it is a bijection.

Proposition 5.3.8. Let A be a 1-category, X be a free ω -category and $f : X \to A$ be an ω -functor. The functor

$$A \to \omega \mathbf{Cat}$$
$$a \mapsto X/a$$

is a cofibrant object for the projective model structure on $\omega Cat(A)$ induced by the folk model structure on $\omega Cat(2.5.7)$.

Proof. Recall that the set

 $\{i_n: \mathbb{S}_{n-1} \to \mathbb{D}_n | n \in \mathbb{N}\}\$

is a set of generating folk cofibrations. From Lemmas 5.3.5 and 5.3.7 we deduce that for every object a of A and every $n \in \mathbb{N}$, the canonical square

is cocartesian. It is straightforward to check that this square is natural in a in an obvious sense, which means that we have a cocartesian square in $\omega Cat(A)$:

(see 2.5.5 for notations). From the second part of Proposition 2.5.6, we deduce that for every $n \ge 0$,

$$\operatorname{sk}_{n-1}(X/-) \to \operatorname{sk}_n(X/-)$$

is a cofibration for the projective model structure on $\omega Cat(A)$. Thus, the transfinite composition

$$\emptyset \to \operatorname{sk}_0(X/-) \to \operatorname{sk}_1(X/) \to \cdots \to \operatorname{sk}_n(X/-) \to \cdots$$

which is canonically isomorphic to $\emptyset \to X/-$ (see Lemma 1.1.6), is also a cofibration for the projective model structure.

Corollary 5.3.9. Let A be a 1-category and $f : X \to A$ be an ω -functor. The canonical arrow of $ho(\omega Cat^{folk})$

$$\operatorname{hocolim}_{a \in A}^{\text{folk}}(X/a) \to X,$$

induced by the co-cone $(X/a \rightarrow X)_{a \in Ob(A)}$, is an isomorphism.

Beware that in the previous corollary, we did *not* suppose that X was free.

Proof. Let P be a free ω -category and $g: P \to X$ a folk trivial fibration and consider the following commutative diagram of $ho(\omega Cat^{folk})$

$$\begin{array}{cccc} & \underset{a \in A}{\operatorname{hocolim}}(P/a) \longrightarrow \underset{a \in A}{\operatorname{colim}}(P/a) \longrightarrow P \\ & \downarrow & \downarrow & \downarrow \\ & & \downarrow & \downarrow & \downarrow \\ & \underset{a \in A}{\operatorname{folk}}(X/a) \longrightarrow \underset{a \in A}{\operatorname{colim}}(X/a) \longrightarrow X \end{array}$$

$$(5.2)$$

where the middle and most left vertical arrows are induced by the arrows

$$g/a: P/a \to X/a,$$

and the most right vertical arrow is induced by g. Since trivial fibrations are stable by pullback, g/a is a trivial fibration. This proves that the most left vertical arrow of diagram (5.2) is an isomorphism.

Now, from Proposition 5.3.8 and Corollary 2.5.10, we deduce that the arrow

$$\operatorname{hocolim}_{a \in A}^{\text{folk}}(P/a) \to \operatorname{colim}_{a \in A}(P/a)$$

is an isomorphism. Moreover, from Lemma 5.3.3, we know that the arrows

$$\operatorname{colim}_{a \in A}(P/a) \to P$$

and

$$\operatorname{colim}_{a \in A}(X/a) \to X$$

are isomorphisms.

Finally, since g is a folk weak equivalence, the most right vertical arrow of diagram (5.2) is an isomorphism and by an immediate 2-out-of-3 property this proves that all arrows of (5.2) are isomorphisms. In particular, so is the composition of the two bottom horizontal arrows, which is what we desired to show.

We now move on to the next step needed to prove that every 1-category is homologically coherent. For that purpose, let us recall a construction commonly referred to as the "Grothendieck construction".

5.3.10. Let A be a 1-category and $F : A \to Cat$ a functor. We denote by $\int F$ or $\int_{a \in A} F(a)$ the category such that:

- An object of $\int F$ is a pair (a, x) where a is an object of A and x is an object of F(a).
- An arrow $(a, x) \rightarrow (a', x')$ of $\int F$ is a pair (f, k) where $f: a \rightarrow a'$

is an arrow of A, and

$$k: F(f)(x) \to x'$$

The unit on (a, x) is the pair $(1_a, 1_x)$ and the composition of $(f, k) : (a, x) \to (a', x')$ and $(f', k') : (a', x') \to (a'', x'')$ is given by:

$$(f', k') \circ (f, k) = (f' \circ f, k' \circ F(f')(k)).$$

Every natural transformation

$$A \underbrace{\bigoplus_{G}^{F}}_{G} \mathbf{Cat}$$

induces a functor

$$\int \alpha : \int F \to \int G$$
$$(a, x) \mapsto (a, \alpha_a(x))$$

Altogether, this defines a functor

$$\int : \mathbf{Cat}(A) \to \mathbf{Cat}$$
$$F \mapsto \int F,$$

where Cat(A) is the category of functors from A to Cat.

We now recall an important theorem due to Thomason.

Theorem 5.3.11 (Thomason). The functor $\int : Cat(A) \to Cat$ sends the pointwise Thomason equivalences (2.2.5) to Thomason equivalences and the induced functor

$$\overline{\int} : \operatorname{ho}(\mathbf{Cat}^{\mathrm{Th}}(A)) \to \operatorname{ho}(\mathbf{Cat}^{\mathrm{Th}})$$

is canonically isomorphic to the homotopy colimit functor

$$\operatorname{hocolim}_{A}^{\operatorname{Th}} : \operatorname{ho}(\mathbf{Cat}^{\operatorname{Th}}(A)) \to \operatorname{ho}(\mathbf{Cat}^{\operatorname{Th}}).$$

Proof. The original source for this Theorem is [Tho79]. However, the definition of homotopy colimit used by Thomason, albeit equivalent, is not the same as the one we used in this dissertation and is slightly outdated. A more modern proof of the theorem can be found in [Mal05, Proposition 2.3.1 and Théorème 1.3.7].

Corollary 5.3.12. Let A be a 1-category. The canonical map

$$\operatorname{hocolim}_{a \in A}^{\operatorname{Th}}(A/a) \to A$$

induced by the co-cone $(A/a \to A)_{a \in Ob(A)}$, is an isomorphism of $ho(\mathbf{Cat}^{\mathrm{Th}})$.

Proof. For every object a of A, the canonical map to the terminal category

$$A/a \to \mathbb{D}_0$$

is a Thomason equivalence. This comes from the fact that A/a is oplax contractible (Proposition 5.1.4), or from [Qui73, Section 1, Corollary 2] and the fact that A/a has a terminal object.

In particular, the morphism of functors

$$A/(-) \Rightarrow k_{\mathbb{D}_0},$$

where $k_{\mathbb{D}_0}$ is the constant functor $A \to \mathbf{Cat}$ with value the terminal category \mathbb{D}_0 , is a pointwise Thomason equivalence. It follows from the first part of Theorem 5.3.11 that

$$\int_{a\in A} A/a \to \int_{a\in A} k_{\mathbb{D}_0}$$

is a Thomason equivalence and an immediate computation shows that

$$\int_{a\in A} k_{\mathbb{D}_0} \simeq A.$$

From the second part of Theorem 5.3.11, we have that

$$\operatorname{hocolim}_{a \in A}^{\operatorname{Th}}(A/a) \simeq A.$$

A thorough analysis of all the isomorphisms involved shows that this last isomorphism is indeed induced by the co-cone $(A/a \rightarrow A)_{a \in Ob(A)}$.

Remark 5.3.13. It is possible to extend the previous corollary to prove that for every functor $f : X \to A$ (X and A being 1-categories), we have

$$\operatorname{hocolim}_{a \in A} (X/a) \simeq X.$$

However, to prove that it is also the case when X is an ω -category and f an ω -functor, as in Corollary 5.3.9, one would need to extend the Grothendieck construction to functors with value in ω Cat and to prove an ω -categorical analogue of Theorem 5.3.11. Such results, while being highly plausible, go beyond the scope of this dissertation.

Putting all the pieces together, we are now able to prove the awaited theorem.

Theorem 5.3.14. Every 1-category is homologically coherent.

Proof. All the arguments of the proof have already been given and we sum them up here essentially for the sake of clarity. Let A be a 1-category. Consider the diagram

$$\begin{array}{c} A \to \omega \mathbf{Cat} \\ a \mapsto A/a \end{array}$$

and the co-cone

$$A/a \to A)_{a \in \mathrm{Ob}(A)}.$$

- The canonical map of $ho(\omega Cat^{folk})$

$$\operatorname{hocolim}_{a \in A} (A/a) \to A$$

is an isomorphism thanks to Corollary 5.3.9 applied to $id_A : A \to A$.

- The canonical map of $ho(\omega Cat^{Th})$

$$\operatorname{hocolim}_{a \in A}^{\operatorname{Th}}(A/a) \to A$$

is an isomorphism thanks to Corollary 5.3.12 and the fact that the canonical morphisms of op-prederivators $\mathcal{H}o(\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\omega\mathbf{Cat}^{\mathrm{Th}})$ is homotopy cocontinuous (see 3.2.10).

- Every A/a is homologically coherent thanks to Proposition 5.1.2 and Proposition 5.1.4.

Thus, Proposition 4.5.10 applies and this proves that A is homologically coherent. \Box

CHAPTER 5. CONTRACTIBLE ω -CATEGORIES AND CONSEQUENCES

CHAPTER 6

HOMOTOPY AND HOMOLOGY TYPE OF FREE 2-CATEGORIES

6.1 **PRELIMINARIES: THE CASE OF FREE 1-CATEGORIES**

In this section, we review some homotopical results on free (1-)categories that will be of great help in the sequel.

6.1.1. A reflexive graph G consists of the data of two sets G_0 and G_1 together with

- a "source" map $s: G_1 \to G_0$,
- a "target" map $t: G_1 \to G_0$,
- a "unit" map $1_{(-)}: G_0 \to G_1$,

such that for every $x \in G_0$,

$$(1_x) = \operatorname{t}(1_x) = x.$$

S

The vocabulary of categories is used: elements of G_0 are *objects* or 0-*cells*, elements of G_1 are *arrows* or 1-*cells*, arrows of the form 1_x with x an object are *units*, etc. A *morphism of reflexive graphs* $f : G \to G'$ consists of maps $f_0: G_0 \to G'_0$ and $f_1 :$ $G_1 \to G'_1$ that commute with sources, targets and units in an obvious sense. This defines the category **Rgrph** of reflexive graphs. Later we will make use of monomorphisms in the category **Rgrph**; they are the morphisms $f : G \to G'$ that are injective on objects and on arrows, i.e. such that $f_0: G_0 \to G'_0$ and $f_1: G_1 \to G'_1$ are injective.

There is an "underlying reflexive graph" functor

$$U: \mathbf{Cat} \to \mathbf{Rgrph},$$

which has a left adjoint

$$L: \mathbf{Rgrph} \to \mathbf{Cat}.$$

For a reflexive graph G, the objects of L(G) are exactly the objects of G and an arrow f of L(G) is a chain

$$X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_{n-1} \xrightarrow{f_n} X_n$$

of arrows of G, such that *none* of the f_k are units. The integer n is referred to as the *length* of f and is denoted by $\ell(f)$. Composition is given by concatenation of chains.

Lemma 6.1.2. A category C is free in the sense of 1.2.5 if and only if there exists a reflexive graph G such that

$$C \simeq L(G).$$

Proof. If C is free, consider the reflexive graph G such that $G_0 = C_0$ and G_1 is the subset of C_1 whose elements are either generating 1-cells of C or units. It is straightforward to check that $C \simeq L(G)$.

Conversely, if $C \simeq L(G)$ for some reflexive graph G, then the description of the arrows of L(G) given in the previous paragraph shows that C is free and that its set of generating 1-cells is (isomorphic to) the set of non unital 1-cells of G.

Remark 6.1.3. In other words, a category is free on a graph if and only if it is free on a reflexive graph. The difference between these two notions is at the level of morphisms: there are more morphisms of reflexive graphs because (generating) 1-cells may be sent to units. Hence, for a morphism of reflexive graphs $f : G \to G'$, the induced functor L(f) is not necessarily rigid in the sense of Definition 1.2.11.

6.1.4. There is another important description of the category Rgrph. Write $\Delta_{\leq 1}$ for the full subcategory of Δ spanned by [0] and [1]. The category Rgrph is nothing but $\widehat{\Delta_{\leq 1}}$, the category of pre-sheaves on $\Delta_{\leq 1}$. In particular, the canonical inclusion $i : \Delta_{\leq 1} \to \Delta$ induces by pre-composition a functor

$$i^*:\widehat{\Delta}\to\mathbf{Rgrph},$$

which, by the usual technique of Kan extensions, has a left adjoint

$$i_!: \mathbf{Rgrph} \to \Delta.$$

For a graph G, the simplicial set $i_!(G)$ has G_0 as its set of 0-simplices, G_1 as its set of 1-simplices and all k-simplices are degenerate for k > 1. For future reference, we put here the following lemma.

Lemma 6.1.5. The functor $i_!$: **Rgrph** $\rightarrow \widehat{\Delta}$ preserves monomorphisms.

Proof. What we need to show is that, given a morphism of simplicial sets

$$f: X \to Y,$$

if $f_0 : X_0 \to Y_0$ and $f_1 : X_1 \to Y_1$ are monomorphisms and if all *n*-simplices of X are degenerate for $n \ge 2$, then f is a monomorphism. A proof of this assertion is contained in [GZ67, Paragraph 3.4]. The key argument is the Eilenberg–Zilber Lemma (Proposition 3.1 of op. cit.).

6.1.6. Let us denote by $N : \widehat{\Delta} \to \mathbf{Cat}$ (instead of N_1 as in Paragraph 3.1.3) the usual nerve of categories and by $c : \mathbf{Cat} \to \widehat{\Delta}$ its left adjoint. Recall that for a (small) category C, an *n*-simplex of N(C) is a chain

$$X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_{n-1} \xrightarrow{f_n} X_n$$

of arrows of C. Such an n-simplex is degenerate if and only if at least one of the f_k is a unit. It is straightforward to check that the composite of

$$\mathbf{Cat} \stackrel{N}{\to} \widehat{\Delta} \stackrel{i^*}{\to} \mathbf{Rgrph}$$

is nothing but the forgetful functor $U : Cat \to Rgrph$. Thus, the functor $L : Rgrph \to Cat$ is (isomorphic to) the composite of

$$\mathbf{Rgrph} \stackrel{i_!}{
ightarrow} \widehat{\Delta} \stackrel{c}{
ightarrow} \mathbf{Cat}.$$

We now review a construction due to Dwyer and Kan ([DK80]). Let G be a reflexive graph. For every $k \ge 1$, we define the simplicial set $N^k(G)$ as the sub-simplicial set of N(L(G)) whose n-simplices are chains

$$X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_{n-1} \xrightarrow{f_n} X_n$$

of arrows of L(G) such that

$$\sum_{1 \le i \le n} \ell(f_i) \le k.$$

In particular, we have

$$N^1(G) = i_!(G)$$

and the transfinite composition of

$$i_!(G) = N^1(G) \hookrightarrow N^2(G) \hookrightarrow \cdots \hookrightarrow N^k(G) \hookrightarrow N^{k+1}(G) \hookrightarrow \cdots$$

is easily seen to be the map

$$\eta_{i_!(G)}: i_!(G) \to Nci_!(G),$$

where η is the unit of the adjunction $c \dashv N$.

Lemma 6.1.7 (Dwyer–Kan). For every $k \ge 1$, the canonical inclusion map

$$N^k(G) \to N^{k+1}(G)$$

is a trivial cofibration of simplicial sets.

Proof. Let $A_{k+1} = \text{Im}(\partial_0) \cup \text{Im}(\partial_{k+1})$ be the union of the first and last face of the standard (k+1)-simplex Δ_{k+1} . Notice that the canonical inclusion

$$A_{k+1} \hookrightarrow \Delta_{k+1}$$

is a trivial cofibration. Let I_{k+1} be the set of chains

$$f = X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_{k-1} \xrightarrow{f_k} X_k \xrightarrow{f_{k+1}} X_{k+1}$$

of arrows of L(G) such that for every $1 \le i \le k+1$

$$\ell(f_i) = 1,$$

i.e. each f_i is a non-unit arrow of G. For every $f \in I_{k+1}$, we define a morphism $\varphi_f : A_{k+1} \to N^k(G)$ in the following fashion:

- $\varphi_f|_{\operatorname{Im}(\partial_0)}$ is the k-simplex of $N^k(G)$

$$X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_k \xrightarrow{f_{k+1}} X_{k+1},$$

- $\varphi_f|_{\operatorname{Im}(\partial_{k+1})}$ is the k-simplex of $N^k(G)$

$$X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} X_2 \longrightarrow \cdots \longrightarrow X_{k-1} \xrightarrow{f_k} X_k.$$

All in all, we have a cocartesian square

which proves that the right vertical arrow is a trivial cofibration.

From this lemma, we deduce the following proposition.

Proposition 6.1.8. *Let G be a reflexive graph. The map*

$$\eta_{i_!(G)}: i_!(G) \to Nci_!(G),$$

where η is the unit of the adjunction $c \dashv N$, is a trivial cofibration of simplicial sets.

Proof. This follows from the fact that trivial cofibrations are stable by transfinite composition. \Box

From the previous proposition, we deduce the following very useful corollary.

Corollary 6.1.9. Let

$$\begin{array}{ccc} A & \stackrel{\beta}{\longrightarrow} & B \\ \downarrow^{\alpha} & & \downarrow^{\delta} \\ C & \stackrel{\gamma}{\longrightarrow} & D \end{array}$$

be a cocartesian square in **Rgrph**. If either α or β is a monomorphism, then the induced square of **Cat**

$$L(A) \xrightarrow{L(\beta)} L(B)$$
$$\downarrow^{L(\alpha)} \qquad \downarrow^{L(\delta)}$$
$$L(C) \xrightarrow{L(\gamma)} L(D)$$

is Thomason homotopy cocartesian.

Proof. Since the nerve N induces an equivalence of op-prederivators

$$\mathcal{H}o(\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta}),$$

it suffices to prove that the induced square of simplicial sets

$$NL(A) \xrightarrow{NL(\beta)} NL(B)$$
$$\downarrow^{NL(\alpha)} \qquad \qquad \downarrow^{NL(\delta)}$$
$$NL(C) \xrightarrow{NL(\gamma)} NL(D)$$

is homotopy cocartesian. But, since $L \simeq c \circ i_1$, it follows from Lemma 6.1.9 that this last square is weakly equivalent to the square of simplicial sets

$$\begin{split} i_!(A) & \xrightarrow{i_!(\beta)} i_!(B) \\ & \downarrow_{i_!(\alpha)} & \downarrow_{i_!(\delta)} \\ i_!(C) & \xrightarrow{i_!(\gamma)} i_!(D). \end{split}$$

This square is cocartesian because i_1 is a left adjoint. Since i_1 preserves monomorphisms (Lemma 6.1.5), the result follows from Lemma 2.5.12 and the fact that the monomorphisms are the cofibrations of the standard Quillen model structure on simplicial sets.

6.1.10. By working a little more, we obtain the more general result stated in the proposition below. Let us say that a morphism of reflexive graphs $\alpha : A \to B$ is *quasi-injective* on arrows when for all arrows f and g of A, if

$$\alpha(f) = \alpha(g),$$

then either f = g, or f and g are both units. In other words, α never sends a non-unit arrow to a unit arrow and α never identifies two non-unit arrows. It follows that if α is quasi-injective on arrows and injective on objects, then it is also injective on arrows and hence, a monomorphism of **Rgrph**.

Proposition 6.1.11. Let

$$\begin{array}{ccc} A & \stackrel{\beta}{\longrightarrow} & B \\ \downarrow^{\alpha} & & \downarrow^{\delta} \\ C & \stackrel{\gamma}{\longrightarrow} & D \end{array}$$

be a cocartesian square in **Rgrph**. Suppose that the following two conditions are satisfied

- a) Either α or β is injective on objects.
- b) Either α or β is quasi-injective on arrows.

Then, the induced square of Cat

$$L(A) \xrightarrow{L(\beta)} L(B)$$
$$\downarrow^{L(\alpha)} \qquad \downarrow^{L(\delta)}$$
$$L(C) \xrightarrow{L(\gamma)} L(D)$$

is Thomason homotopy cocartesian.

Proof. The case where α or β is both injective on objects and quasi-injective on arrows is Corollary 6.1.9. Hence, we only have to treat the case when α is injective on objects and β is quasi-injective on arrows; the remaining case being symmetric.

Let use denote by E the set of objects of B that are in the image of β . We consider this set as well as the set A_0 of objects of A as discrete reflexive graphs, i.e. reflexive graphs with no non-unit arrows. Now, let G be the reflexive graph defined by the following cocartesian square

$$\begin{array}{ccc} A_0 & \longrightarrow & E \\ \downarrow & & & \downarrow \\ A & \longrightarrow & G, \end{array}$$

where the morphism

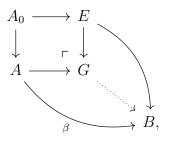
$$A_0 \to A$$

is the canonical inclusion, and the morphism

$$A_0 \to E$$

is induced by the restriction of β on objects. In other words, G is obtained from A by collapsing the objects that are identified through β . It admits the following explicit description: G_0 is (isomorphic to) E and the set of non-unit arrows of G is (isomorphic to) the set of non-unit arrows of A; the source (resp. target) of a non-unit arrow f of G is the source (resp. target) of $\beta(f)$. This completely describes G.

Now, we have the following solid arrow commutative diagram



where the arrow $E \to B$ is the canonical inclusion. Hence, by universal property, the dotted arrow exists and makes the whole diagram commute. A thorough verification easily shows that, because β is quasi-injective on arrows, the morphism $G \to B$ is a monomorphism of **Rgrph**.

By forming successive cocartesian squares and combining with the square obtained earlier, we obtain a diagram of three cocartesian squares:

$$\begin{array}{cccc} A_0 & \longrightarrow & E \\ & & & & \\ \downarrow & & & & \\ A & \longrightarrow & G & \longrightarrow & B \\ & & & & & \\ \downarrow^{\alpha} & & & & \\ C & \longrightarrow & H & \longrightarrow & D. \end{array}$$

What we want to prove is that the image by the functor L of the pasting of squares (2) and (3) is homotopy cocartesian. Since the morphism $G \rightarrow B$ is a monomorphism, we deduce from Corollary 6.1.9 that the image by the functor L of square (3) is homotopy cocartesian. Hence, in virtue of Lemma 2.4.9, all we have to show is that the image by L of square (2) is homotopy cocartesian. On the other hand, the morphisms

$$A_0 \to A$$

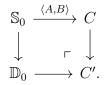
and

$$A_0 \to C$$

are monomorphisms and thus, using Corollary 6.1.9, we deduce that the image by L of square (1) and of the pasting of squares (1) and (2) are homotopy cocartesian. By Lemma 2.4.9 again, this proves that the image by L of square (2) is homotopy cocartesian.

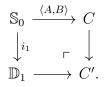
We now apply Corollary 6.1.9 and Proposition 6.1.11 to a few examples.

Example 6.1.12 (Identifying two objects). Let C be a free category, A and B be two objects of C with $A \neq B$ and let C' be the category obtained from C by identifying A and B, i.e. defined by the following cocartesian square



Then, this square is Thomason homotopy cocartesian. Indeed, it is obviously the image by the functor L of a cocartesian square of **Rgrph** and the top morphism is a monomorphism. Hence, we can apply Corollary 6.1.9.

Example 6.1.13 (Adding a generator). Let C be a free category, A and B two objects of C (possibly equal) and let C' be the category obtained from C by adding a generator $A \rightarrow B$, i.e. defined by the following cocartesian square:



Then, this square is Thomason homotopy cocartesian. Indeed, it obviously is the image of a square of **Rgrph** by the functor L and the morphism $i_1 : \mathbb{S}_0 \to \mathbb{D}_1$ comes from a monomorphism of **Rgrph**. Hence, we can apply Corollary 6.1.9.

Remark 6.1.14. Since $i_1 : \mathbb{S}_0 \to \mathbb{D}_1$ is a folk cofibration and since every free category is obtained by recursively adding generators starting from a set of objects (seen as a 0-category), the previous example yields another proof that *free* (1-)categories are homologically coherent (which we already knew since we have seen that *all* (1-)categories are homologically coherent).

Example 6.1.15 (Identifying two generators). Let C be a free category and let $f, g : A \to B$ be parallel generating arrows of C such that $f \neq g$. Now consider the category C' obtained from C by "identifying" f and g, i.e. defined by the following cocartesian square

$$\begin{array}{c} \mathbb{S}_1 \xrightarrow{\langle f,g \rangle} C \\ \downarrow & & \downarrow \\ \mathbb{D}_1 \longrightarrow C', \end{array}$$

where the morphism $\mathbb{S}_1 \to \mathbb{D}_1$ is the one that sends the two generating arrows of \mathbb{S}_1 to the unique generating arrow of \mathbb{D}_1 . Then this square is Thomason homotopy cocartesian. Indeed, it is the image by the functor L of a cocartesian square in **Rgrph**, the morphism $\mathbb{S}_1 \to \mathbb{D}_1$ is injective on objects and the morphism $\mathbb{S}_1 \to C$ is quasi-injective on arrows. Hence, we can apply Proposition 6.1.11. Note that since we did *not* suppose that $A \neq B$, the top morphism of the previous square is not necessarily a monomorphism and we cannot always apply Corollary 6.1.9.

Example 6.1.16 (Killing a generator). Let C be a free category and let $f : A \to B$ be one of its generating arrows such that $A \neq B$. Now consider the category C' obtained from C by "killing" f, i.e. defined by the following cocartesian square:

Then, this square is Thomason homotopy cocartesian. Indeed, it obviously is the image of a cocartesian square in **Rgrph** by the functor L and since the source and target of f are different, the top map comes from a monomorphism of **Rgrph**. Hence, we can apply Corollary 6.1.9.

Remark 6.1.17. Note that in the previous example, we see that it was useful to consider the category of reflexive graphs and not only the category of graphs because the map $\mathbb{D}_1 \to \mathbb{D}_0$ does not come from a morphism in the category of graphs.

Note also that the hypothesis that $A \neq B$ was fundamental in the previous example as for A = B the square is *not* Thomason homotopy cocartesian.

6.2 **PRELIMINARIES: BISIMPLICIAL SETS**

6.2.1. A *bisimplicial set* is a presheaf over the category $\Delta \times \Delta$,

$$X: \Delta^{\mathrm{op}} \times \Delta^{\mathrm{op}} \to \mathbf{Set}.$$

In a similar fashion as for simplicial sets (3.1.1), for $n, m \ge 0$, we use the notations

$$X_{n,m} := X([n], [m])$$

$$\partial_i^v := X(\delta^i, \mathrm{id}) : X_{n+1,m} \to X_{n,m}$$

$$\partial_j^h := X(\mathrm{id}, \delta^j) : X_{n,m+1} \to X_{n,m}$$

$$s_i^v := X(\sigma^i, \mathrm{id}) : X_{n,m} \to X_{n+1,m}$$

$$s_i^h := X(\mathrm{id}, \sigma^j) : X_{n,m} \to X_{n,m+1}$$

The maps ∂_i^h and s_i^h will be referred to as the *horizontal* face and degeneracy operators; and ∂_i^v and s_i^v as the *vertical* face and degeneracy operators.

Note that for every $n \ge 0$, we have simplicial sets

$$X_{\bullet,n} : \Delta^{\mathrm{op}} \to \mathbf{Set}$$
$$[k] \mapsto X_{k,n}$$

and

$$X_{n,\bullet} : \Delta^{\mathrm{op}} \to \mathbf{Set}$$

 $[k] \mapsto X_{n,k}.$

The category of bisimplicial sets is denoted by $\widehat{\Delta \times \Delta}$.

6.2.2. The functor

$$\delta: \Delta \to \Delta \times \Delta$$
$$[n] \mapsto ([n], [n])$$

induces by pre-composition a functor

$$\delta^*:\widehat{\Delta\times\Delta}\to\widehat{\Delta}.$$

By the usual calculus of Kan extensions, δ^* admits a left adjoint δ_1 and a right adjoint δ_*

$$\delta_! \dashv \delta^* \dashv \delta_*.$$

We say that a morphism $f : X \to Y$ of bisimplicial sets is a *diagonal weak equivalence* (resp. *diagonal fibration*) when $\delta^*(f)$ is a weak equivalence (resp. fibration) of simplicial sets. By definition, δ^* induces a morphism of op-prederivators

$$\overline{\delta^*}: \mathcal{H}o(\widehat{\Delta \times \Delta}^{\operatorname{diag}}) \to \mathcal{H}o(\widehat{\Delta}),$$

where $\mathcal{H}_0(\widehat{\Delta \times \Delta}^{\text{diag}})$ is the homotopy op-prederivator of $\widehat{\Delta \times \Delta}$ equipped with diagonal weak equivalences. Recall from [Moe89, Proposition 1.2] that the category of bisimplicial sets can be equipped with a model structure whose weak equivalences are the diagonal weak equivalences and whose fibrations are the diagonal fibrations. We shall refer to this model structure as the *diagonal model structure*.

Proposition 6.2.3. Consider that $\overline{\Delta} \times \overline{\Delta}$ is equipped with the diagonal model structure. Then, the adjunction

$$\delta_!:\widehat{\Delta} \xleftarrow{} \widehat{\Delta \times \Delta}: \delta^*,$$

is a Quillen equivalence.

Proof. By definition δ^* preserves weak equivalences and fibrations and thus, the adjunction is a Quillen adjunction. The fact that δ^* induces an equivalence at the level of homotopy categories is [Moe89, Proposition 1.2].

6.2.4. In particular, the morphism of op-prederivators

$$\overline{\delta^*}: \mathcal{H}o(\widehat{\Delta \times \Delta}^{\operatorname{diag}}) \to \mathcal{H}o(\widehat{\Delta})$$

is actually an equivalence of op-prederivators.

Diagonal weak equivalences are not the only interesting weak equivalences for bisimplicial sets.

6.2.5. A morphism $f: X \to Y$ of bisimplicial sets is a *vertical (resp. horizontal) weak* equivalence when for every $n \ge 0$, the induced morphism of simplicial sets

$$f_{\bullet,n}: X_{\bullet,n} \to Y_{\bullet,n}$$

(resp.

$$f_{n,\bullet}: X_{n,\bullet} \to Y_{n,\bullet})$$

is a weak equivalence of simplicial sets. Recall now a very useful lemma.

Lemma 6.2.6. Let $f : X \to Y$ be a morphism of bisimplicial sets. If f is a vertical or horizontal weak equivalence then it is a diagonal weak equivalence.

Proof. See for example [BK72, Chapter XII,4.3] or [Cis04, Proposition 2.1.7]. \Box

6.2.7. In particular, the identity functor of the category of bisimplicial sets induces the morphisms of op-prederivators:

$$\mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{vert}}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{diag}})$$

and

$$\mathcal{H}o(\widehat{\Delta \times \Delta}^{hor}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{diag}).$$

Proposition 6.2.8. The morphisms of op-prederivators

$$\mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{vert}}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{diag}})$$

and

$$\mathcal{H}o(\widehat{\Delta \times \Delta}^{hor}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{diag})$$

are homotopy cocontinuous.

Proof. Recall that the category of bisimplicial sets can be equipped with a model structure where the weak equivalences are the vertical (resp. horizontal) weak equivalences and the cofibrations are the monomorphisms (see for example [GJ09, Chapter IV] or [Cis04]). We respectively refer to these model structures as the *vertical model structure* and *horizontal model structure*. Since the functor $\delta^* : \widehat{\Delta \times \Delta} \to \widehat{\Delta}$ preserves monomorphisms, it follows from Lemma 6.2.6 that the adjunction

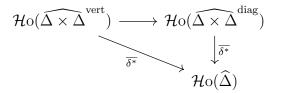
$$\delta^*:\widehat{\Delta\times\Delta} \longleftrightarrow \widehat{\Delta}:\delta_*$$

is a Quillen adjunction when $\widehat{\Delta} \times \widetilde{\Delta}$ is equipped with either the vertical model structure or the horizontal model structure. In particular, the induced morphisms of opprederivators $\overline{\delta^*} : \mathcal{H}o(\widehat{\Delta \times \Delta}^{vert}) \to \mathcal{H}o(\widehat{\Delta})$

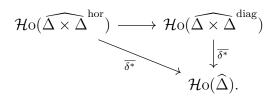
and

$$\overline{\delta^*}: \mathcal{H}_{O}(\widehat{\Delta \times \Delta}^{\mathrm{hor}}) \to \mathcal{H}_{O}(\widehat{\Delta})$$

are homotopy cocontinuous. Now, the obvious identity $\delta^* = \delta^* \circ id_{\widehat{\Delta \times \Delta}}$ implies that we have commutative triangles



and



The result follows then from the fact that $\overline{\delta^*} : \mathcal{H}o(\widehat{\Delta \times \Delta}^{\operatorname{diag}}) \to \mathcal{H}o(\widehat{\Delta})$ is an equivalence of op-prederivators.

In practice, we will use the following corollary.

Corollary 6.2.9. Let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ \downarrow^{f} & & \downarrow^{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array}$$

be a commutative square in the category of bisimplicial sets satisfying at least one of the two following conditions:

(a) For every $n \ge 0$, the square of simplicial sets

$$\begin{array}{ccc} A_{\bullet,n} & \stackrel{u_{\bullet,n}}{\longrightarrow} & B_{\bullet,n} \\ & & \downarrow_{f_{\bullet,n}} & & \downarrow_{g_{\bullet,n}} \\ C_{\bullet,n} & \stackrel{v_{\bullet,n}}{\longrightarrow} & D_{\bullet,n} \end{array}$$

is homotopy cocartesian.

(b) For every $n \ge 0$, the square of simplicial sets

$$\begin{array}{ccc} A_{n,\bullet} & \stackrel{u_{n,\bullet}}{\longrightarrow} & B_{n,\bullet} \\ & & \downarrow_{f_{n,\bullet}} & & \downarrow_{g_{n,\bullet}} \\ C_{n,\bullet} & \stackrel{v_{n,\bullet}}{\longrightarrow} & D_{n,\bullet} \end{array}$$

is homotopy cocartesian.

Then, the square

$$\delta^{*}(A) \xrightarrow{\delta^{*}(u)} \delta^{*}(B)$$

$$\downarrow^{\delta^{*}(f)} \qquad \qquad \downarrow^{\delta^{*}(g)}$$

$$\delta^{*}(C) \xrightarrow{\delta^{*}(v)} \delta^{*}(D)$$

is a homotopy cocartesian square of simplicial sets.

Proof. From [Gro16, Corollary 10.3.10(i)] we know that the square of bisimplicial sets



is homotopy cocartesian with respect to the vertical weak equivalences if and only if for every $n \ge 0$, the square

$$\begin{array}{ccc} A_{\bullet,n} & \stackrel{u_{\bullet,n}}{\longrightarrow} & B_{\bullet,n} \\ & & \downarrow_{f_{\bullet,n}} & & \downarrow_{g_{\bullet,n}} \\ C_{\bullet,n} & \stackrel{v_{\bullet,n}}{\longrightarrow} & D_{\bullet,n} \end{array}$$

is a homotopy cocartesian square of simplicial sets and similarly for horizontal weak equivalences. The result follows then from Proposition 6.2.8.

6.3 **BISIMPLICIAL NERVE FOR 2-CATEGORIES**

We shall now describe a "nerve" for 2-categories with values in bisimplicial sets and recall a few results that shows that this nerve is, in some sense, equivalent to the nerve defined in 3.1.3.

Notation 6.3.1. - Once again, we write $N : \mathbf{Cat} \to \widehat{\Delta}$ instead of N_1 for the usual nerve of categories. Moreover, using the usual notation for the set of k-simplices of a simplicial set, if C is a (small) category, then

 $N(C)_k$

is the set of k-simplices of the nerve of C.

- Similarly, we write $N : 2\mathbf{Cat} \to \widehat{\Delta}$ instead of N_2 for the nerve of 2-categories. This makes sense since the nerve for categories is the restriction of the nerve for 2-categories.
- For 2-categories, we refer to the *-composition of 2-cells as the *horizontal composition* and the *-composition of 2-cells as the *vertical composition*.
- For a 2-category C and x and y objects of C, we denote by

C(x,y)

the category whose objects are the 1-cells of C with x as source and y as target, and whose arrows are the 2-cells of C with x as 0-source and y as 0-target. Composition is induced by vertical composition in C.

6.3.2. Every 2-category C defines a simplicial object in Cat,

$$S(C): \Delta^{\mathrm{op}} \to \mathbf{Cat},$$

where, for each $n \ge 0$,

$$S_n(C) := \coprod_{(x_0,\cdots,x_n)\in Ob(C)^{\times (n+1)}} C(x_0,x_1) \times \cdots \times C(x_{n-1},x_n).$$

Note that for n = 0, the above formula reads $S_0(C) = C_0$. For n > 0, the face operators $\partial_i : S_n(C) \to S_{n-1}(C)$ are induced by horizontal composition for 0 < i < n and by projection for i = 0 or n. The degeneracy operators $s_i : S_n(C) \to S_{n+1}(C)$ are induced by the units for the horizontal composition.

Post-composing S(C) with the nerve functor $N : \mathbf{Cat} \to \widehat{\Delta}$, we obtain a functor

$$NS(C): \Delta^{\mathrm{op}} \to \widehat{\Delta}.$$

Remark 6.3.3. When C is a 1-category, the simplicial object S(C) is nothing but the usual nerve of C where, for each $n \ge 0$, $S_n(C)$ is seen as a discrete category.

Definition 6.3.4. The *bisimplicial nerve* of a 2-category C is the bisimplicial set $N_{\Delta \times \Delta}(C)$ defined as

$$N_{\Delta \times \Delta}(C)_{n,m} := N(S_n(C))_m,$$

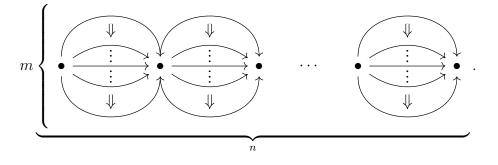
for all $n, m \ge 0$.

6.3.5. In other words, the bisimplicial nerve of C is obtained by "un-currying" the functor $NS(C) : \Delta^{\text{op}} \to \widehat{\Delta}$.

Since the nerve N commutes with products and sums, we obtain the formula

$$N_{\Delta \times \Delta}(C)_{n,m} = \prod_{(x_0, \cdots, x_n) \in Ob(C)^{\times (n+1)}} N(C(x_0, x_1))_m \times \cdots \times N(C(x_{n-1}, x_n))_m.$$
(6.1)

More intuitively, an element of $N_{\Delta \times \Delta}(C)_{n,m}$ consists of a "pasting scheme" in C that looks like



In the definition of the bisimplicial nerve of a 2-category we gave, one direction of the bisimplicial set is privileged over the other. We now give an equivalent definition of the bisimplicial nerve which uses the other direction.

6.3.6. Let C be a 2-category. For every $k \ge 1$, we define a 1-category $V_k(C)$ in the following fashion:

- The objects of $V_k(C)$ are the objects of C.
- A morphism α is a sequence

$$\alpha = (\alpha_1, \alpha_2, \cdots, \alpha_k)$$

of vertically composable 2-cells of C, i.e. such that for every $1 \le i \le k - 1$, we have

$$\mathbf{s}(\alpha_i) = \mathbf{t}(\alpha_{i+1}).$$

The source and target of α are given by

$$s(\alpha) := s_0(\alpha_1)$$
 and $t(\alpha) := t_0(\alpha_1)$.

(Note that we could have used any of the α_i instead of α_1 since they all have the same 0-source and 0-target.)

- Composition is given by

$$(\alpha_1, \alpha_2, \cdots, \alpha_k) \circ (\beta_1, \beta_2, \cdots, \beta_k) := (\alpha_1 * \beta_1, \alpha_2 * \beta_2, \cdots, \alpha_k * \beta_k)$$

and the unit on an object x is the sequence

$$(1_x^2,\cdots,1_x^2).$$

For k = 0, we define $V_0(C)$ to be the category obtained from C by simply forgetting the 2-cells (which is nothing but $\tau_{\leq 1}^s(C)$ with the notations of 1.1.4). The correspondence $n \mapsto V_n(C)$ defines a simplicial object in Cat

$$V(C): \Delta^{\mathrm{op}} \to \mathbf{Cat},$$

where the face operators are induced by the vertical composition and the degeneracy operators are induced by the units for the vertical composition.

Lemma 6.3.7. Let C be a 2-category. For every $n \ge 0$, we have

$$N(V_m(C))_n = (N_{\Delta \times \Delta}(C))_{n,m}.$$

Proof. This is simply a reformulation of the formula given in Paragraph 6.3.5. \Box

6.3.8. The bisimplicial nerve canonically defines a functor

$$N_{\Delta \times \Delta} : 2\mathbf{Cat} \to \bar{\Delta} \times \bar{\Delta}$$

which enables us to compare the homotopy theory of 2Cat with the homotopy theory of bisimplicial sets.

Lemma 6.3.9. A 2-functor $F : C \to D$ is a Thomason equivalence if and only if $N_{\Delta \times \Delta}(F)$ is a diagonal weak equivalence of bisimplicial sets.

Proof. It follows from what is shown in [BC03, Section 2] that there is a zigzag of weak equivalence of simplicial sets

$$\delta^*(N_{\Delta \times \Delta}(C)) \leftarrow \cdots \rightarrow N(C)$$

which is natural in C. This implies what we wanted to show. See also [AM20a, Théorème 3.13].

From this lemma, we deduce two useful criteria to detect Thomason equivalences of 2-categories.

Corollary 6.3.10. Let $F : C \to D$ be a 2-functor. If

a) $F_0: C_0 \rightarrow D_0$ is a bijection,

and

b) for all objects x, y of C, the functor

$$C(x,y) \to D(F(x),F(y))$$

induced by F is a Thomason equivalence of 1-categories,

then F is a Thomason equivalence of 2-categories.

Proof. By definition, for every 2-category C and every $m \ge 0$, we have

$$(N_{\Delta \times \Delta}(C))_{\bullet,m} = NS(C).$$

The result then follows from Lemma 6.2.6 and the fact that the weak equivalences of simplicial sets are stable by coproducts and finite products. \Box

Corollary 6.3.11. Let $F : C \to D$ be a 2-functor. If for every $k \ge 0$,

$$V_k(F): V_k(C) \to V_k(D)$$

is a Thomason equivalence of 1-categories, then F is a Thomason equivalence of 2-categories.

Proof. From Lemma 6.3.7, we now that for every $m \ge 0$,

$$N_{\Delta \times \Delta}(C)_{\bullet,m} = N(V_m(C)).$$

The result then follows from Lemma 6.2.6.

6.3.12. It also follows from Lemma 6.3.9 that the bisimplicial nerve induces a morphism of op-prederivators

$$\overline{N_{\Delta \times \Delta}} : \mathcal{H}o(2\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{diag}}).$$

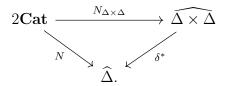
Proposition 6.3.13. The morphism of op-prederivators

$$\overline{N_{\Delta \times \Delta}} : \mathcal{H}o(2\mathbf{Cat}^{\mathrm{Th}}) \to \mathcal{H}o(\widehat{\Delta \times \Delta}^{\mathrm{diag}})$$

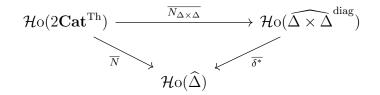
1.

is an equivalence of op-prederivators.

Proof. Consider the following triangle of functors



Even though it is *not* commutative (even up to an isomorphism), it follows from the results contained in [BC03, Section2] that the induced triangle



is commutative up to a canonical isomorphism. The result follows then from the fact that $\overline{\delta^*}$ and \overline{N} are equivalences of op-prederivators (Proposition 6.2.3 and Theorem 3.2.5 respectively).

From Proposition 6.3.13, we deduce the proposition below which contains two useful criteria to detect when a commutative square of 2Cat is Thomason homotopy cocartesian.

Proposition 6.3.14. Let

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ \downarrow_{f} & \qquad \downarrow_{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array} \tag{(*)}$$

be a commutative square in 2Cat satisfying at least one of the two following conditions:

(a) For every $n \ge 0$, the commutative square of Cat

$$V_n(A) \xrightarrow{V_n(u)} V_n(B)$$

$$V_n(f) \downarrow \qquad \qquad \downarrow V_n(g)$$

$$V_n(C) \xrightarrow{V_n(v)} V_n(D)$$

is Thomason homotopy cocartesian,

(b) For every $n \ge 0$, the commutative square of Cat

is Thomason homotopy cocartesian.

Then, square (*) is Thomason homotopy cocartesian.

Proof. This is an immediate consequence of Proposition 6.3.13 and Corollary 6.2.9. \Box

6.4 ZOOLOGY OF 2-CATEGORIES: BASIC EXAMPLES

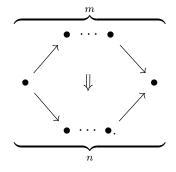
6.4.1. Before embarking on computations of homology and homotopy types of 2-categories, let us recall the following particular case of Corollary 4.5.11. Suppose given a cocartesian square

$$\begin{array}{ccc} A & \stackrel{u}{\longrightarrow} & B \\ & & & \downarrow^{f} & & \downarrow^{g} \\ C & \stackrel{v}{\longrightarrow} & D \end{array}$$

of 2-categories. If A, B and C are free and homologically coherent, if at least one of u or f is a folk cofibration and if the square is Thomason homotopy cocartesian, then D is homologically coherent.

6.4.2. Let $n, m \ge 0$. We denote by $A_{(m,n)}$ the free 2-category with only one generating

2-cell whose source is a chain of length m and whose target is a chain of length n:



More formally, $A_{(m,n)}$ is described in the following way:

- generating 0-cells: $A_0, \dots, A_m, B_1, \dots, B_{n-1}$

(and for convenience, we also set $B_0 := A_0$ and $B_n := A_m$)

- generating 1-cells: $\begin{cases} f_{i+1} : A_i \to A_{i+1} & \text{ for } 0 \le i \le m-1 \\ g_{j+1} : B_j \to B_{j+1} & \text{ for } 0 \le j \le n-1 \end{cases}$
- generating 2-cell: $\alpha : f_m \circ \cdots \circ f_1 \Rightarrow g_n \circ \cdots \circ g_1$.

Notice that $A_{(1,1)}$ is nothing but \mathbb{D}_2 . We are going to prove that if $n \neq 0$ or $m \neq 0$, then $A_{(m,n)}$ is homologically coherent and has the homotopy type of a point. When $m \neq 0$ and $n \neq 0$, this result is not surprising, but when n = 0 or m = 0 (but not both), it is a priori less clear what the homotopy type of $A_{(m,n)}$ is and whether it is homologically coherent or not. For example, $A_{(1,0)}$ can be pictured as



and has many non trivial 2-cells, such as $f \underset{0}{*} \alpha \underset{0}{*} f$.

Note that when m = 0 and n = 0, then the 2-category $A_{(0,0)}$ is nothing but the 2-category $B^2\mathbb{N}$ and we have already seen that it is *not* homologically coherent (see 4.5.2).

6.4.3. For $n \ge 0$, we write Δ_n for the linear order $\{0 \le \cdots \le n\}$ seen as a small category. Let $i : \Delta_1 \to \Delta_n$ be the unique functor such that

$$i(0) = 0$$
 and $i(1) = n$

Lemma 6.4.4. For $n \neq 0$, the functor $i : \Delta_1 \rightarrow \Delta_n$ is a strong oplax deformation retract (3.4.5).

Proof. Let $r: \Delta_n \to \Delta_1$ the unique functor such that

$$r(0) = 0$$
 and $r(k) = 1$ for $k > 0$.

By definition we have $r \circ i = 1_{\Delta_1}$. Now, the natural order on Δ_n induces a natural transformation

$$\alpha: \mathrm{id}_{\Delta_n} \Rightarrow i \circ r,$$

and it is straightforward to check that $\alpha * i = id_i$.

6.4.5. For every $n \ge 0$, consider the following cocartesian square

$$\begin{array}{ccc} \Delta_1 & & \stackrel{i}{\longrightarrow} & \Delta_n \\ \downarrow^{\tau} & & \downarrow \\ A_{(1,1)} & & \stackrel{i}{\longrightarrow} & A_{(1,n)}, \end{array}$$

where $\tau : \Delta_1 \to A_{(1,1)}$ is the 2-functor that sends the unique non-trivial 1-cell of Δ_1 to the target of the generating 2-cell of $A_{(1,1)}$. It is not hard to check that τ is strong oplax deformation retract and thus, a co-universal Thomason equivalence (Lemma 3.4.6). Hence, the morphism $\Delta_n \to A_{(1,n)}$ is also a (co-universal) Thomason equivalence and the square is Thomason homotopy cocartesian (Lemma 2.4.8). Now, the morphism $\tau : \Delta_1 \to A_{(1,1)}$ is also a folk cofibration and since Δ_1, Δ_n and $A_{(1,1)}$ are homologically coherent, it follows from what we said in 6.4.1 that $A_{(1,n)}$ is homologically coherent. Finally, since Δ_1, Δ_n and $A_{(1,1)}$ have the homotopy type of a point, the fact that the previous square is Thomason homotopy cocartesian implies that $A_{(1,n)}$ has the homotopy type of a point.

Similarly, for every $m \ge 0$, by considering the cocartesian square

$$\begin{array}{ccc} \Delta_1 & & \stackrel{i}{\longrightarrow} & \Delta_m \\ \downarrow^{\sigma} & & \downarrow \\ A_{(1,1)} & & \longrightarrow & A_{(m,1)}, \end{array}$$

where $\sigma : \Delta_1 \to A_{(1,1)}$ is the 2-functor that sends the unique non trivial 1-cell of Δ_1 to the source of the generating 2-cell of $A_{(1,1)}$, we can prove that $A_{(m,1)}$ is homologically coherent and has the homotopy type of a point.

Now, let $m \ge 0$ and n > 0 and consider the cocartesian square

$$\begin{array}{ccc} \Delta_1 & & \stackrel{i}{\longrightarrow} & \Delta_n \\ \downarrow^{\tau} & & \downarrow \\ A_{(m,1)} & \longrightarrow & A_{(m,n)}, \end{array}$$

where τ is the 2-functor that sends the unique non-trivial 1-cell of Δ_1 to the target of the generating 2-cell of $A_{(m,1)}$. This 2-functor is once again a folk cofibration, but it is *not* in general a co-universal Thomason equivalence (it would be if we had made the hypothesis that $m \neq 0$, but we did not). However, since we made the hypothesis that $n \neq 0$, it follows from Lemma 6.4.4 that $i : \Delta_1 \to \Delta_n$ is a co-universal Thomason equivalence. Hence, the previous square is Thomason homotopy cocartesian and $A_{(m,n)}$ has the homotopy type of a point. Since $A_{(m,1)}$, Δ_1 and Δ_n are homologically coherent, this shows that for $m \geq 0$ and n > 0, $A_{(m,n)}$ is homologically coherent.

Similarly, if m > 0 and $n \ge 0$, then $A_{(m,n)}$ has the homotopy type of a point and is homologically coherent.

Combined with the result of Paragraph 4.5.2, we have proved the following proposition.

Proposition 6.4.6. Let $m, n \ge 0$ and consider the 2-category $A_{(m,n)}$. If $m \ne 0$ or $n \ne 0$, then $A_{(m,n)}$ is homologically coherent and has the homotopy type of a point. If n = m = 0, then $A_{(0,0)}$ is not homologically coherent and has the homotopy type of a $K(\mathbb{Z}, 2)$.

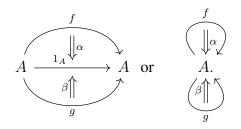
6.5 ZOOLOGY OF 2-CATEGORIES: MORE EXAMPLES

As a warm-up, let us begin with an example which is a direct consequence of the results at the end of the previous section.

6.5.1. Let *P* the free 2-category defined as follows:

- generating 0-cell: A,
- generating 1-cells: $f, g: A \to A$,
- generating 2-cells: $\alpha : f \Rightarrow 1_A, \beta : g \Rightarrow 1_A$.

In pictures, this gives



Notice that this category has many non-trivial 2-cells and it is not *a priori* clear what its homotopy type is and whether or not it is homologically coherent. Observe that P is obtained as the following amalgamated sum

Since \mathbb{D}_0 , $A_{(1,0)}$ are free and homologically coherent and since $\mathbb{D}_0 \to A_{(1,0)}$ is a folk cofibration, all we have to show to prove that P is homologically coherent is that the above square is Thomason homotopy cocartesian. Notice that the 2-category $A_{(1,0)}$ is obtained as the following amalgamated sum

$$\begin{array}{ccc} \Delta_1 & \stackrel{\tau}{\longrightarrow} & A_{(1,1)} \\ \downarrow & & \downarrow \\ \Delta_0 & \stackrel{-}{\longrightarrow} & A_{(1,0)}, \end{array}$$

where $\tau : \Delta_1 \to A_{(1,1)}$ has already been defined in 6.4.5. We have seen that τ is a co-universal Thomason equivalence and thus, so is $\mathbb{D}_0 \to A_{(0,1)}$ (as \mathbb{D}_0 and Δ_0 are two different names for the same category). Hence, square (6.2) is Thomason homotopy cocartesian and this proves that P is homologically coherent and has the homotopy type of a point.

All the variations by reversing the direction of α or β work exactly the same way.

Let us now get into more sophisticated examples.

6.5.2 (Variations of spheres). Let *P* the free 2-category defined as follows:

- generating 0-cells: A, B,
- generating 1-cells: $f, g : A \rightarrow B$,
- generating 2-cells: $\alpha : f \Rightarrow g, \beta : g \Rightarrow f$.

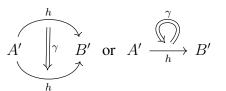
In pictures, this gives



Let P' be the free 2-category defined as follows:

- generating 0-cells: A', B',
- generating 1-cell: $h: A' \to B'$,
- generating 2-cell: $\gamma : h \Rightarrow h$,

which can be pictured as



and let $F: P \to P'$ be the unique 2-functor such that

We wish to prove that this 2-functor is a Thomason equivalence. Since it is an isomorphism on objects, it suffices to prove that the functors induced by F

$$F_{A,A} : P(A, A) \to P'(F(A), F(A)),$$

$$F_{B,B} : P(B, B) \to P'(F(B), F(B)),$$

$$F_{B,A} : P(B, A) \to P'(F(B), F(A))$$

and

$$F_{A,B}: P(A,B) \to P'(F(A),F(B))$$

are Thomason equivalences of categories (Corollary 6.3.10). For the first two ones, this follows trivially from the fact that the categories P(A, A), P'(A', A'), P(B, B) and P'(B', B') are all isomorphic to \mathbb{D}_0 . For the third one, this follows trivially from the fact that the categories P(B, A) and P'(B', A') are the empty category. For the fourth one, this can be seen as follows. The category P(A, B) is the free category on the graph

$$f \xrightarrow[\beta]{\alpha} g$$

(2-cells of P become 1-cells of P(A, B) and 1-cells of P become 0-cells of P(A, B)) and the category P'(F(A), F(B)) = P'(A', B') is the free category on the graph

$$\stackrel{\gamma}{\bigcap}$$

 $h.$

The functor $F_{A,B}$ comes from a morphism of reflexive graphs and is obtained by "killing the generator β " (see Example 6.1.16). In particular, the square

$$\mathbb{D}_{1} \xrightarrow{\langle \beta \rangle} P(A, B)$$

$$\downarrow \qquad \qquad \qquad \downarrow^{F_{A,B}}$$

$$\mathbb{D}_{0} \xrightarrow{\langle h \rangle} P'(A', B')$$

is Thomason homotopy cocartesian and thus, $F_{(A,B)}$ is a Thomason equivalence.

Now consider (a copy of) \mathbb{S}_2 labelled as follows:



and let $G:\mathbb{S}_2\to P'$ be the unique 2-functor such that

-
$$G(\overline{A}) = A'$$
 and $G(\overline{B}) = B'$,

$$- G(i) = G(j) = h,$$

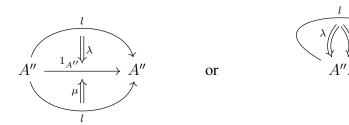
-
$$G(\delta) = \gamma$$
 and $G(\epsilon) = 1_h$.

For similar reasons as for F, the 2-functor G is a Thomason equivalence. This proves that both P' and P have the homotopy type of \mathbb{S}_2 .

Now, let P'' be the free 2-category defined as follows:

- generating 0-cell: A'',
- generating 1-cell: $l: A'' \to A''$,
- generating 2-cells: $\lambda : l \Rightarrow 1_{A''}$ and $\mu : l \Rightarrow 1_{A''}$.

In pictures, this gives



Let $H: \mathbb{S}_2 \to P''$ be the unique 2-functor such that:

- $H(\overline{A}) = H(\overline{B}) = A'',$
- H(i) = l and $H(j) = 1_{A''}$,
- $H(\delta) = \lambda$ and $H(\epsilon) = \mu$.

Let us prove that H is a Thomason equivalence using Corollary 6.3.11. In order to do so, we have to compute $V_k(H) : V_k(\mathbb{S}_2) \to V_k(P'')$ for every $k \ge 0$. For k = 0, the category $V_0(\mathbb{S}_2)$ is the free category on the graph

$$\overline{A} \xrightarrow{i}_{j} \overline{B},$$

the category $V_0(P'')$ is the free category on the graph

and
$$V_0(H)$$
 comes from a morphism of reflexive graphs obtained by "killing the generator j ". Hence, it is a Thomason equivalence of categories. For $k > 0$, the category $V_k(\mathbb{S}_2)$ has two objects \overline{A} and \overline{B} and an arrow $\overline{A} \to \overline{B}$ is a k -tuple of one of the following forms:

 \widehat{A}

(1_i, · · · , 1_i, δ, 1_j, · · · , 1_j),
(1_i, · · · , 1_i, ε, 1_j, · · · , 1_j),
(1_i, · · · , 1_i),
(1_j, · · · , 1_j),

and these are the only non-trivial arrows. In other words, $V_k(\mathbb{S}_2)$ is the free category on the graph with two objects and 2k + 2 parallel arrows between these two objects. In order to compute $V_k(P'')$, let us first notice that every 2-cell of P'' (except for $1_{A''}^{(2)}$) is uniquely encoded as a finite word on the alphabet that has three symbols : 1_l , λ and μ . Concatenation corresponding to the 0-composition of these cells. This means exactly that $V_1(P'')$ is free on the graph that has one object and three arrows. More generally, it is a tedious but harmless exercise to prove that for every k > 0, the category $V_k(P'')$ is the free category on the graph that has one object A'' and 2k + 1 arrows which are of one of the following forms:

- $(1_l, \cdots, 1_l, \lambda, 1^2_{A''}, \cdots, 1^2_{A''}),$ - $(1_l, \cdots, 1_l, \mu, 1^2_{A''}, \cdots, 1^2_{A''}),$ - $(1_l, \cdots, 1_l)$.

Once again, the functor $V_k(H)$ comes from a morphism of reflexive graphs and is obtained by "killing the generator $(1_j, \dots, 1_j)$ ". Hence, it is a Thomason equivalence and thus, so is H. This proves that P'' has the homotopy type of \mathbb{S}_2 .

Finally, consider the commutative diagram of $ho(\mathbf{Ch}_{\geq 0})$

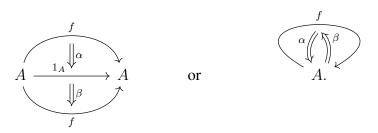
$$\mathbb{H}^{\mathrm{Sing}}(P) \xrightarrow{\mathbb{H}^{\mathrm{Sing}}(F)} \mathbb{H}^{\mathrm{Sing}}(P') \xleftarrow{\mathbb{H}^{\mathrm{Sing}}(G)} \mathbb{H}^{\mathrm{Sing}}(\mathbb{S}_{2}) \xrightarrow{\mathbb{H}^{\mathrm{Sing}}(H)} \mathbb{H}^{\mathrm{Sing}}(P'') \downarrow_{\pi_{P}} \qquad \qquad \downarrow_{\pi_{P'}} \qquad \qquad \downarrow_{\pi_{\mathbb{S}_{2}}} \qquad \qquad \downarrow_{\pi_{P''}} \\ \mathbb{H}^{\mathrm{pol}}(P) \xrightarrow{\mathbb{H}^{\mathrm{pol}}(F)} \mathbb{H}^{\mathrm{pol}}(P') \xleftarrow{\mathbb{H}^{\mathrm{pol}}(G)} \mathbb{H}^{\mathrm{pol}}(\mathbb{S}_{2}) \xrightarrow{\mathbb{H}^{\mathrm{pol}}(H)} \mathbb{H}^{\mathrm{pol}}(P'').$$

Since F, G and H are Thomason equivalences, the three top horizontal morphisms are isomorphisms. Besides, a simple computation using Proposition 4.2.8, which we leave to the reader, shows that the three bottom horizontal morphisms are also isomorphisms. Since S_2 is homologically coherent (Proposition 5.2.3), the morphism π_{S_2} is an isomorphism. This implies that $\pi_{P'}$, $\pi_{P'}$ and $\pi_{P''}$ are isomorphisms, which means by definition that P, P' and P'' are homologically coherent.

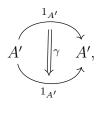
6.5.3. Let *P* be the free 2-category defined as follows:

- generating 0-cell: A,
- generating 1-cell: $f : A \to A$,
- generating 2-cells: $\alpha : f \Rightarrow 1_A$ and $\beta : 1_A \Rightarrow f$.

In pictures, this gives



Now, let P' be a copy of $B^2\mathbb{N}$ labelled as follows



which can be also pictured as

$\widehat{\bigwedge}_{A'.}$

Let $F: P \to P'$ be the unique 2-functor such that:

-
$$F(A) = A'$$
,
- $F(f) = 1_{A'}$
- $F(\alpha) = 1_{A'}^{(2)}$ and $F(\beta) = \gamma$,

and $G: P' \to P$ be the unique 2-functor such that:

-
$$G(A') = A$$
,
- $G(\gamma) = \alpha * \beta$.

Notice that we have $F \circ G = id_{P'}$ and that we have an oplax transformation

$$h: \operatorname{id}_P \Rightarrow G \circ F$$

defined as

-
$$h_A := 1_A$$
.
- $h_f := \alpha$.

Hence, F is a Thomason equivalence (Proposition 3.4.4) and P has the homotopy type of a $K(\mathbb{Z}, 2)$ (see 4.5.2). In particular, it has non-trivial singular homology groups in all even dimension; but since it is a free 2-category, all its polygraphic homology groups are trivial strictly above dimension 2, which means that P is *not* homologically coherent.

6.5.4. All the results from 6.5.2 and 6.5.3 are summed up by the following table.

2-category	homologically coherent?	homotopy type
	yes	\mathbb{S}_2
	yes	\mathbb{S}_2
$\bullet \xrightarrow{\bigcirc} \bullet$	yes	\mathbb{S}_2
	yes	\mathbb{S}_2
	no	$K(\mathbb{Z},2)$
	no	$K(\mathbb{Z},2)$

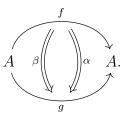
Notice that the fourth and fifth entries of this table only differ by the direction of a generating 2-cell but the homotopy types are not the same.

Let us now move on to bouquets of spheres.

6.5.5. Let *P* be the free 2-category defined as follows:

- generating 0-cell: A,
- generating 1-cells: $f, g : A \rightarrow A$,
- generating 2-cells: $\alpha, \beta: f \Rightarrow g$.

In pictures, this gives:



Now let \mathbb{S}_2 be labelled as follows:



Notice that P is obtained as the following amalgamated sum:

Let us prove that this square is Thomason homotopy cocartesian using the first part of Proposition 6.3.14. This means that we have to show that the induced square of Cat

$$\begin{array}{cccc}
V_k(\mathbb{S}_0) & \stackrel{V_k(\langle C, D \rangle)}{\longrightarrow} & V_k(\mathbb{S}_2) \\
& \downarrow & & \downarrow \\
V_{k}(\mathbb{D}_0) & \longrightarrow & V_k(P)
\end{array}$$
(6.4)

is Thomason homotopy cocartesian for every $k \ge 0$. Notice first that we trivially have $V_k(\mathbb{S}_0) \simeq \mathbb{S}_0$ and $V_k(\mathbb{D}_0) \simeq \mathbb{D}_0$ for every $k \ge 0$ and that $V_0(\mathbb{S}_2)$ is the free category on the graph

$$C \xrightarrow{f} D.$$

Besides, $V_0(P)$ is the free category on the graph

$$\begin{array}{c} f \\ \bigcirc \\ A. \\ \bigcirc \\ g \end{array}$$

In particular, square (6.4) is cocartesian for k = 0 and we are in the situation of identification of two objects of a free category (see Example 6.1.12). Hence, square (6.4) is Thomason cocartesian for k = 0. Similarly, for k > 0, we have already seen that $V_k(\mathbb{S}_2)$ is the free category on the graph that has 2 objects and 2k + 2 parallel arrows between these two objects and we leave as an easy exercise to the reader to check that the category $V_k(P)$ is the free category on the graph that has one object and 2k + 2 arrows, which are the k-tuples of one of the following forms:

- $(1_f, \cdots, \alpha, \cdots, 1_g)$,
- $(1_f, \cdots, \beta, \cdots, 1_g),$

-
$$(1_f, \cdots, 1_f),$$

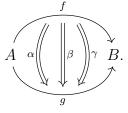
-
$$(1_g, \cdots, 1_g)$$
.

In particular, square (6.4) is again a cocartesian square of identification of two objects of a free category, and thus, it is Thomason homotopy cocartesian. This implies that square (6.3) is Thomason homotopy cocartesian. Since \mathbb{S}_0 , \mathbb{D}_0 and \mathbb{S}_2 are homologically coherent and since $\langle C, D \rangle : \mathbb{S}_0 \to \mathbb{S}_2$ is a folk cofibration, this proves that P is homologically coherent and has the homotopy type of the bouquet of a 1-sphere with a 2-sphere.

6.5.6. Let *P* be the free 2-category defined as follows:

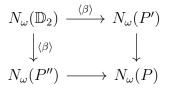
- generating 0-cells: A and B,
- generating 1-cells: $f, g: A \rightarrow B$,
- generating 2-cells: $\alpha, \beta, \gamma : f \Rightarrow g$.

In pictures, this gives



Now let P' be the sub-2-category of P spanned by A, B, f, g, α and β , and let P'' be the sub-2-category of P spanned by A, B, f, g, β and γ . These 2-categories are simply copies of \mathbb{S}_2 . Notice that we have a cocartesian square

and by reasoning as in the proof of Lemma 5.2.2, one can show that the square induced by the nerve

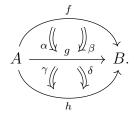


is also cocartesian. Since $\langle \beta \rangle : \mathbb{D}_2 \to P'$ and $\langle \beta \rangle : \mathbb{D}_2 \to P''$ are monomorphisms and N_{ω} preserves monomorphisms, it follows from Lemma 2.5.12 that square (6.5) is Thomason homotopy cocartesian and in particular that P has the homotopy type of a bouquet of two 2-spheres. Since \mathbb{D}_2 , P' and P'' are free and homologically coherent and since $\langle \beta \rangle : \mathbb{D}_2 \to P'$ and $\langle \beta \rangle : \mathbb{D}_2 \to P''$ are folk cofibrations, this also proves that P is homologically coherent (see 6.4.1).

6.5.7. Let *P* be the free 2-category defined as follows:

- generating 0-cells: A and B,
- generating 1-cells: $f, g, h : A \rightarrow B$,
- generating 2-cells: $\alpha, \beta : f \Rightarrow g$ and $\delta, \gamma : g \Rightarrow h$.

In pictures, this gives:



Let us prove that this 2-category is homologically coherent. Let P_0 be the sub-1-category of P spanned by A, B and g, let P_1 be the sub-2-category of P spanned by A, B, g, h, γ and δ and let P_2 be the sub-2-category of P spanned by A, B, f, g, α and β . The 2-categories P_1 and P_2 are copies of \mathbb{S}_2 , and P_0 is a copy of \mathbb{D}_1 . Moreover, we have a cocartesian square of inclusions

Let us prove that this square is Thomason homotopy cocartesian using the second part of Corollary 6.3.14. This means that we have to show that for every $k \ge 0$, the induced square of Cat

is Thomason homotopy cocartesian. For k = 0, this is obvious since all of the morphisms of square (6.6) are isomorphisms at the level of objects and the functor S_0 is the functor that sends a 2-category to its set of objects (seen as a discrete category). Now, notice that the categories $P_i(A, A)$ and $P_i(B, B)$ for $0 \le i \le 2$ are all isomorphic to the terminal category \mathbb{D}_0 and the categories $P_i(B, A)$ for $0 \le i \le 2$ are all the empty category. It follows that for k > 0, we have

$$S_k(P_i) \simeq \mathbb{D}_0 \amalg \left(\prod_{E_k} P_i(A, B) \right) \amalg \mathbb{D}_0$$

where E_k is the set of all k-tuples of the form

$$(A, \cdots, A, B, \cdots, B).$$

The set E_1 is empty and thus all of the morphisms of square (6.7) for the value k = 1 are isomorphisms. This makes square (6.7) Thomason homotopy cocartesian for k = 1. For k > 1, notice that the categories P(A, B), $P_2(A, B)$ and $P_1(A, B)$ are respectively free on the graphs

$$f \xrightarrow{\alpha}_{\beta} g \xrightarrow{\gamma}_{\delta} h$$
$$f \xrightarrow{\alpha}_{\beta} g,$$

and

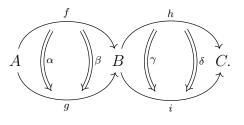
$$g \xrightarrow[]{\gamma}{\longrightarrow} h.$$

It is then straightforward to check that we are in a situation where Corollary 6.1.9 applies and thus square (6.7) is Thomason homotopy cocartesian for $k \ge 1$. Altogether, this proves that square (6.6) is Thomason homotopy cocartesian. Since the inclusions $P_0 \hookrightarrow P_1$ and $P_0 \hookrightarrow P_2$ are folk cofibrations and since P_0 , P_1 and P_2 are homologically coherent, this proves that P is homologically coherent and has the homotopy type of a bouquet of two 2-spheres.

6.5.8. Let *P* be the free 2-category defined as follows:

- generating 0-cells: A, B and C,
- generating 1-cells: $f, g : A \to B$ and $h, i : B \to C$,
- generating 2-cells: $\alpha, \beta : f \Rightarrow g$ and $\gamma, \delta : h \Rightarrow i$.

In pictures, this gives:



Let us prove that P is homologically coherent. Let P' be the sub-2-category of P spanned by A, B, f, g, α and β and let P" be the sub-2-category of P spanned by B, C, h, i, γ and δ . These two 2-categories are copies of \mathbb{S}_2 and we have a cocartesian square

where the anonymous arrows are the canonical inclusions. Let us prove that this cocartesian square is Thomason homotopy cocartesian using the first part of Proposition 6.3.14. This means that we need to prove that for every $k \ge 0$, the induced square of Cat

is Thomason homotopy cocartesian. For every $k \ge 0$, we have $V_k(\mathbb{D}_0) \simeq \mathbb{D}_0$ and for k = 0 the categories $V_0(P')$, $V_0(P'')$ and $V_0(P)$ are respectively free on the graphs

$$A \xrightarrow[g]{f} B,$$
$$B \xrightarrow[i]{i} C,$$

and

$$A \xrightarrow[g]{f} B \xrightarrow[i]{h} C.$$

This implies that square (6.9) is cocartesian for k = 0 and in virtue of Corollary 6.1.9 it is also Thomason homotopy cocartesian for this value of k. For k > 0, since P' and P" are both isomorphic to \mathbb{S}_2 , we have already seen in 6.5.2 that $V_k(P')$ and $V_k(P'')$ are (isomorphic) to the free category on the graph that has two objects and 2k + 2 parallel arrows between these two objects. Similarly, the category $V_k(P)$ is free on the graph that has three objects A, B, C, whose arrows from A to B are k-tuples of one of the following form

(1_f,..., 1_f, α, 1_g,..., 1_g),
(1_f,..., 1_f, β, 1_g,..., 1_g),
(1_f,..., 1_f),

- $(1_g, \cdots, 1_g),$

whose arrows from B to C are k-tuple of one of the following form

- $(1_h, \cdots, 1_h, \gamma, 1_i, \cdots, 1_i)$,
- $(1_h, \cdots, 1_h, \delta, 1_i, \cdots, 1_i)$,
- $(1_h, \cdots, 1_h)$,
- $(1_i, \cdots, 1_i),$

and with no other arrows. This implies that square (6.9) is cocartesian for every k > 0and in virtue of Corollary 6.1.9 it is also Thomason homotopy cocartesian for these values of k. Altogether, this proves that square (6.8) is Thomason homotopy cocartesian. Hence, P is homologically coherent and has the homotopy type of a bouquet of two 2-spheres.

Let us end this section with an example of a 2-category that has the homotopy type of the torus.

6.5.9. Let *P* be the free 2-category defined as follows:

- generating 0-cell: A,
- generating 1-cells: $f, g : A \to A$,
- generating 2-cell: $\alpha : g * f \Rightarrow f * g$.

In pictures, this gives:

$$\begin{array}{ccc} A & \stackrel{f}{\longrightarrow} & A \\ g \downarrow & \stackrel{\alpha}{\longrightarrow} & \downarrow g \\ A & \stackrel{f}{\longrightarrow} & A. \end{array}$$

From now on, we will use concatenation instead of the symbol $*_0$ for the 0-composition. For example, fg will stand for f*g. With this notation, the set 1-cells of P is canonically isomorphic to the set of finite words in the alphabet $\{f, g\}$ and the set of 2-cells of P is canonically isomorphic to the set of finite words in the alphabet $\{f, g, \alpha\}$. For a 1-cell w such that f appears n times in w and g appears m times in w, it is a simple exercise left to the reader to show that there exists a unique 2-cell of P from w to the word

$$f \cdots f g \cdots g$$

where f is repeated n times and g is repeated m times. Recall that we write $B^1(\mathbb{N} \times \mathbb{N})$ for the monoid $\mathbb{N} \times \mathbb{N}$ considered as a category with only one object, and let

$$F: P \to B^1(\mathbb{N} \times \mathbb{N})$$

be the unique 2-functor such that:

This last equation makes sense since (1, 1) = (0, 1) + (1, 0) = (1, 0) + (0, 1). For every 1-cell w of P (encoded as a finite words in the alphabet $\{f, g\}$) such that f appears n times and g appears m times, we have F(w) = (n, m). Let us prove that F is a Thomason equivalence using a dual of [AM20c, Corollaire 5.26] (see Remark 5.20 of op. cit.). If we write \star for the only object of $B^1(\mathbb{N} \times \mathbb{N})$, what we need to show is that the canonical 2-functor from P/\star (see 3.7.3) to the terminal 2-category

$$P/\star \to \mathbb{D}_0$$

is a Thomason equivalence. The 2-category P/\star is described as follows:

- A 0-cell is a 1-cell of $B^1(\mathbb{N} \times \mathbb{N})$.
- For (n,m) and (n',m') two 0-cells of P/\star , a 1-cell from (n,m) to (n',m') is a 1-cell w of P such that the triangle

$$\star \xrightarrow{F(w)} \star \\ (n,m) \xrightarrow{} \star \swarrow (n',m')$$

is commutative. More explicitly, if F(w) = (n'', m''), the commutativity of the previous triangle means

$$n' + n'' = n$$
 and $m' + m'' = m$.

- Given two parallel 1-cells w and w' of P/\star , a 2-cell of P/\star from w to w' is simply a 2-cell of P from w to w' seen as 1-cells of P.

From what we said earlier on the 1-cells and 2-cells of P, it follows easily that for every 0-cell (n, m) of P/\star , the category

$$P/\star((m,n),(0,0))$$

has a terminal object, which is given by

$$f \cdots f g \cdots g$$

where f is repeated n times and g is repeated m times. Then, it follows from [AM20c, Théorème 5.27 and Remarque 5.28] that $P/\star \to \mathbb{D}_0$ is a Thomason equivalence of 2-categories and this proves that F is a Thomason equivalence. Since $B^1(\mathbb{N} \times \mathbb{N}) \simeq$ $B^1(\mathbb{N}) \times B^1(\mathbb{N})$ and $B^1(\mathbb{N})$ has the homotopy type of \mathbb{S}_1 , we conclude that P has the homotopy type of $\mathbb{S}_1 \times \mathbb{S}_1$, i.e. the homotopy type of the torus.

Consider now the commutative square

$$\mathbb{H}^{\operatorname{Sing}}(P) \xrightarrow{\mathbb{H}^{\operatorname{Sing}}(F)} \mathbb{H}^{\operatorname{Sing}}(B^{1}(\mathbb{N} \times \mathbb{N}))
 \downarrow_{\pi_{P}} \qquad \qquad \qquad \downarrow_{\pi_{B^{1}(\mathbb{N} \times \mathbb{N})}}
 \mathbb{H}^{\operatorname{pol}}(P) \xrightarrow{\mathbb{H}^{\operatorname{pol}}(F)} \mathbb{H}^{\operatorname{pol}}(B^{1}(\mathbb{N} \times \mathbb{N}))$$

Since F is a Thomason equivalence, the top horizontal arrow is an isomorphism and since $B^1(\mathbb{N} \times \mathbb{N})$ is a 1-category, it is homologically coherent (Theorem 5.3.14), which means that the right vertical arrow is an isomorphism. The 1-category $B^1(\mathbb{N} \times \mathbb{N})$ is not free but since it has the homotopy type of the torus, we have

$$H_k^{\operatorname{Sing}}(B^1(\mathbb{N}\times\mathbb{N})) = 0 = H_k^{\operatorname{pol}}(B^1(\mathbb{N}\times\mathbb{N}))$$

for $k \ge 2$ and it follows then from Corollary 4.6.17 and Paragraph 4.6.18 that the map canonical map

$$\alpha_{B^1(\mathbb{N}\times\mathbb{N})}^{\mathrm{pol}}:\mathbb{H}^{\mathrm{pol}}(B^1(\mathbb{N}\times\mathbb{N}))\to\lambda(B^1(\mathbb{N}\times\mathbb{N}))$$

is a quasi-isomorphism. Since P is free, it follows that the map $\mathbb{H}^{\text{pol}}(F)$ can be identified with the image in $\text{ho}(\mathbf{Ch}_{>0})$ of the map

$$\lambda(F):\lambda(P)\to\lambda(B^1(\mathbb{N}\times\mathbb{N})),$$

which is easily checked to be a quasi-isomorphism. By a 2-out-of-3 property, we deduce that $\pi_P : \mathbb{H}^{\text{pol}}(P) \to \mathbb{H}^{\text{Sing}}(P)$ is an isomorphism, which means by definition that P is homologically coherent.

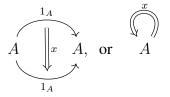
6.6 THE "BUBBLE-FREE" CONJECTURE

Definition 6.6.1. Let C be a 2-category. A *bubble* (in C) is a 2-cell x of C such that:

- x is not a unit,
- $s_0(x) = t_0(x)$,

- $t_1(x) = s_1(x) = 1_{s_0(x)}$.

6.6.2. In pictures, a bubble x is represented as



where $A = s_0(x) = t_0(x)$.

Definition 6.6.3. A 2-category is said to be *bubble-free* if it has no bubbles.

6.6.4. The archetypal example of a 2-category that is *not* bubble-free is $B^2\mathbb{N}$. Another non-bubble 2-category is the one from Paragraph 6.5.3. It is remarkable that of all the free 2-categories we have seen so far, these are the only examples that are non-homologically coherent. This motivates the following conjecture.

Conjecture 6.6.5. *A free* 2-*category is homologically coherent if and only if it is bubblefree.*

6.6.6. At the time of writing, I do not have a real hint towards a proof of the above conjecture. Yet, in light of all the examples seen in the previous section, it seems very likely to be true. Note that we have also conjectured in Paragraph 4.6.24 that for every ω -category C, we have

$$H_2^{\mathrm{pol}}(C) \simeq H_2^{\mathrm{Sing}}(C).$$

If this conjecture on the second homology group is true, then conjecture 6.6.5 may be reformulated as: A free 2-category P has trivial singular homology groups strictly above dimension 2 if and only if it is bubble-free.

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